# Frame Building Problem

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Stephen Fulwider, Nadeem Mohsin NP-Completeness Proof of FBP

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## Outline



- Description
- Examples



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Description Examples

## Outline



Examples



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# Informal Description

- Need frames to build a greenhouse.
- We have some boards we can use.
- Boards can be cut to make smaller frames.
- Want to minimize extra wood needed.

Description Examples

## Formal Definition: FBP

#### Given

A set F of frames  $\{f_i\}$  and a set B of boards  $\{b_i\} \cup \{E\}$ , where E is the length of the extra board.

#### Question

Does there exist an assignment of every frame to a board such that the sum of the lengths of frames assigned to a board is no greater than the length of the board?

Description Examples

### Formal-er Definition: FBP

#### Given

A set F of frames  $\{f_i\}$  and a set B of boards  $\{b_i\} \cup \{E\}$ , where E is the length of the extra board.

#### Question

Does there exist a total mapping  $M: F \rightarrow B$  such that

$$\forall b_k \in B : \sum_{M(j)=k} f_j \leq b_k$$

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### Examples

### Example 1

$$B = \{20, 10\} \cup \{5\}$$
  
F = \{10, 5, 4, 3, 3, 3\}

### Example 2

$$B = \{20, 10\} \cup \{5\}$$
  
F =  $\{10, 5, 4, 4, 3, 3, 3\}$ 

### Example 3

 $B = \{20, 10\} \cup \{5\}$  $F = \{10, 8, 5, 4, 3, 3, 3\}$ 

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# Failed Approaches

### • Multidimensional Knapsack

- Similar flavor
- $(v_i, w_i)$  versus  $f_i$
- Partial versus total mapping
- Zero-One Integer Programming
  - Define the family of indicator variables  $a_{i,j}$  that indicate whether  $f_i$  is cut from  $b_j$ .

• 
$$\sum_{i} a_{ij} f_i \leq b_j$$
 subject to  $\sum_{i,j} a_{ij} = |F|$ .

# BinPacking[1]

#### Given

Finite set U of items, a size  $s(u) \in Z^+$ for each  $u \in U$ , a positive integer bin capacity B, and a positive integer K.

#### Question

Is there a partition of U into disjoint sets  $U_1, U_2, \ldots, U_k$  such that the sum of the sizes of the items in each  $U_i$  is B or less?

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### Formal Definition: FBP<sub>K</sub>

#### Given

A set F of frames  $\{f_i\}$  and a set B of boards  $\{b_i\} \cup \{E\}$ , where E is the length of the extra board, and an integer K.

#### Question

Does there exist an assignment of every frame to a board such that the sum of the lengths of frames assigned to a board is no greater than the length of the board, using at most K boards?

 $BinPacking \propto FBP_K$ 

- Restrict *FBP<sub>K</sub>* to instances where all boards have the same length.
- $b_1 = b_2 = \cdots = b_{|B|}$
- Boards correspond to bins and frames correspond to items.

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### • Reduction from $FBP_K$ to FBP.

Input to FBP <sub>K</sub>	
• F	
• B	
• E	
• K	

### Input to FBP

- *F*′ = *F*
- $B' = \{ Largest \ K \text{ elements from } B \}$

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• E' =Largest element from B

# $FBP_K \propto FBP$

- $Yes(FBP) \rightarrow Yes(FBP_K)$ 
  - Clearly, this is true, since the created instance of FBP contains exactly K boards.
- $Yes(FBP_K) \rightarrow Yes(FBP)$ 
  - If the solution to  $FBP_K$  uses only elements of B', then this is clearly true.
  - Else,  $\exists b_i, b_j \ni b_i$  is used,  $b_j$  is not used,  $b_i \leq b_j$  and  $b_j \in B'$ .
  - Then every frame assigned to  $b_i$  can be assigned to  $b_j$ .
  - By induction, this transforms any solution to a Yes-instance of  $FBP_K$  to a Yes-instance of FBP.

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### $BinPacking \propto FBP_K \propto FBP$

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