

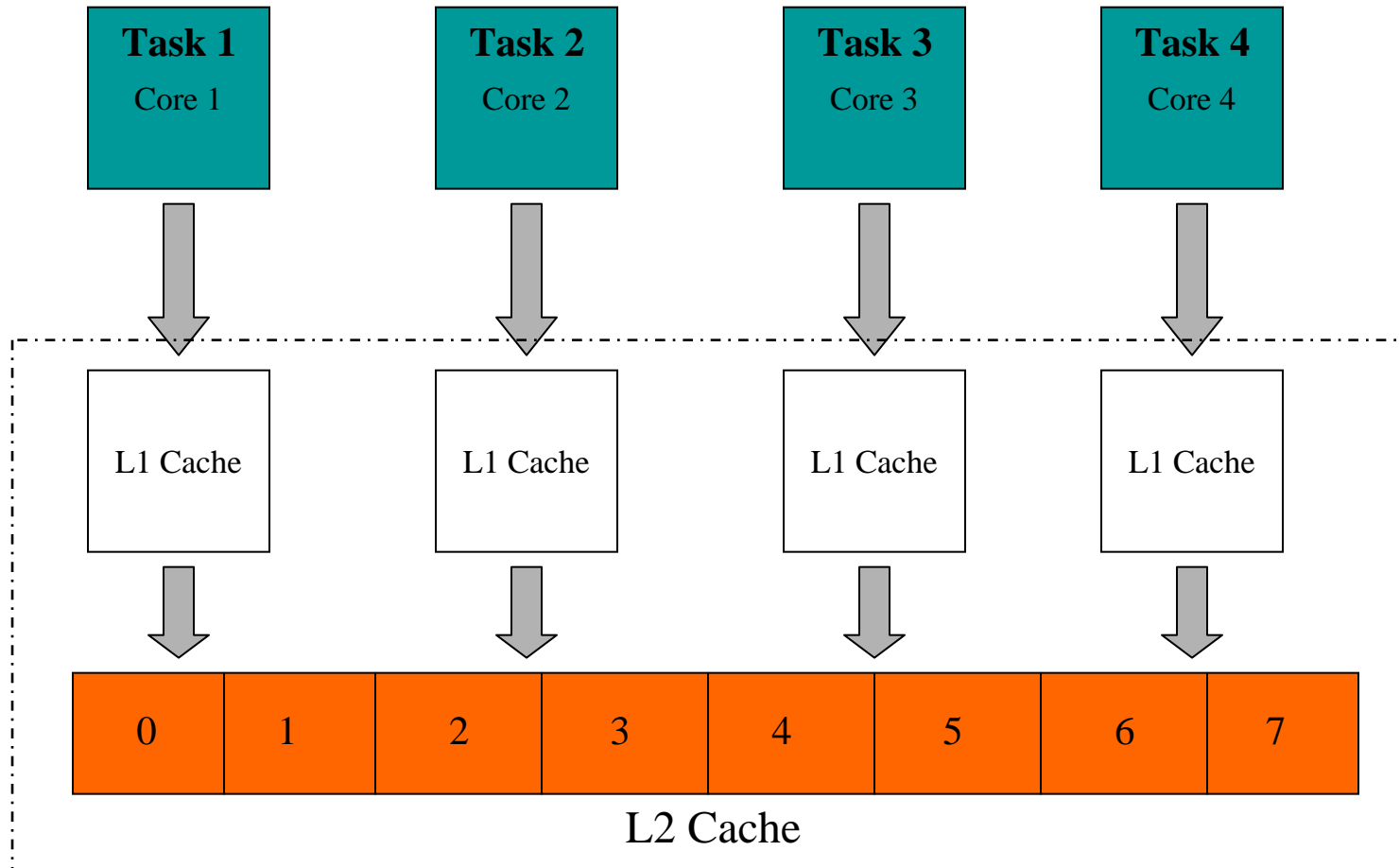


# Cache Partition Problem for Multi-core System

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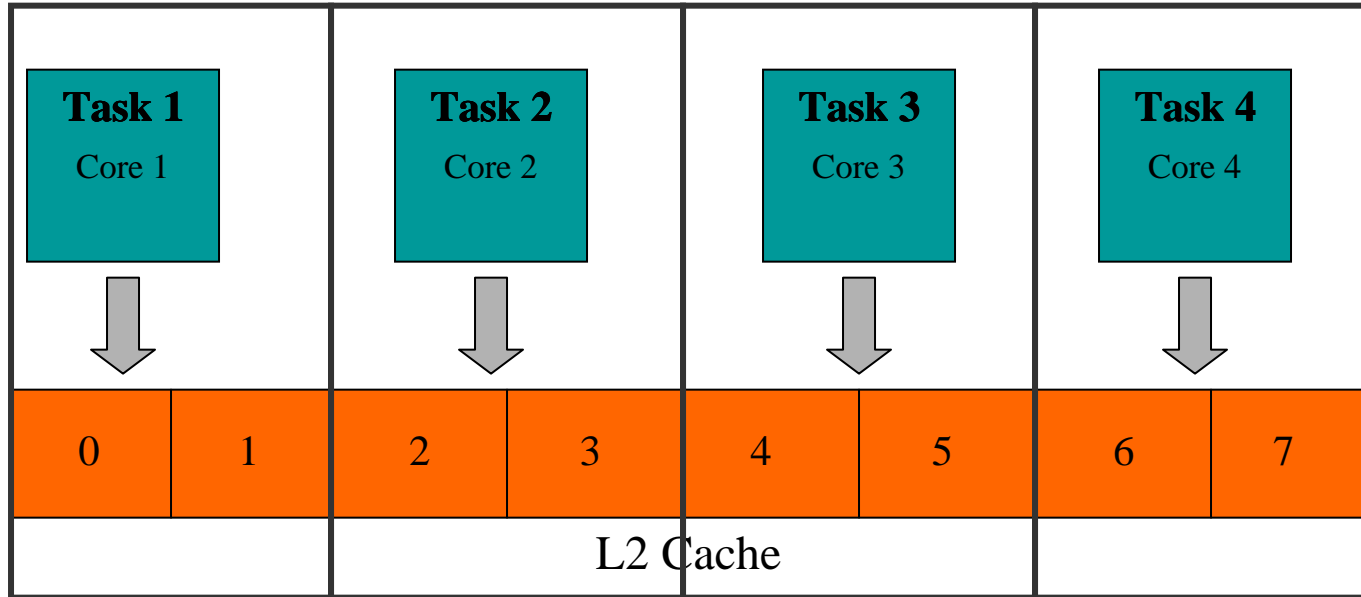
