Computer Science Foundation Exam

May 5, 2006

Computer Science

Section 1A

Name: _____

SSN: _____

Q1		KNW	
Q2		KNW	
Q3		KNW	
Q4		ANL,DSN	
Q5		KNW	
Q6		DSN	
Total			

You have to do all the 6 problems in this section of the exam.
Partial credit cannot be given unless all work is shown and is readable.

Be complete, yet concise, and above all be neat. Do Your rough work on the last page.

1. Do the following problems using summations. All quantities are integers.
[2 + 2 + 3 + 6 pts]

a) What would be the value of *final* after the following segment has been executed?

```
final = 0;
for ( i= 0 ; i <= 50 ; i++)  final += 20;
```

$$\sum_{i=0}^{50} 20$$

= 20 (50 – 0 + 1)

= 20 (51)

= 1020 (2 POINTS)

AWARD ONLY 1 POINT IF SUMMATION STARTS FROM 1 and ANSWER IS 1000

b) What would be the value of *num* after the following segment has been executed?

```
total = 0;
k  = 20;
num = 0;
for ( i= 1 ; i < 5 ; i++) {
     for ( j = 1; j <= 30 ;  j++)
          num   += 6 * k;
          total += 3*j – 6 ;
}
```

num = $\sum_{i=1}^{4} \sum_{j=1}^{30} 6k$

num = $\sum_{i=1}^{4} 180k$

= 720 k

= 720 (20)

= 14400 **(3 POINTS)**

AWARD ONLY 1 POINT IF FIRST SUMMATION GOES UPTO 5 AND RESULT IS 18000

c) What would be the value of *count* after the following segment has been executed?

```
count = 0;
sum = 0;
for ( i= 1 ; i < =5 ;i++) {
     for ( k = 1;   k <= 10 ;  k++)
          count+ = 2 * i;
          sum+ = 3*k – 6 ;
}
```

count = $\sum_{i=1}^{5} \sum_{k=1}^{10} 2i$

$= \sum_{i=1}^{5} 20i$ **AWARD 1 POINT IF THIS IS CORRECT**

= 20 (5)(6)/2

= 300

AWARD ANOTHER 2 POINTS IF THIS PART IS CORRECT OTHERWISE AWARD ZERO

d) Work out an expression for *total* after the following segment has been executed.
Note that n is a even integer (n > 0).

```
total = 0;
for ( k= 1 ; k < 2 ; k++) {
     for ( i = n/2; i <= n ; i++)
          total+ = 16*i - 12*n ;
}
```

First 'for ' loop goes only once, so need to consider it separately

$$\sum_{i=n/2}^{n} 16i \quad -12n$$

$$\sum_{i=n/2}^{n} 16i \quad - \sum_{i=n/2}^{n} 12n$$

AWARD 1 POINT IF SPLIT INTO 2 SUMMATIONS IS PROPER

THE FIRST SUMMATION SHOULD BE SPLIT UP AS SHOWN BELOW.(A COMMON MISTAKE IS TO TAKE IT UPTO n/2)

$$= \sum_{i=1}^{n} 16i \quad - \sum_{i=1}^{n/2-1} 16i \quad -12n[\, n - n/2 + 1]$$

AWARD 3 POINTS IF SECOND SUMMATION GOES UPTO n/2 -1
SOLUTION PROCEEDS AS FOLLOWS IN THAT CASE

$$= \frac{16n(n+1)}{2} \quad - \frac{16(n/2)(n/2\ -1)}{2} \quad -12n(n/2 \quad +1)$$

$$= 8n(n+1) \quad -2(n)(n-2) \quad -6n(n+2)$$

$$= \ 8n^2 + 8n - 2n^2 + 4n - 6n^2 - 12n$$

$$= \ \boldsymbol{0 \ (correct \ answer)}$$

AWARD last 2 POINTS ONLY IF RESULT IS ZERO

IF SECOND SUMMATION has been taken upto n/2 AWARD 1 POINT AT THIS STAGE.
the solution should PROCEED AS FOLLOWS:

$= \frac{16n(n+1)}{2} - \frac{16(n/2)(n/2+1)}{2} - 12n(n/2+1)$

$= 8n(n+1) - 2(n)(n+2) - 6n(n+2)$

$= 8n^2 + 8n - 2n^2 - 4n - 6n^2 - 12n$

= *-8n (answer)*
AWARD last 2 POINTS ONLY IF RESULT NOW IS -8n

2. **[8 pts]** **:1 POINT FOR EACH CORRECT ANSWER**

a) A very large array contains *m* elements sorted in the *reverse order* (largest to smallest). What would be the time complexity in terms of Big-O, when the following sorting algorithms are applied on this array? Quick sort uses first element as the pivot.

