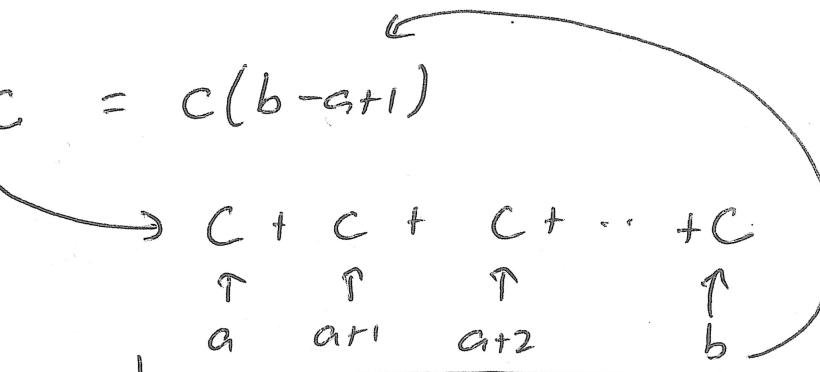


Sums continued  $\rightarrow$  purpose to help determine @ Big-Oh run time of iterative code.

$$\sum_{\substack{i=a \\ a \leq b}}^b c = c(b-a+1)$$



$$\begin{array}{r}
 1 \\
 + 2 \\
 \hline
 3 \\
 + 3 \\
 \hline
 6 \\
 + 4 \\
 \hline
 10 \\
 \vdots
 \end{array}$$

$$S = \sum_{i=1}^n i = \underbrace{(1 + 2 + 3 + \dots + n)}_{f(n) + 1}$$

$$S = \sum_{i=1}^n (n+1-i) = 100 + 99 + \dots + 1$$

$$f(n) + 1$$

$$+ 98$$

$$2S = \sum_{i=1}^n i + (n+1-i) \underbrace{(n+1)(n+1)(n+1)}_{+(n+1)}$$

how many times is  $n+1$  added?

$$\text{Ans} = n$$

$$2S = \sum_{i=1}^n (n+1)$$

$$\frac{n(n+1)}{2}$$

$2 \leftarrow$  each # added twice

$$2S = n(n+1)$$

$$S = \frac{n(n+1)}{2}$$

$1, 3, 6, 10, \dots$

(triangle numbers...)



$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4}$$

$$\sum_{i=a}^b f(i) = \sum_{i=1}^b f(i) - \sum_{i=1}^{a-1} f(i)$$

$a$

$a < i \leq b$

$f(a) + f(a+1) + f(a+2) + \dots + f(b)$      $f(1) + f(2) + \dots + f(a-1) + f(a) + f(a+1) + \dots + f(b)$

$n=3$

$$\begin{aligned}\sum_{i=n+1}^{2n} i &= \sum_{i=1}^{2n} i - \sum_{i=1}^n i \\&= \frac{(2n)(2n+1)}{2} - \frac{n(n+1)}{2} \\&= \frac{n}{2} \left( (2n+1) - (n+1) \right) \\&= \frac{n}{2} \left( 2n+1 - n - 1 \right) \\&= \frac{n}{2} \left( 2n - n \right) = \frac{n^2}{2} \\&= \frac{n}{2} (4n+2 - n - 1) \\&= \frac{n}{2} (3n+1)\end{aligned}$$

~~$$\begin{aligned}\sum_{i=4}^6 i &= 4+5+6=15 \\3 \cdot \frac{6}{2} &= (2 \cdot 3^2 - 1) \\&= \frac{3}{2} (17)\end{aligned}$$~~

$$\sum_{i=a}^b (f(i) + g(i)) = \sum_{i=a}^b f(i) + \sum_{i=a}^b g(i)$$

$\rightarrow f(a) + f(a+1) + f(a+2) + \dots + f(b)$

$\rightarrow g(a) + g(a+1) + g(a+2) + \dots + g(b)$

$f(a) + g(a)$   
 $f(a+1) + g(a+1)$   
 $\vdots$   
 $f(b) + g(b)$

$$\sum_{i=a}^b f(i) \times g(i) \neq \left( \sum_{i=a}^b f(i) \right) \times \left( \sum_{i=a}^b g(i) \right)$$

$$\boxed{\begin{aligned} & f(a)g(a) + \\ & f(a+1)g(a+1) + \\ & f(a+2)g(a+2) + \\ & \vdots \\ & f(b)g(b) \end{aligned}}$$

$$(f(a) + f(a+1) + f(a+2) + \dots + f(b)) \times (g(a) + g(a+1) + \dots + g(b))$$

$$\begin{aligned} & (f(a) + f(a+1))(g(a) + g(a+1)) \\ & = f(a)g(a) + f(a)g(a+1) + f(a+1)g(a) + \\ & \quad f(a+1)g(a+1) \end{aligned}$$

