

# Using Isabelle to Help Verify Code that Uses Abstract Data Types

Jason Kirschenbaum

The Ohio State University

Bruce M. Adcock

Derek Bronish

Paolo Bucci

Bruce W. Weide

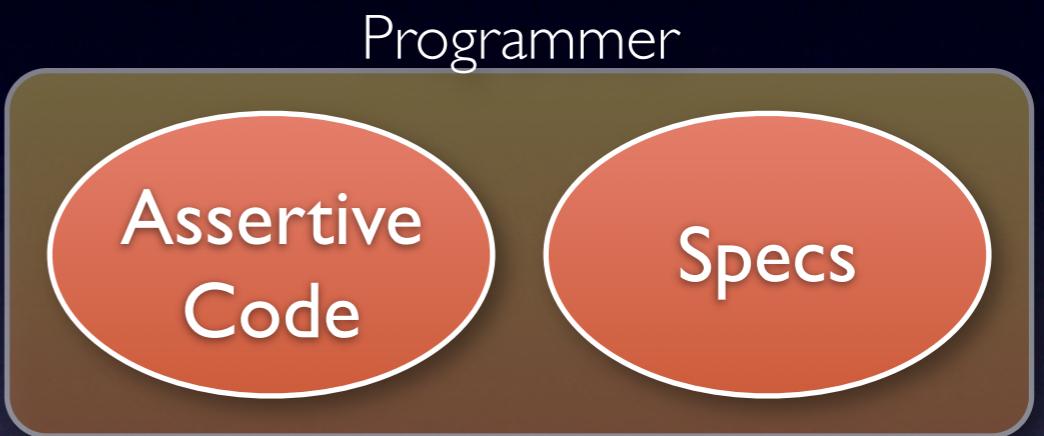
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# The Grand Challenge

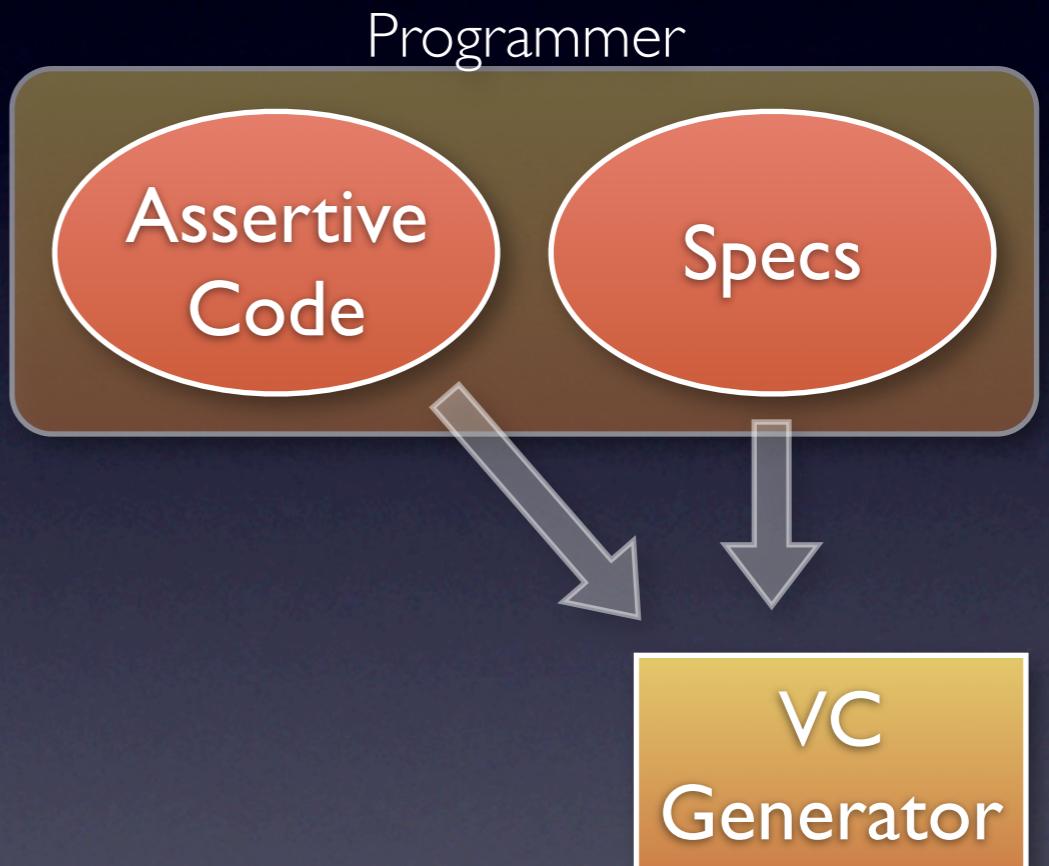
“ I revive an old challenge: the construction and application of a **verifying compiler** that guarantees correctness of a program before running it. ”

Tony Hoare, The Verifying Compiler: A Grand Challenge for Computing Research, 2003

