Sample Question#1. Part a

1. Prove that the following are equivalent

a) S is an infinite recursive (decidable) set.

b) S is the range of a monotonically increasing total recursive function.
Note: f is monotonically increasing means that $\forall x\ f(x+1) > f(x)$.

a) Implies b)

Let $x \in S \Leftrightarrow \chi_S(x)$

Define $f_R(0) = \mu\ x\ \chi_S(x)$; $f_R(y+1) = \mu\ x\ [\chi_S(x)\ \&\&\ (x > f_R\ (y))]$

Clearly, since S is non-empty, it has a least one value and so there exist a smallest value such that $\chi_S(x)$; we will enumerate this as $f_R(0) = \mu\ x\ \chi_S(x)$.

Assume we have enumerated the y-th value in S as $f_R\ (y)$. Since S is infinite, there will be values in S greater than $f_R\ (y)$ and our search $\mu\ x\ [\chi_S(x)\ \&\&\ (x > f_R\ (y))$ will find the next largest value for which $\chi_S(x)$. Thus, inductively, we will enumerate the elements of S in increasing order, as desired.

Sample Question#1 Part b

1. Prove that the following are equivalent

a) S is an infinite recursive (decidable) set.

b) S is the range of a monotonically increasing total recursive function.
Note: f is monotonically increasing means that $\forall x\ f(x+1) > f(x)$.

b) Implies a)

Let S be enumerated by the monotonically increasing algorithm f_S. Define χ_S by

$\chi_S(x) = (f_R\ ((\mu\ z\ [f_R\ (z) \geq x]) == x)$

Clearly, if x is enumerated, it must appear before any values greater than it are enumerated and consequently this is a bounded search to find the first element listed that is at least as large as x. If this element is x, then x is in S, else it is not. The fact that f_R is monotonically increasing makes S infinite. The fact that it has a characteristic function makes it decidable.

Sample Question#2

2. Let A and B be re sets. For each of the following, either prove that the set is re, or give a counterexample that results in some known non-re set.

Let A be semi decided by f_A and B by f_B

a) $A \cup B$: must be re as it is semi-decided by

$$f_{A \cup B}(x) = \exists t\, [stp(f_A, x, t) \,||\, stp(f_B, x, t)]$$

b) $A \cap B$: must be re as it is semi-decided by

$$f_{A \cap B}(x) = \exists t\, [stp(f_A, x, t) \,\&\&\, stp(f_B, x, t)]$$

c) ~A: can be non-re. If ~A is always re, then all re are recursive as any set that is re and whose complement is re is decidable. However, A = K is a non-rec, re set and so ~A is not re.

Sample Question#3

3. Present a demonstration that the *even* function is primitive recursive.
   even(x) = 1 if x is even
   even(x) = 0 if x is odd
   You may assume only that the base functions are prf and that prf's are closed under a finite number of applications of composition and primitive recursion.

DONE in class.

Sample Question#4

4. Given that the predicate **STP** and the function **VALUE** are prf's, show that we can semi-decide
   { f | φ_f evaluates to 0 for some input}
   This can be shown re by the predicate
   {f | $\exists$<x,t> [stp(f,x,t) && value(f,x,t) = 0] }

Sample Question#5

5. Let **S** be an re (recursively enumerable), non-recursive set, and **T** be re, non-empty, possibly recursive set. Let

   E = { z | z = x + y, where x $\in$ S and y $\in$ T }.

   (a) Can **E** be non re? **No as we can let S and T be semi-decided by f_S and f_T, resp., E is then semi-dec. by**
   f_E (z) = $\exists$<x,y,t> [stp(f_S, x, t) && stp(f_T, y, t) &&
   (z = value(f_S, x, t) + value(f_T, y, t))]

   (b) Can **E** be re non-recursive? **Yes, just let T = {0}, then E = S which is known to be re, non-rec.**

   (c) Can **E** be recursive? **Yes, let T = $\aleph$, then E = { x | x ≥ min (S) } which is a co-finite set and hence rec.**

Sample Question#6

6. Assuming **TOTAL** is undecidable, use reduction to show the undecidability of
Incr = { f | $\forall x\ \varphi_f(x+1) > \varphi_f(x)$ }
Let f be arb.
Define $G_f(x) = \varphi_f(x) - \varphi_f(x) + x$
$f \in$ TOTAL iff $\forall x \varphi_f(x)\downarrow$ iff $\forall x\ G_f(x)\downarrow$ iff
$\forall x\ \varphi_f(x) - \varphi_f(x) + x = x$ iff $G_f \in$ Incr

Sample Question#7

7. Let **Incr = { f | $\forall$x, φ_f(x+1)>φ_f(x) }**.
   Let **TOT = { f | $\forall$x, φ_f(x)$\downarrow$ }**.
   Prove that **Incr $\equiv_m$ TOT**. Note Q#6 starts this one.
   Let f be arb.
   Define G_f (x) = $\exists$t[stp(f,x,t) && stp(f,x+1,t) && (value(f,x+1,t) > value(f,x,t))]
   f $\in$ Incr iff $\forall$x φ_f(x+1)>φ_f(x) iff
   $\forall$x G_f (x)$\downarrow$ iff $G_f \in$ TOT

Sample Question#8

8. Let **Incr = { f | $\forall x\ \varphi_f(x+1)>\varphi_f(x)$ }**. Use Rice's theorem to show **Incr** is not recursive.
   Non-Trivial as
   $C_0(x)=0 \notin$ Incr; $S(x)=x+1 \in$ Incr
   Let f,g be arb. Such that $\forall x\ \varphi_f(x)=\varphi_g(x)$
   $f \in$ Incr iff $\forall x\ \varphi_f(x+1)>\varphi_f(x)$ iff
   $\forall x\ \varphi_g(x+1)>\varphi_g(x)$ iff $g \in$ Incr

Sample Question#9

9. Let **S** be a recursive (decidable set), what can we say about the complexity (recursive, re non-recursive, non-re) of **T**, where **T** $\subset$ **S**?

   Nothing. Just let S = $\aleph$, then T could be any subset of $\aleph$. There are an uncountable number of such subsets and some are clearly in each of the categories above.

Sample Question#10

10. Define the pairing function $\mathbf{<x,y>}$ and its two inverses $\mathbf{<z>_1}$ and $\mathbf{<z>_2}$, where if $\mathbf{z = <x,y>}$, then $\mathbf{x = <z>_1}$ and $\mathbf{y = <z>_2}$.

Right out of Notes.

Sample Question#11

11. Assume **A** $\leq_m$ **B** and **B** $\leq_m$ **C**.
    Prove **A** $\leq_m$ **C**.

    Done in class

Sample Question#12

12. Let **P = { f | ∃ x [STP(f, x, x)] }**. Why does Rice's theorem not tell us anything about the undecidability of **P**?

This is not an I/O property as we can have implementations of C_0 that are efficient and satisfy P and others that do not.