FOIL

$$\boxed{\sum_{i=a}^b cf(i) = c \sum_{i=a}^b f(i)}$$

$$\begin{aligned} \sum_{i=n-1}^{2n+3} (3i - 5) &= \sum_{i=1}^{2n+3} (3i - 5) - \sum_{i=1}^{n-2} (3i - 5) \\ &= \sum_{i=1}^{2n+3} 3i - \sum_{i=1}^{2n+3} 5 - \left( \sum_{i=1}^{n-2} 3i - \sum_{i=1}^{n-2} 5 \right) \\ &= 3 \sum_{i=1}^{2n+3} i - 5(2n+3) - 3 \sum_{i=1}^{n-2} i + 5(n-2) \\ &= \frac{3(2n+3)(2n+4)}{2} - 10n - 15 - \frac{3(n-2)(n-1)}{2} + 5n - 10 \\ &= \frac{3(2n+3)2(n+2)}{2} - \frac{3(n-2)(n-1)}{2} - 5n - 25 \\ &= 3(2n^2 + 7n + 6) - \frac{3}{2}(n^2 - 3n + 2) - 5n - 25 \\ &= \frac{6n^2 + 21n + 18}{2} - \frac{3}{2}n^2 + \frac{9}{2}n - 3 - 5n - 25 \\ &= \frac{9}{2}n^2 + \frac{41}{2}n - 10 \end{aligned}$$

$$S = 3 + \frac{3}{2} + \frac{3}{4} + \frac{3}{8} + \dots \quad \text{Infinite geo sum}$$

$$\begin{array}{rcl} S & = & \cancel{a} + \cancel{ax} + \cancel{ax^2} + \cancel{ax^3} + \dots \\ -xS & = & \cancel{ax} + \cancel{ax^2} + \cancel{ax^3} + \dots \\ \hline S - xS & = & a \end{array} \quad |x| < 1$$

$$S - xS = a$$

$$S(1-x) = a$$

$$S = \frac{a}{1-x}, \text{ infinite geo sum} \quad \frac{\text{1st term}}{1-\text{common ratio}}$$

$$\begin{array}{rcl} S = \sum_{i=0}^{n-1} ax^i & = & (\cancel{a}) + \cancel{ax} + \cancel{ax^2} + \dots \quad \cancel{ax}^{n-1} \\ -xS & = & \cancel{ax} + \cancel{ax^2} + \dots \quad \cancel{ax}^{n-1} + \cancel{ax^n} \end{array}$$

$$S - xS = a - ax^n$$

$$S(1-x) = a(1-x^n)$$

$$S = \frac{a(1-x^n)}{1-x} = \frac{a(x^n - 1)}{x-1}, \quad x \neq 1.$$

```

for (int i = 0; i < n; i++)
    for (int j = 0; j < m; j++)
        sum++;

```

$$\sum_{i=1}^n m = nm$$

$O(nm)$

$$\sum_{i=1}^n \left( \sum_{j=1}^m 1 \right)$$

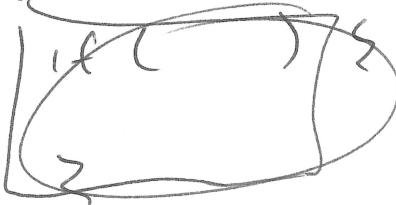
$$= \sum_{i=1}^n m$$

$$= nm$$

```

for (int i = 0; i < n; i++) {
    for (int j = 0; j < i; j++) {

```



Const

}  
}

$$\begin{aligned}
& \sum_{i=0}^{n-1} \left( \sum_{j=0}^{i-1} 1 \right) \\
&= \sum_{i=0}^{n-1} i = \frac{(n-1)n}{2} \\
&= O(n^2)
\end{aligned}$$

$i = n, res = 0$

while ( $n > 0$ ) {

    res = res +  $n \% 2$ ;

$n = n / 2$ ;

}