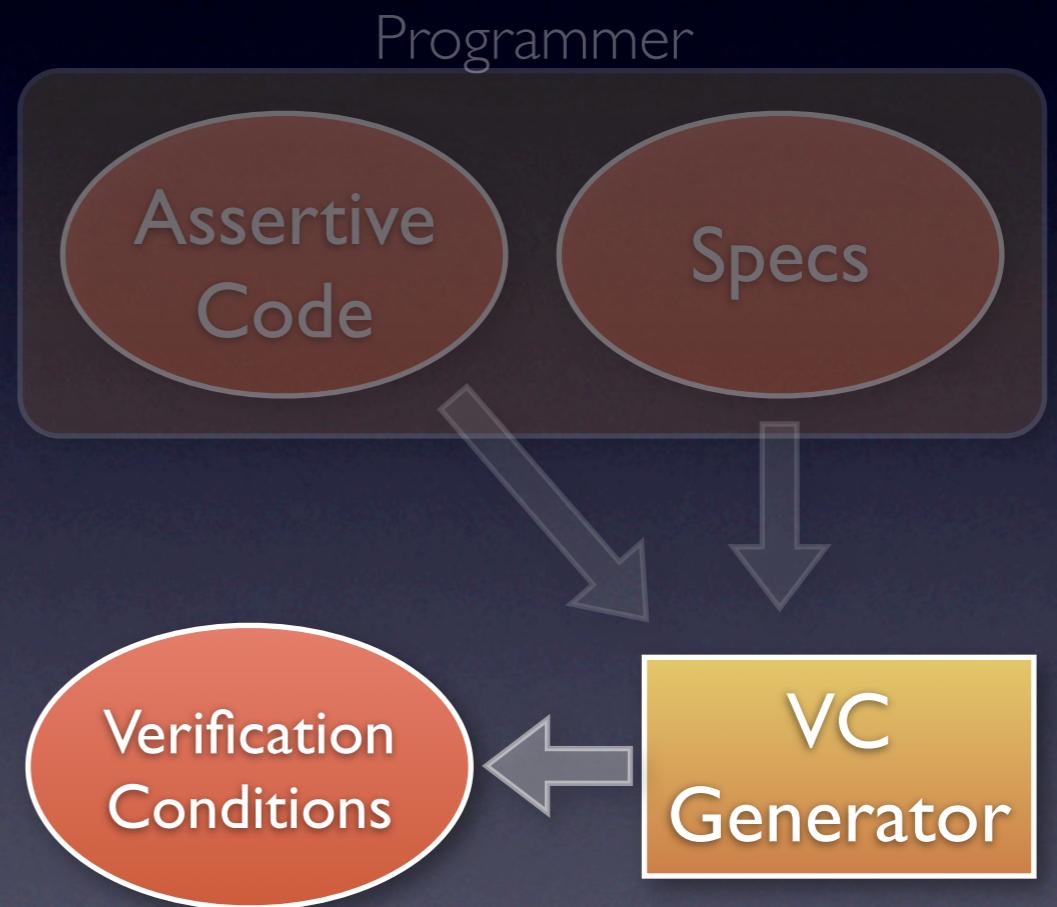
# Organization of verification system



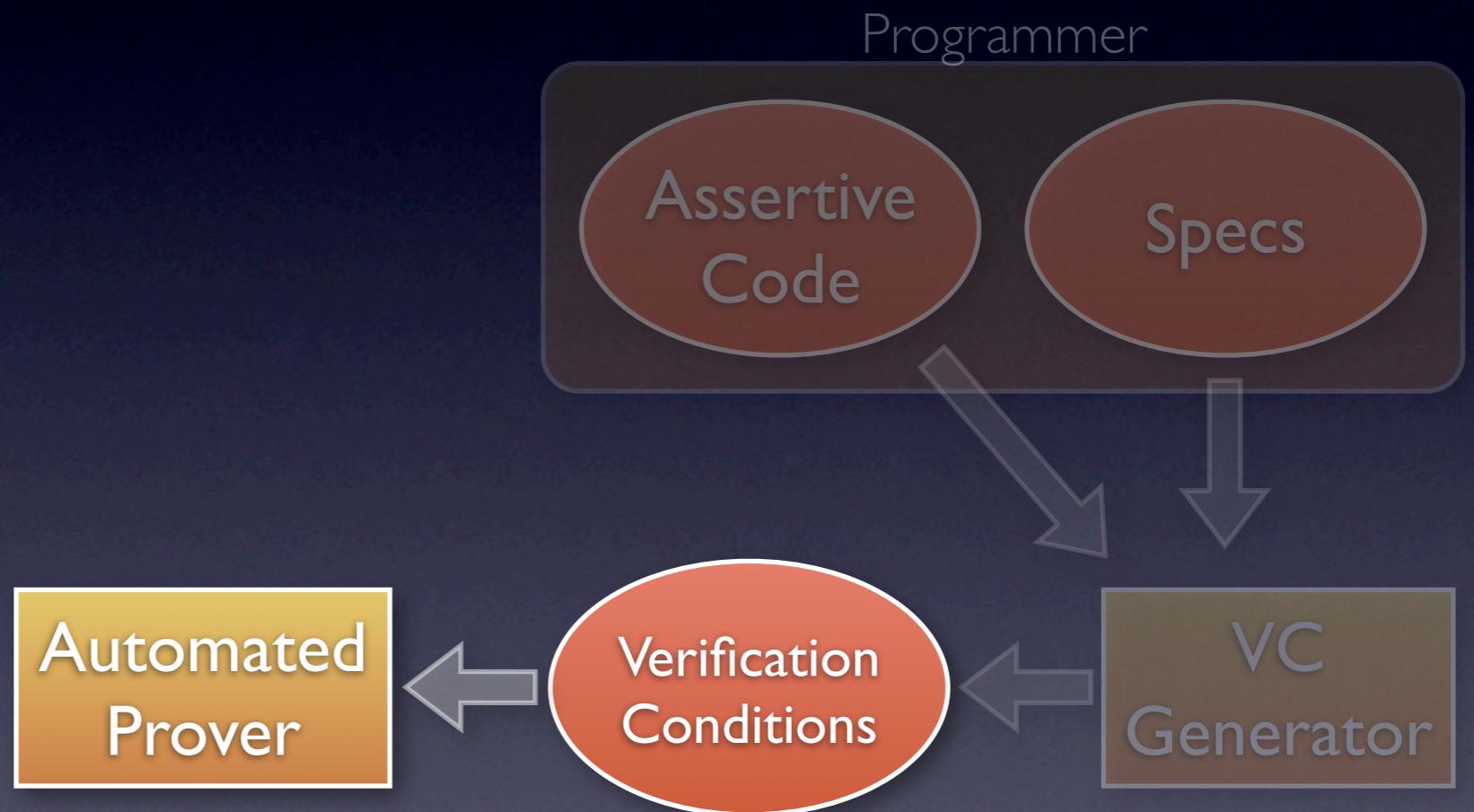