# Introduction to the Multi-core System



Memory System



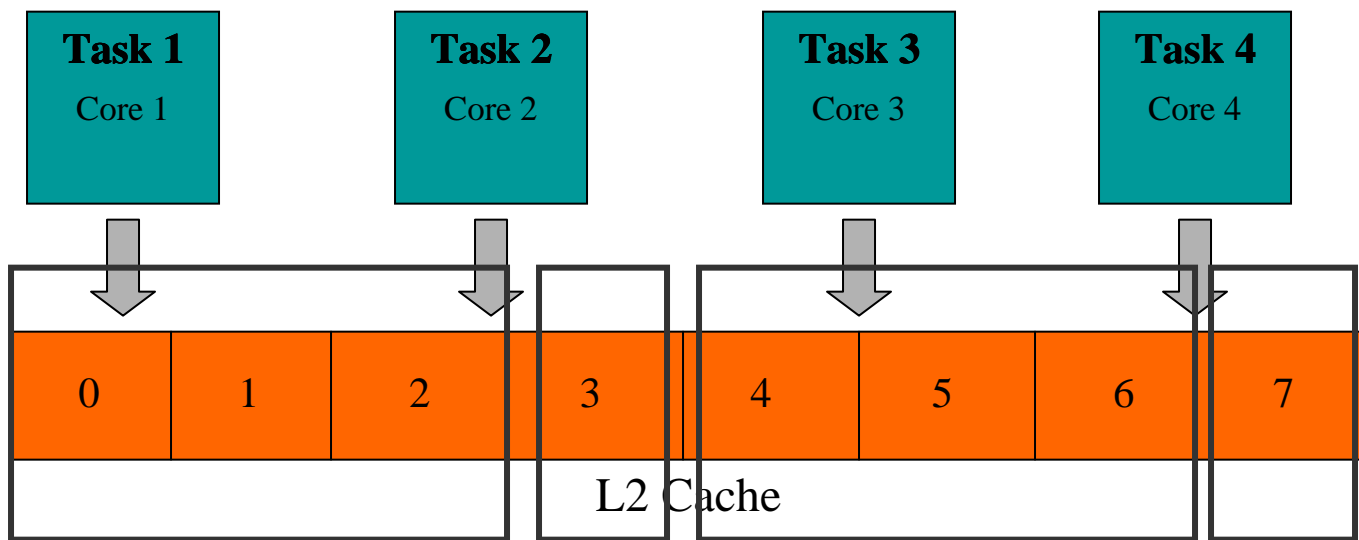
# Introduction to the Multi-core System



- Application Category:
  - Cache-sensitive application
  - Cache-insensitive application



