

## COP 3502 Recitation Sheet: Recurrence Relations Solutions

1) Use the iteration technique to solve the following recurrence relation in terms of  $n$ :

$$T(n) = 2T(n/2) + 1, \text{ for all integers } n > 1$$
$$T(1) = 1$$

Find a tight Big-Oh answer.

### Solution

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$
$$T\left(\frac{n}{2}\right) = 2T\left(\frac{n}{4}\right) + 1$$
$$T(n) = 2\left(2T\left(\frac{n}{4}\right) + 1\right) + 1$$
$$T(n) = 4T\left(\frac{n}{4}\right) + 2 + 1$$
$$T(n) = 4T\left(\frac{n}{4}\right) + 3$$
$$T\left(\frac{n}{4}\right) = 2T\left(\frac{n}{8}\right) + 1$$
$$T(n) = 4\left(2T\left(\frac{n}{8}\right) + 1\right) + 3$$
$$T(n) = 8T\left(\frac{n}{8}\right) + 4 + 3$$
$$T(n) = 8T\left(\frac{n}{8}\right) + 7$$

Based on these three iterations, we see that after  $k$  iterations, the recurrence is

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + (2^k - 1)$$

Plug in the value of  $k$  such that  $\frac{n}{2^k} = 1$  to this recurrence. This means that  $2^k = n$ . Substituting, we get:

$$T(n) = nT(1) + (n - 1)$$
$$T(n) = n + (n - 1)$$
$$T(n) = 2n - 1$$

It follows that  $T(n) = O(n)$ .

2) What is the closed form solution to the following recurrence relation? Please use the iteration technique, show all of your work and provide your final answer in Big-Oh notation.

$$T(1) = 1$$
$$T(n) = 2T(n/4) + 1$$

**Solution**

Iterate the recurrence three times:

$$T(n) = 2T\left(\frac{n}{4}\right) + 1 \quad (\text{one iteration})$$

$$T(n) = 2(2T\left(\frac{n}{16}\right) + 1) + 1$$

$$T(n) = 4T\left(\frac{n}{16}\right) + 3 \quad (\text{two iterations})$$

$$T(n) = 4(2T\left(\frac{n}{64}\right) + 1) + 3$$

$$T(n) = 8T\left(\frac{n}{64}\right) + 7 \quad (\text{three iterations})$$

Now, let's make a guess as to the form of the recurrence after iterating k times based on the first three iterations:

$$T(n) = 2^k T\left(\frac{n}{4^k}\right) + (2^k - 1)$$

Since we know  $T(1)$ , we want to plug in the value of k for which  $\frac{n}{4^k} = 1$ , in for k. Solving, we find that  $n = 4^k$ . Taking the square root of both sides, we find  $\sqrt{n} = \sqrt{4^k} = \sqrt{2^{2k}} = (2^{2k})^{\frac{1}{2}} = 2^k$ . Substituting for both  $4^k$  and  $2^k$ , in the right hand of the recurrence, we get:

$$T(n) = \sqrt{n}T\left(\frac{4^k}{4^k}\right) + (\sqrt{n} - 1) = \sqrt{n}T(1) + (\sqrt{n} - 1) = \sqrt{n} + \sqrt{n} - 1 \in \mathbf{O}(\sqrt{n})$$

3) Use the iteration technique to determine a close form solution for the recurrence relation  $T(n)$  defined below. Note: due to the nature of this recurrence, it's possible to get an exact solution for  $T(n)$ , so please try to do that instead of just getting a Big-Oh bound.

$$T(n) = 2T(n-1) + 2^n$$
$$T(1) = 2$$

### Solution

Iterate as follows:

$$\begin{aligned}T(n) &= 2T(n-1) + 2^n \\&= 2(2T(n-2) + 2^{n-1}) + 2^n \\&= 4T(n-2) + 2^n + 2^n \\&= 4T(n-2) + 2(2^n) \\&= 4(2T(n-3) + 2^{n-2}) + 2(2^n) \\&= 8T(n-3) + 2^n + 2(2^n) \\&= 8T(n-3) + 3(2^n)\end{aligned}$$