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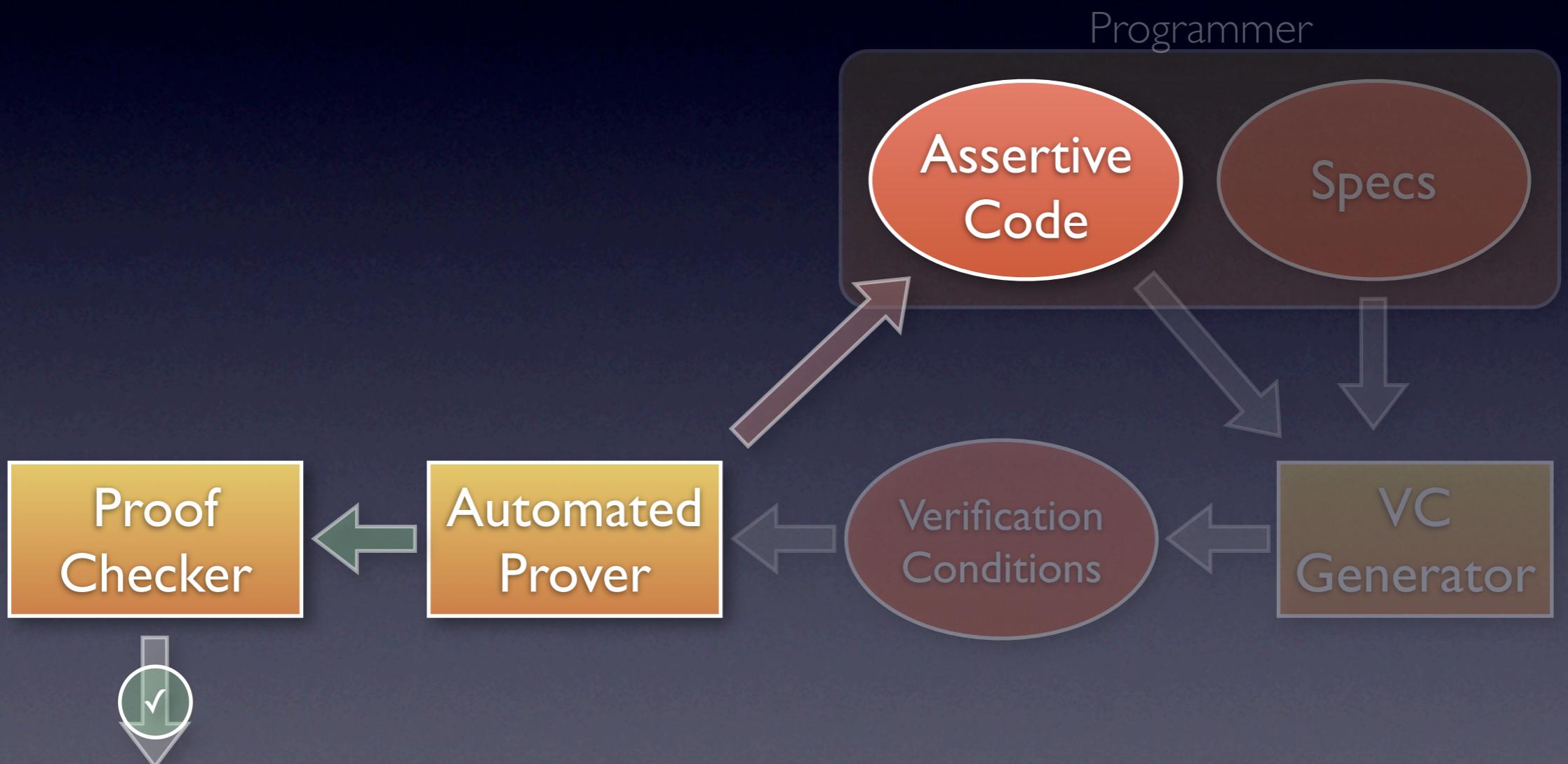
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