COT 6410 Fall 2010

Generally useful information.

- The notation $\mathbf{z = <x,y>}$ denotes the pairing function with inverses $\mathbf{x = <z>_1}$ and $\mathbf{y = <z>_2}$.
- The minimization notation μ **y [P(…,y)]** means the least **y** (starting at **0**) such that **P(…,y)** is true. The bounded minimization (acceptable in primitive recursive functions) notation μ **y (u≤y≤v) [P(…,y)]** means the least **y** (starting at **u** and ending at **v**) such that **P(…,y)** is true. Unlike the text, I find it convenient to define μ **y (u≤y≤v) [P(…,y)]** to be **v+1**, when no **y** satisfies this bounded minimization.
- The tilde symbol, **~,** means the complement. Thus, set **~S** is the set complement of set **S**, and predicate **~P(x)** is the logical complement of predicate **P(x).**
- A function **P** is a predicate if it is a logical function that returns either **1** (**true**) or **0** (**false**). Thus, **P(x)** means **P** evaluates to **true** on **x**, but we can also take advantage of the fact that **true** is **1** and **false** is **0** in formulas like $\mathbf{y \times P(x)}$, which would evaluate to either **y** (if **P(x)**) or **0** (if **~P(x)**).
- A set **S** is recursive if **S** has a total recursive characteristic function χ_S, such that $\mathbf{x \in S \Leftrightarrow \chi_S(x)}$. Note χ_S is a predicate. Thus, it evaluates to **0** (**false**), if $\mathbf{x \notin S}$.
- When I say a set **S** is re, unless I explicitly say otherwise, you may assume any of the following equivalent characterizations:
  1. **S** is either empty or the range of a total recursive function $\mathbf{f_S}$.
  2. **S** is the domain of a partial recursive function $\mathbf{g_S}$.
- If I say a function **g** is partially computable, then there is an index **g** (I know that's overloading, but that's okay as long as we understand each other), such that $\mathbf{\Phi_g(x) = \Phi(x, g) = g(x)}$. Here $\mathbf{\Phi}$ is a universal partially recursive function.
  Moreover, there is a primitive recursive function **STP**, such that
  STP(x, g, t) is **1** (true), just in case **g**, started on **x**, halts in **t** or fewer steps.
  STP(x, g, t) is **0** (false), otherwise.
  Finally, there is another primitive recursive function **VALUE**, such that
  VALUE(x, g, t) is **g(x)**, whenever **STP(x, g, t)**.
  VALUE(x, g, t) is defined but meaningless if **~STP(x, g, t)**.
- The notation $\mathbf{f(x)\downarrow}$ means that **f** converges when computing with input **x**, but we don't care about the value produced. In effect, this just means that **x** is in the domain of **f**.
- The notation $\mathbf{f(x)\uparrow}$ means **f** diverges when computing with input **x**. In effect, this just means that **x** is **<u>not</u>** in the domain of **f**.
- The **Halting Problem** for any effective computational system is the problem to determine of an arbitrary effective procedure **f** and input **x**, whether or not $\mathbf{f(x)\downarrow}$. The set of all such pairs, $\mathbf{K_0}$, is a classic re non-recursive one.
- The **Uniform Halting Problem** is the problem to determine of an arbitrary effective procedure **f**, whether or not **f** is an algorithm (halts on all input). The set of all such function indices is a classic non re one.
- $\mathbf{A \le_m B}$ (**A** many-one reduces to **B**) means that there exists a total recursive function **f** such that $\mathbf{x \in A \Leftrightarrow f(x) \in B}$. If $\mathbf{A \le_m B}$ and $\mathbf{B \le_m A}$ then we say that $\mathbf{A \equiv_m B}$ (**A** is many-one equivalent to **B**). If the reducing function is 1-1, then we say $\mathbf{A \le_1 B}$ (**A** one-one reduces to **B**) and $\mathbf{A \equiv_1 B}$ (**A** is one-one equivalent to **B**).

COT 6410 Fall 2010 Sample Midterm#1 Name:___SAMPLE____

1. Choosing from among **(REC) recursive**, **(RE) re non-recursive, (coRE) co-re non-recursive**, **(NR) non-re/non-co-re**, categorize each of the sets in a) through d). Justify your answer by showing some minimal quantification of some known recursive predicate.

a.) { f | domain(f) is finite } ____________________

b.) { f | domain(f) is empty } ____________________

c.) { <f,x> | f(x) converges in at most 20 steps } ____________________

d.) { f | domain(f) converges in at most 20 steps for some input x } ____________________

2. Let set **A** be recursive, **B** be re non-recursive and **C** be non-re. Choosing from among **(REC) recursive**, **(RE) re non-recursive**, **(NR) non-re**, categorize the set **D** in each of a) through d) by listing **all** possible categories. No justification is required.

a.) D = ~C ____

b.) D ⊆ A ∪ C ____

c.) D = ~B ____

d.) D = B – A ____

3. Prove that the **Halting Problem** (the set **HALT = K_0 = L_u**) is not recursive (decidable) within any formal model of computation. (Hint: A diagonalization proof is required here.)

4. Using reduction from the known undecidable **HasZero, HZ = { f | ∃x f(x) = 0 }**, show the non-recursiveness (undecidability) of the problem to decide if an arbitrary primitive recursive function **g** has the property **IsZero, Z = { f | ∀x f(x) = 0 }**,. Hint: there is a very simple construction that uses **STP** to do this. **Just giving that construction is not sufficient; you must also explain why it satisfies the desired properties of the reduction**.

5. Define **RANGE_ALL = (f | range(f) = ℵ }**.

a.) Show some minimal quantification of some known recursive predicate that provides an upper bound for the complexity of this set. (Hint: Look at **c.)** and **d.)** to get a clue as to what this must be.)

b.) Use Rice's Theorem to prove that **RANGE_ALL** is undecidable.

c.) Show that **TOTAL** $\leq_m$ **RANGE_ALL**, where **TOTAL = { f |** $\forall y\ \varphi_f(y)\downarrow$ **}**.

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d.) Show that **RANGE_ALL** $\leq_m$ **TOTAL**.

e.) From **a.)** through **d.)** what can you conclude about the complexity of **RANGE_ALL**?

6. This is a simple question concerning Rice's Theorem.

a.) State the strong form of Rice's Theorem. Cover all conditions for it to apply; don't skimp on details.

b.) Describe a set of partial recursive functions whose membership cannot be shown undecidable through Rice's Theorem. What condition is violated by your example?

7. Using the definition that **S** is recursively enumerable iff **S** is either empty or the range of some algorithm $\mathbf{f_S}$ (total recursive function), prove that if both **S** and its complement **~S** are recursively enumerable then **S** is decidable. To get full credit, you must show the characteristic function for **S**, χ_S, in all cases. Be careful to handle the (two) extreme cases. Hint: This is not an empty suggestion.