In general, after  $k$  iterations, we'll have:

$$= 2^k T(n-k) + k(2^n)$$

Plug in  $k = n-1$  to obtain

$$\begin{aligned}T(n) &= 2^{n-1}T(n-(n-1)) + (n-1)(2^n) \\&= 2^{n-1}T(1) + (n-1)(2^n) \\&= 2^{n-1}(2) + (n-1)(2^n) \\&= 2^n + (n-1)(2^n) \\&= \mathbf{n2^n}\end{aligned}$$

4) Using the iteration technique, find a Big-Oh bound for the following recurrence relation, in terms of n:

$$T(n) = 2T\left(\frac{n}{2}\right) + n^2, \text{ for } n > 1$$

$$T(1) = 1$$

**Solution**

$$T(n) = 2T\left(\frac{n}{2}\right) + n^2$$

$$T(n) = 2\left[2T\left(\frac{n}{4}\right) + \left(\frac{n}{2}\right)^2\right] + n^2$$

$$T(n) = 4T\left(\frac{n}{4}\right) + 2 \times \frac{n^2}{4} + n^2$$

$$T(n) = 4T\left(\frac{n}{4}\right) + \frac{3n^2}{2}$$

$$T(n) = 4\left[2T\left(\frac{n}{8}\right) + \left(\frac{n}{4}\right)^2\right] + \frac{3n^2}{2}$$

$$T(n) = 8T\left(\frac{n}{8}\right) + 4 \times \frac{n^2}{16} + \frac{3n^2}{2}$$

$$T(n) = 8T\left(\frac{n}{8}\right) + \frac{7n^2}{4}$$

In general, after k steps, we see that our formula will iterate to:

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + \frac{(2^k - 1)n^2}{2^{k-1}}$$

Plug in  $\frac{n}{2^k} = 1$ , so let  $n = 2^k$  and  $k = \log_2 n$  to obtain:

$$T(n) = nT(1) + \frac{(n-1)n^2}{n/2}$$

$$T(n) = n + 2n(n-1)$$

$$T(n) = 2n^2 - n = \mathbf{O(n^2)}$$

5) Use the iteration technique to determine a Big-Oh solution for the following recurrence relation:

$$T(n) = 4T\left(\frac{n}{2}\right) + n^2, T(1) = 1$$

**Solution**

$$\begin{aligned}T(n) &= 4T\left(\frac{n}{2}\right) + n^2 \\T(n) &= 4\left(4T\left(\frac{n}{4}\right) + \left(\frac{n}{2}\right)^2\right) + n^2 \\T(n) &= 16T\left(\frac{n}{4}\right) + 4\left(\frac{n^2}{4}\right) + n^2 \\T(n) &= 16T\left(\frac{n}{4}\right) + 2n^2 \\T(n) &= 16\left(4T\left(\frac{n}{8}\right) + \left(\frac{n}{4}\right)^2\right) + 2n^2 \\T(n) &= 64T\left(\frac{n}{8}\right) + 16\left(\frac{n^2}{16}\right) + 2n^2 \\T(n) &= 64T\left(\frac{n}{8}\right) + 3n^2\end{aligned}$$

After k iterations, we guess the form of our recurrence to be:

$$T(n) = 4^k T\left(\frac{n}{2^k}\right) + kn^2$$

We plug in a value of k such that  $\frac{n}{2^k} = 1$ . Namely,  $n = 2^k$  and  $k = \log_2 n$ . Note that  $4^k = (2^2)^k = (2^k)^2 = n^2$ :

$$T(n) = n^2 T(1) + (\log_2 n)n^2$$

Plugging in  $T(1) = 1$ , we have

$$T(n) = n^2 + (\log_2 n)n^2 = \mathbf{O(n^2 \lg n)}$$