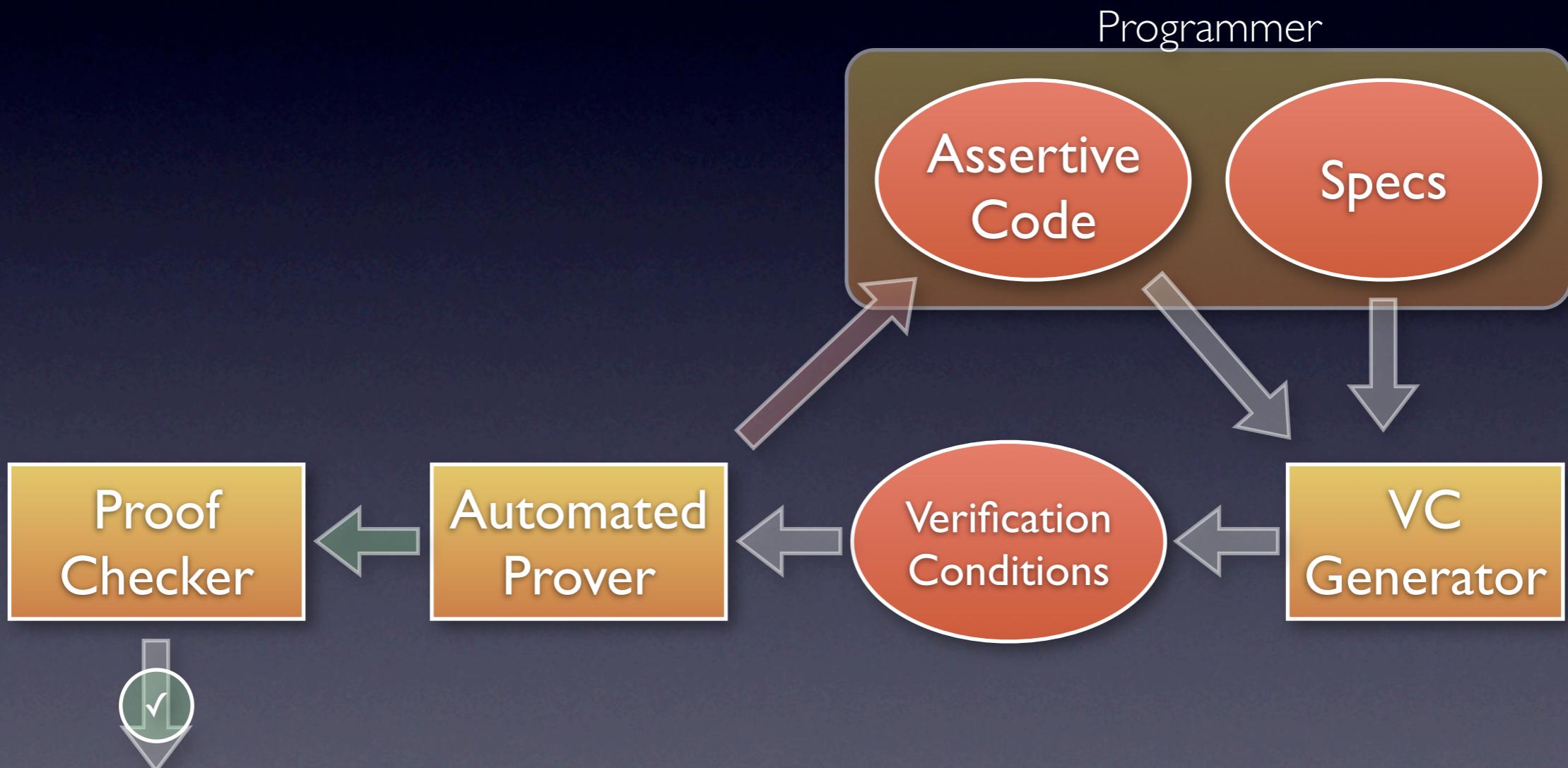
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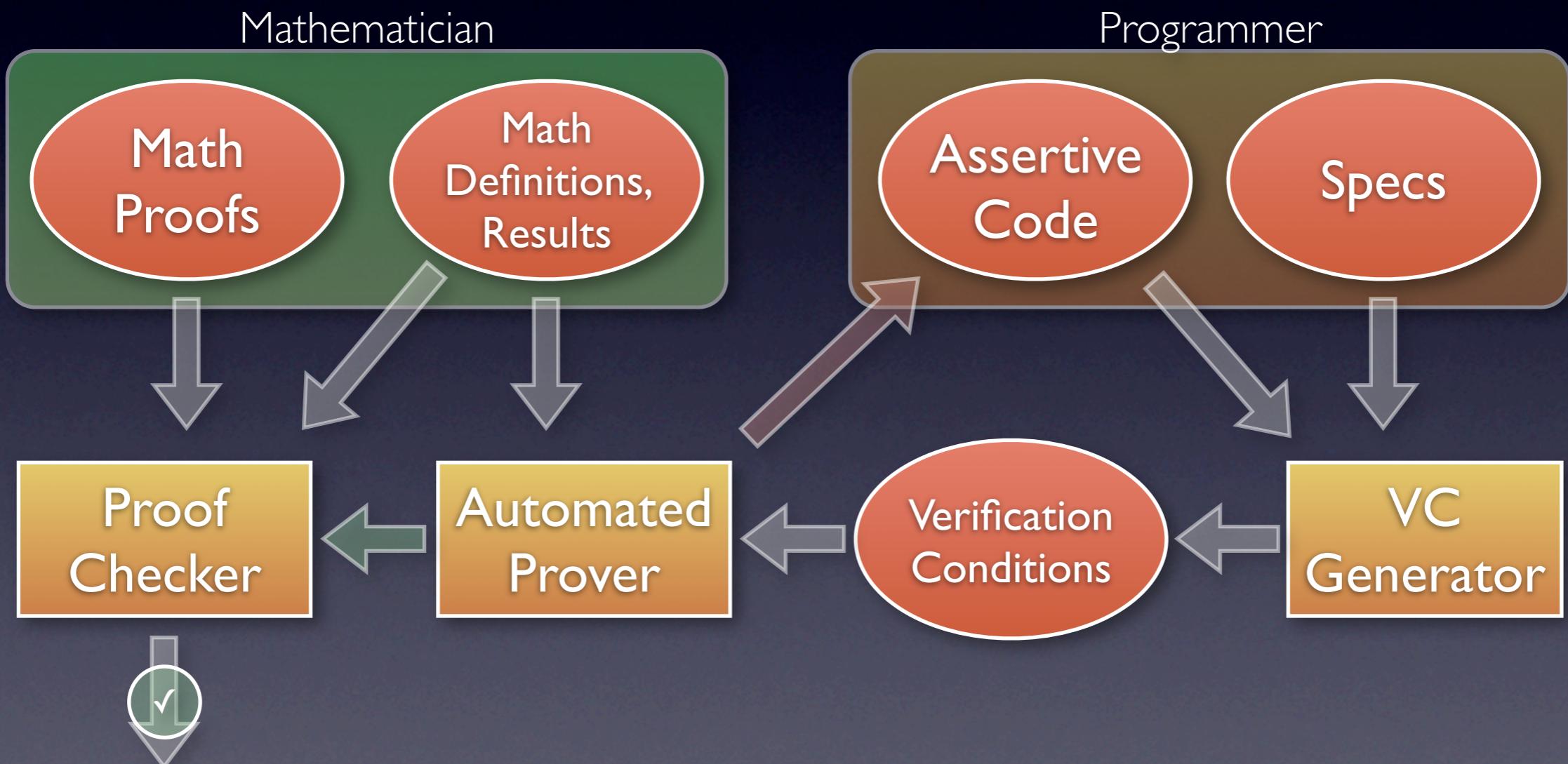