Merge sort_____O(n log n)___ Bubble sort_____O(n^2)_____

Quick sort_____ O(n^2)____ Insertion sort____ O(n^2)_____

b) After the array is sorted in proper order using one of the methods mentioned above the following sorting algorithms are applied on the sorted array. Indicate the time complexity in each case.

Bubble sort____O(n)_____ Selection sort ___ O(n^2)_____

Quick sort _____ O(n^2)_____ Merge sort __O(n log n)_____

3. **[5 pts]** Indicate the **worst** case time complexity of following operations in terms of Big-O.
:1 POINT FOR EACH CORRECT ANSWER

a) Searching for a target on a binary tree __ O(n)_____

b) Searching for a target on a binary search tree ___ O(n)_____

c) Pre-order traversal of a binary search tree ___ O(n)_____

d) Counting leaves of a binary tree ___ O(n)_____

e) Finding height of a binary search tree ___ O(n)_____

4. **[8 pts]** Write a recursive function *int countbin (int D)*, which counts the number of ONES in the binary representation of a given decimal integer D. Work out the time complexity of the function in terms of Big-O.

```
int countbin ( int D)
{
     if (D==0)
          return 0;
else
     return (n%2)+countbin(D/2);

}                                                    :4 POINTS
                              (PARTIAL POINTS MAY BE AWARDED )
```

Recurrence relation:

$T(n) = T(n/2) + 3$
$T(1) = 3$ **:1 POINT**

$T(n/2) = T(n/4) + 3$

$T(n) = T(n/4) + 3.2$

$T(n/4) = T(n/8) + 3$

$T(n) = T(n/8) + 3.3$
Or
$T(n) = T(n/2^3) + 3.3$ **:1 POINT**

General case
$T(n) = T(n/2^k) + 3.k$

Let $n = 2^k$

Which yields $k = \log n$ **:1 POINT**

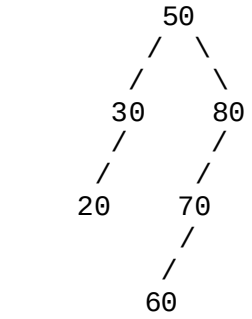
$T(n) = T(1) + 3 \log n$

$= O(\log n)$ **:1 POINT**

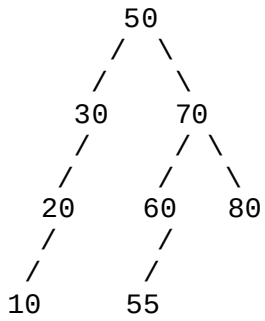
$= O(\log D)$

5. **[8 pts]** a) Draw a binary search tree with the nodes arriving in the order shown below. After every insertion, if the tree gets unbalanced, ***redraw*** the balanced AVL tree.

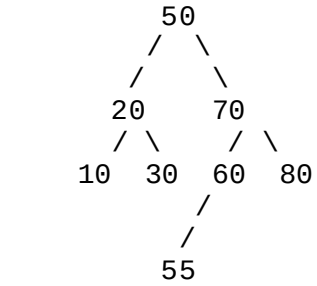
50, 80, 30, 20, 70, 60, 55, 10



:2 POINTS



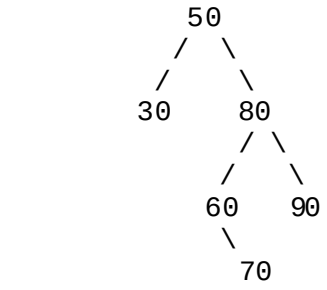
:2 POINTS

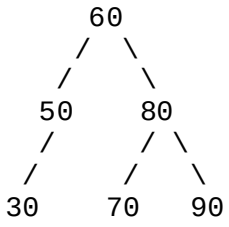


:1 POINT

b) Draw a binary search tree with the nodes arriving in the order shown below. After every insertion, if the tree gets unbalanced, ***redraw*** the balanced AVL tree.

50, 30, 80, 90, 60, 70





:3 POINTS
(NO PARTIAL POINTS)

6. **[8 pts]** Write a function which deletes the largest element on a binary search tree. The nodes of the tree have the following structure:

```
struct node {
     struct node * left;
     int data;
     struct node * right;
};
```

```
struct  node  * largest( struct  node *p)

{
// TAKE CARE OF NULL TREE     : 1 POINT
   if (p==NULL)return NULL;
// TAKE CARE OF ROOT BEING THE LARGEST NODE
                              :3 POINTS
   if (p->right==NULL )

      p= p->left;
 // LARGEST NODE MAY HAVE A LEFT SUBTREE
                              : 2 POINTS
   else
    {
      if (p->right->right==NULL)

          p->right= p->right->left;
      else
// PROPER RECURSION TO FIND THE LARGEST NODE
                              : 2 POINT
          p =   largest(p->right);
     }
}
```