# Introduction to the Multi-core System

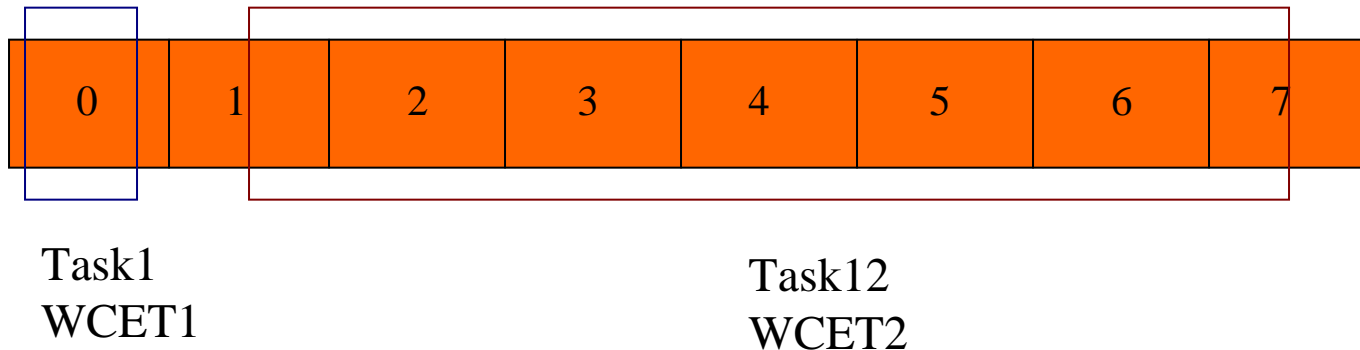


- Property of the task:
  - Task 1 and Task 3 are cache-sensitive
  - Task 2 and Task 4 are cache-insensitive



# Introduction to the Problem

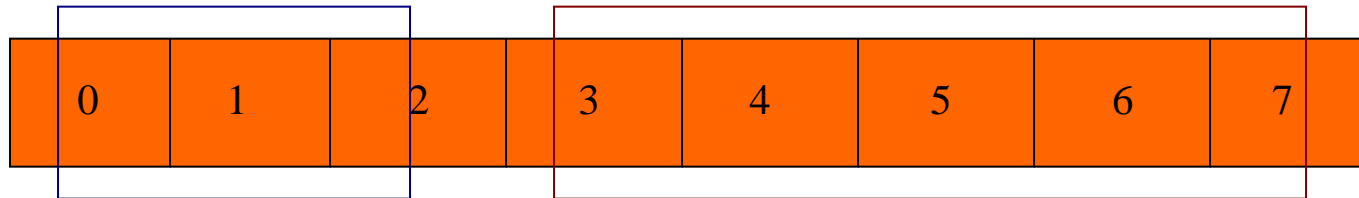
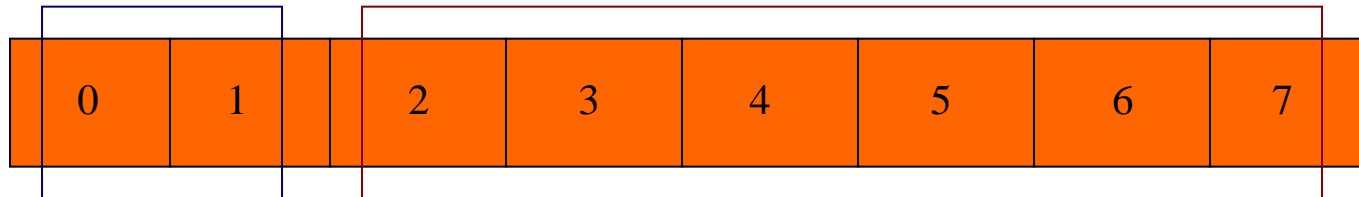
- WCET: Worst Cache Execution Time
- Given the L2 Cache and tasks on each core, is there any partition which makes  $WCET < m$  ? (Good partition strategy) How?
- Assume there are 2 tasks