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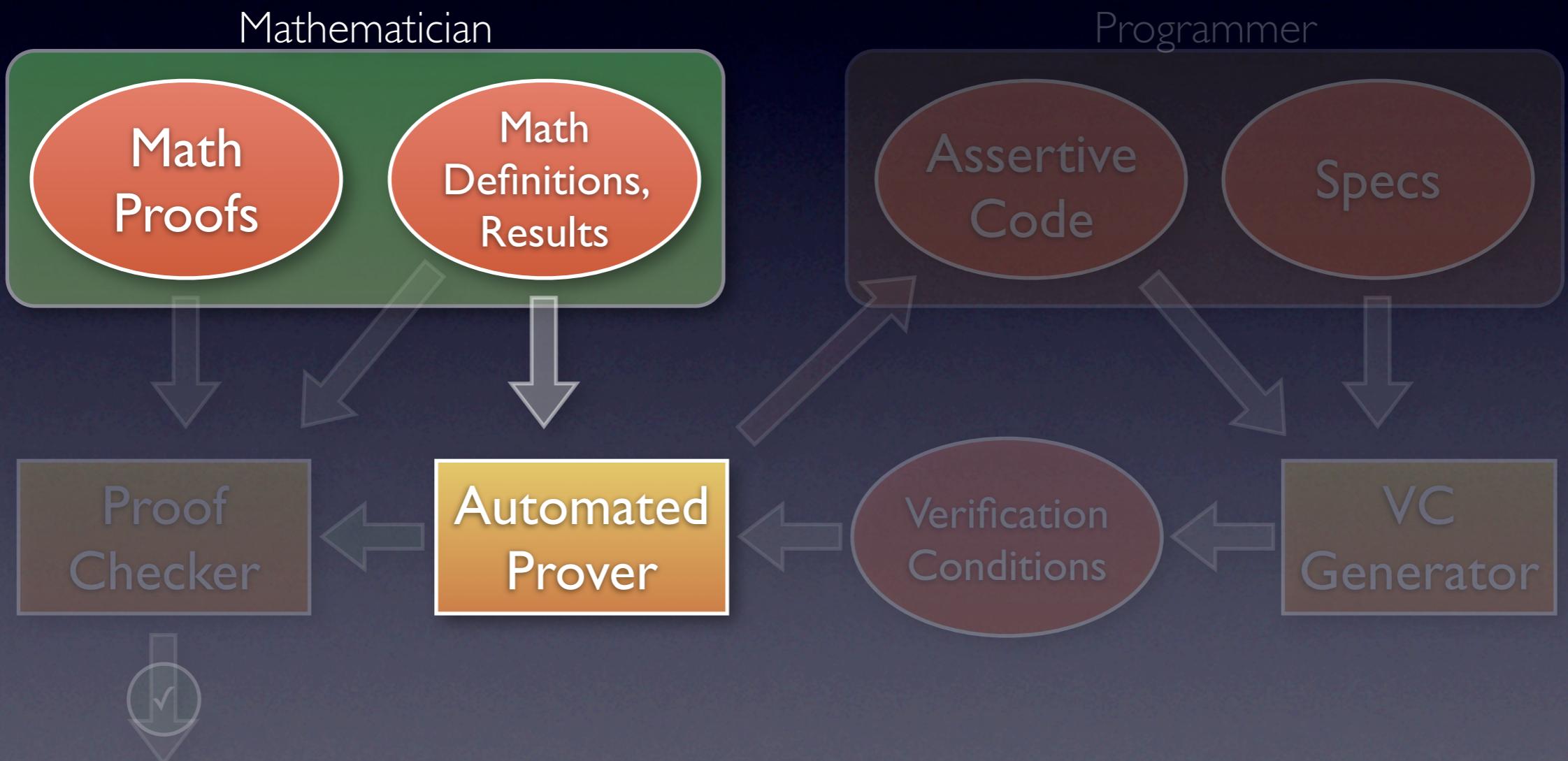
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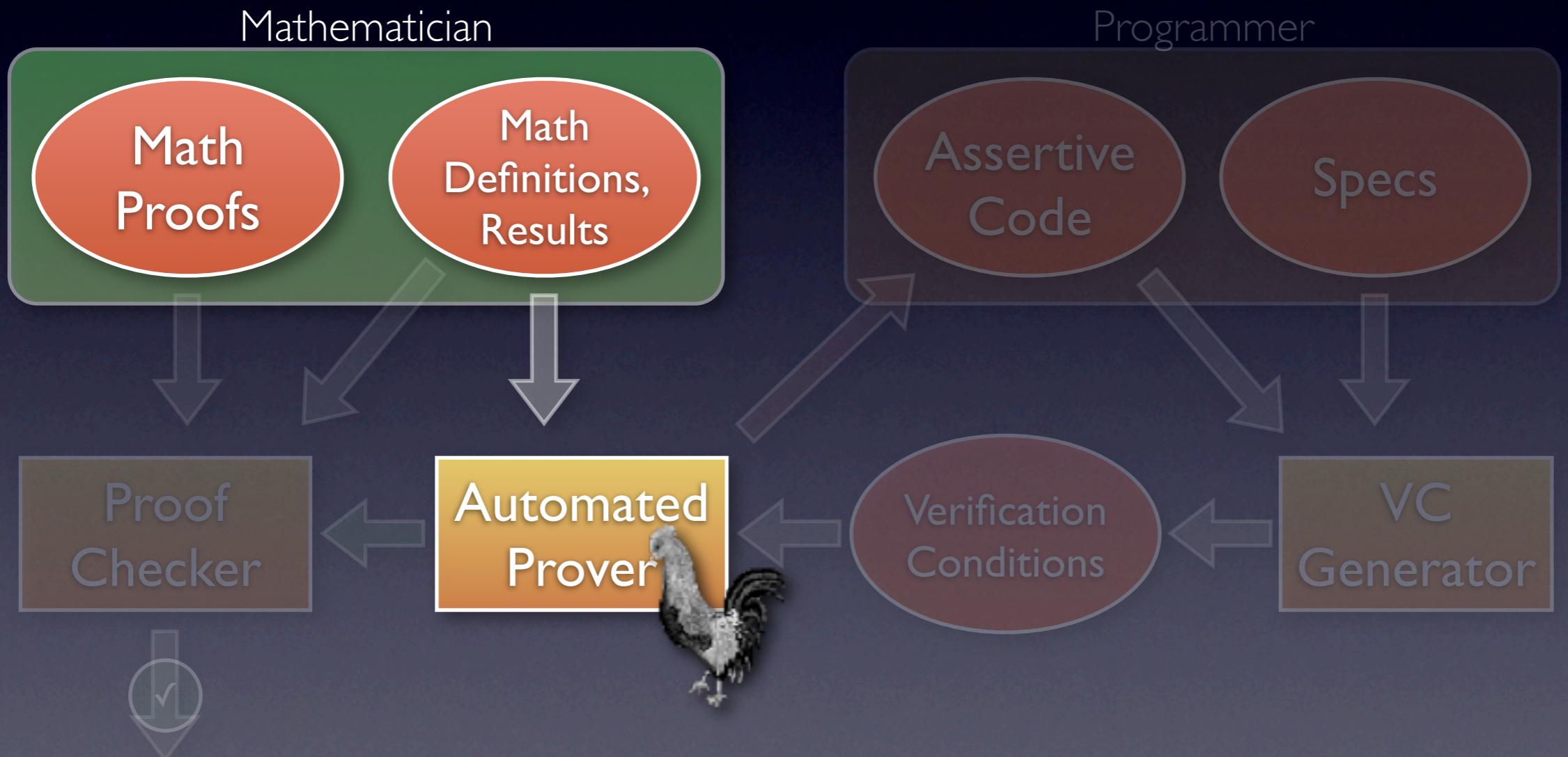
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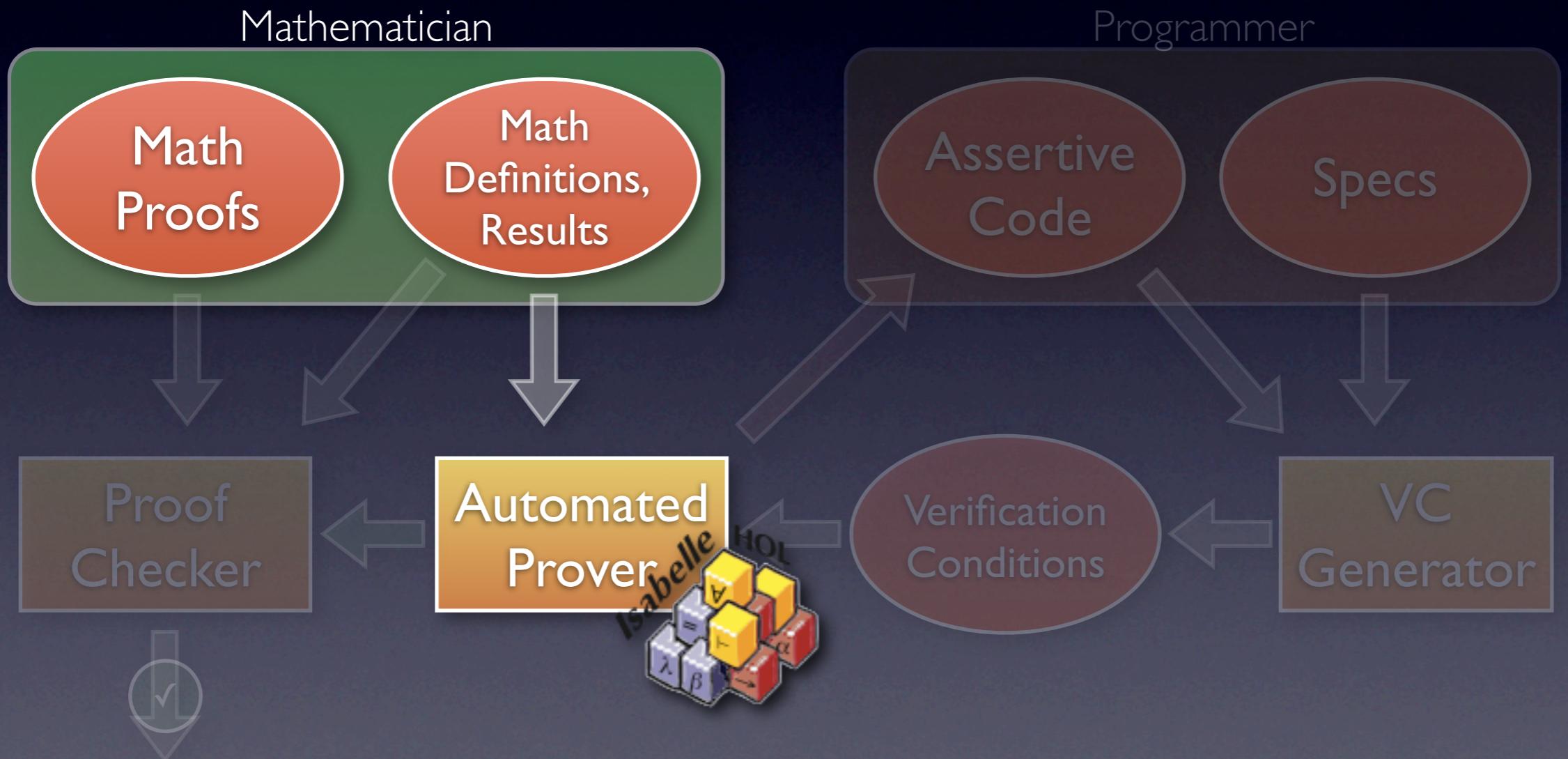