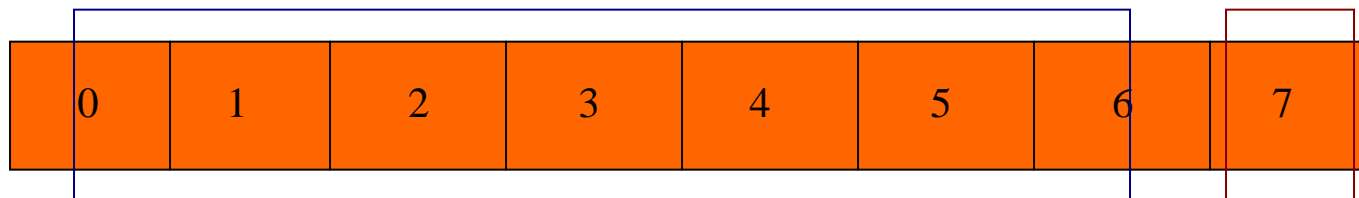
$$WCET1 + WCET2 < m?$$



# Introduction to the Problem



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Task1  
WCET1

Task12  
WCET2

$WCET1 + WCET2 < m$ ? 7 Combinations



# Introduction to the Problem

- 4 Tasks on a 32-way L2 Cache?  
-- 6556 Ways
- 8 Tasks on a 32-way L2 Cache?  
-- 15,380,937 Ways
- Greedy Algorithm



## Problem Definition

- In a Multi-core System, given:  
m-way L2 Cache,  
n tasks on n cores and the number of each task is executed for  $C_i$  times
- Finally there will be n partitions inside the L2 Cache and the corresponding WCET for each partition is  $WCET_1, \dots, WCET_i, \dots, WCET_n$
- $C_1 * WCET_1 + C_2 * WCET_2 + \dots + C_n * WCET_n \leq T$  ?





Subset Sum  Multi-core Cache Partition



## Problem Reformulation

- $s_j \in \{s_1, \dots, s_n\}$ .
- $X_{i,m}$  denote whether a task  $i$  has been assigned to partition  $m$

$$X_{i,1} + X_{i,2} + \dots + X_{i,n} = 1$$

- $x_{1,1} * s_1 + x_{1,2} * s_1 + \dots + x_{2,1} * s_2 + x_{2,2} * s_2 + \dots \leq m$
- $WCET_i = x_{i,1} * WCET_{i,1} + x_{i,2} * WCET_{i,2} + \dots + x_{i,n} * WCET_{i,n}$
- The objective function of the ILP models the WCET of the entire task set for one hyper-period. This overall WCET is defined as:
  - $WCET = c_1 * WCET_1 + c_2 * WCET_2 + \dots + c_n * WCET_n$
- The problem finally is:  
Does  $WCET \leq T$  exist? If so, list all the cases.



# Problem Reformulation

- $x_{i,1} + x_{i,2} + \dots + x_{i,n} = 1$
- $x_{1,1} * s_1 + x_{1,2} * s_1 + \dots + x_{2,1} * s_2 + x_{2,2} * s_2 + \dots \leq m$
- $WCET_i = x_{i,1} * WCET_{i,1} + x_{i,2} * WCET_{i,2} + \dots + x_{i,n} * WCET_{i,n}$
- $WCET = c_1 * WCET_1 + c_2 * WCET_2 + \dots + c_n * WCET_n$
- Does  $WCET \leq T$  exist? If so, list all the cases.

$$WCET = c_1 * WCET_1 + c_2 * WCET_2 + \dots + c_m * WCET_m$$

$$= c_1 * (x_{1,1} * WCET_{1,1} + \dots + x_{1,n} * WCET_{1,n}) + \dots$$

$$+ c_n * (x_{n,1} * WCET_{n,1} + \dots + x_{n,n} * WCET_{n,n})$$

$$= [c_1 * WCET_{1,1} + \dots + c_1 * WCET_{1,n} + \dots + c_n * WCET_{n,1} + \dots + c_n * WCET_{n,n}] * \begin{bmatrix} x_{1,1} \\ \dots \\ x_{1,n} \\ \dots \\ x_{n,1} \\ \dots \\ x_{n,n} \end{bmatrix} \leq T ?$$