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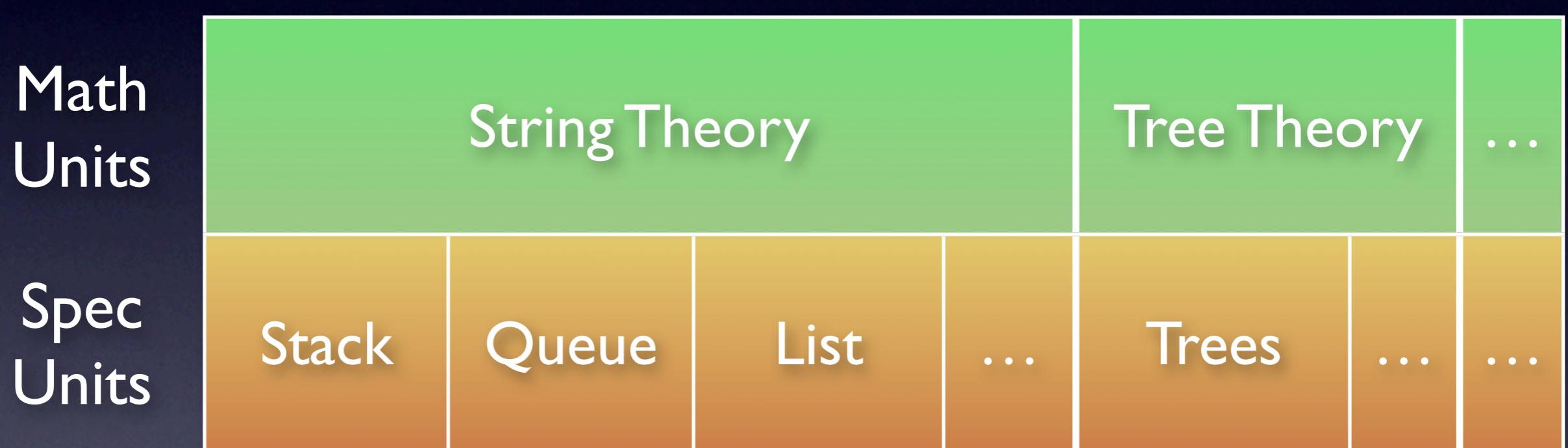
# Organization of verification system