# Problem Reformulation

(1)  $WCET = [c_1 * WCET_{1,1} + \dots + c_1 * WCET_{1,n} + \dots + c_n * WCET_{n,1} + \dots + c_n * WCET_{n,n}] *$

(Given:  $X_{i,1} + X_{i,2} + \dots + X_{i,n} = 1$ )

$\mathbf{W}$

$\mathbf{X}$

$$\begin{bmatrix} x_{1,1} \\ \dots \\ 1 - \sum_{i=1}^{n-1} x_{1,i} \\ \dots \\ x_{n,1} \\ \dots \\ 1 - \sum_{i=1}^{n-1} x_{n,i} \end{bmatrix} \leq T ?$$

(2)  $x_{1,1} * s_1 + x_{1,2} * s_1 + \dots + x_{2,1} * s_2 + x_{2,2} * s_2 + \dots \leq m \Rightarrow$

$\mathbf{P}$

$[s_1, \dots, s_1, s_2, \dots, s_2, \dots, s_n, \dots, s_n] *$

$$\begin{bmatrix} x_{1,1} \\ \dots \\ 1 - \sum_{i=1}^{n-1} x_{1,i} \\ x_{2,1} \\ \dots \\ 1 - \sum_{i=1}^{n-1} x_{2,i} \\ \dots \\ x_{n,1} \\ \dots \\ 1 - \sum_{i=1}^{n-1} x_{n,i} \end{bmatrix} \leq m$$



## Problem Reformulation

- Therefore, based on the previous page, Multi-core Cache Partition problem can be reformulated to:  
$$W^T * X \leq T \text{ subject to } P * X \leq m \text{ (} X_i \text{ can only be 0 or 1)}$$
- The format of this problem looks similar to 0-1 Integer Programming, therefore, we can borrow some idea of the transformation from Subset Sum to 0-1 Integer Programming



# Subset Sum to Multi-core Cache Partition

- Multi-core Cache Partition is in the set NP
- Change the Cache Partition problem into slack form:

$$2x_1 - 3x_2 + 3x_3 \geq m?$$

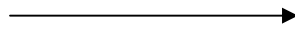
*Constraints :*

$$x_1 + x_2 - x_3 \leq 7;$$

$$-x_1 - x_2 + x_3 \leq -7;$$

$$x_1 - 2x_2 + 2x_3 \leq 4;$$

$$x_1, x_2, x_3 \geq 0$$



$$z = 2x_1 - 3x_2 + 3x_3$$

*Constraints :*

$$x_4 = 7 - x_1 - x_2 + x_3;$$

$$x_5 = -7 + x_1 + x_2 - x_3;$$

$$x_6 = 4 - x_1 + 2x_2 - 2x_3;$$

$$x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$$



# Subset Sum to Multi-core Cache Partition

- Subset Sum:  
INSTANCE: Finite set  $A$ , size  $s(a)$  for each  $a$ ,  
positive integer  $B$ .  
QUESTION: Is there a subset  $A'$  such that the sum of  
the sizes of the elements in  $A'$  is exactly  $B$ ?
- Multi-core Cache Partition  
-- Is there any  $X$  that satisfies  $W^T * X = T$  subject to  $P * X = m$   
( $X_i$  can only be 0 or 1)



## Subset Sum to Multi-core Cache Partition

- Construction:

Input:  $F(S, T)$ :

Output: let  $S = \{S_1, S_2, \dots, S_n\}$ , (Constraint:  $S_1 + \dots + S_n = S$ )

$$T = X_1 * S_1 + X_2 * S_2 + \dots + X_n * S_n$$

- Suppose that  $F\langle S, T \rangle$  is a yes-instance of subset sum.
  - A subset of  $S$  whose elements can sum to  $T$ .
  - Set corresponding  $X_i$
- Suppose that  $\langle s_1 * X_1 + s_2 * X_2 + \dots + s_n * X_n = T \rangle$  is a yes instance of Multi-core Cache Partitioning problem
  - how to set  $X_i$
  - select a subset of  $S$  whose elements sum to  $T$





## References

1. Michael R. Garey, David S. Johnson.  
*Computers and Intractability: A Guide to the Theory of NP-Completeness*, W.H. Freeman and Company, 1979
2. *Utility-Based Cache Partitioning: A Low-Overhead, High-Performance*, Runtime Mechanism to Partition Shared Caches, Moinuddin K. Qureshi and Yale Patt, the 39<sup>th</sup> International Symposium on Microarchitecture, 2006

Thank you and Questions?