# Organization of verification system



# Mathematical Theory Reuse



# Research Question

- Can we use an off-the-shelf theorem prover to prove VCs automatically and still maintain the properties identified earlier?
- We explore this question via an extension to a queue ADT.

# Queue Specification

```
contract QueueTemplate (type Item)

math subtype QUEUE_MODEL is string of Item

type Queue is modeled by QUEUE_MODEL
  exemplar q
  initialization ensures
    q = empty_string

procedure Enqueue (updates q: Queue,
                  clears x: Item)
  ensures
    q = #q * <#x>

procedure Dequeue (updates q: Queue,
                   replaces x: Item)
  requires
    q /= empty_string
  ensures
    #q = <x> * q

function IsEmpty (restores q: Queue): control
  ensures
    IsEmpty = (q = empty_string)

end QueueTemplate
```



Specs

# Queue Specification

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contract QueueTemplate (type Item)

  math subtype QUEUE_MODEL is string of Item

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                     replaces x: Item)
    requires
      q /= empty_string
    ensures
      #q = <x> * q

  function IsEmpty (restores q: Queue): control
    ensures
      IsEmpty = (q = empty_string)

end QueueTemplate
```



Specs

# Queue Specification

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procedure Dequeue (updates q: Queue,
                   replaces x: Item)
  requires
    q /= empty_string
  ensures
    #q = <x> * q

function IsEmpty (restores q: Queue): control
  ensures
    IsEmpty = (q = empty_string)

end QueueTemplate
```

Specs

# Queue Flip Specification

```
contract QueueFlip enhances QueueTemplate
procedure Flip (updates q: Queue)
ensures
    q = reverse (#q)
end QueueFlip
```

Specs

# Iterative Implementation

```
realization Iterative implements QueueFlip

facility StackFacility is StackTemplate (Item)

procedure Flip (updates q: Queue)
    variable s: Stack
    loop
        maintains reverse(#s) * #q = reverse(s) * q
        decreases |q|
        while not IsEmpty (q) do
            variable x: Item
            Dequeue (q, x)
            Push (s, x)
        end loop
    loop
        maintains #q * #s = q * s
        decreases |s|
        while not IsEmpty (s) do
            variable x : Item
            Pop (s, x)
            Enqueue (q, x)
        end loop
    end Flip
end Iterative
```

Assertive  
Code

# Iterative Implementation

```
realization Iterative implements QueueFlip

facility StackFacility is StackTemplate (Item)

procedure Flip (updates q: Queue)
    variable s: Stack
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        maintains reverse(#s) * #q = reverse(s) * q
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            variable x: Item
            Dequeue (q, x)
            Push (s, x)
        end loop
    loop
        maintains #q * #s = q * s
        decreases |s|
        while not IsEmpty (s) do
            variable x : Item
            Pop (s, x)
            Enqueue (q, x)
        end loop
    end Flip
end Iterative
```

Assertive  
Code

# Iterative Implementation

```
realization Iterative implements QueueFlip

facility StackFacility is StackTemplate (Item)

procedure Flip (updates q: Queue)
    variable s: Stack
    loop
        maintains reverse(#s) * #q = reverse(s) * q
        decreases |q|
        while not IsEmpty (q) do
            variable x: Item
            Dequeue (q, x)
            Push (s, x)
        end loop
    loop
        maintains #q * #s = q * s
        decreases |s|
        while not IsEmpty (s) do
            variable x : Item
            Pop (s, x)
            Enqueue (q, x)
        end loop
    end Flip

end Iterative
```

Assertive  
Code

# Iterative Implementation

```
realization Iterative implements QueueFlip

facility StackFacility is StackTemplate (Item)

procedure Flip (updates q: Queue)
    variable s: Stack
    loop
        maintains reverse(#s) * #q = reverse(s) * q
        decreases |q|
        while not IsEmpty (q) do
            variable x: Item
            ... (loop body)
    end loop
end Flip

end Iterative
```

Assertive  
Code

# Recursive Implementation

```
realization Recursive implements QueueFlip

procedure Flip (updates q: Queue)
  decreases |q|
  if not IsEmpty (q) then
    variable x: Item
    Dequeue (q, x)
    Flip (q)
    Enqueue (q, x)
  end if
end Flip

end Recursive
```

Assertive  
Code

# AVC for Recursive Implementation

$\text{is\_initial } (x_2)$

$\text{is\_initial } (x_5)$

$\langle x_3 \rangle \circ q_3 \neq \lambda$

---

$\text{reverse } (q_3) \circ \langle x_3 \rangle = \text{reverse } (\langle x_3 \rangle \circ q_3)$

Verification  
Conditions

# Theory Libraries

- The required theorem was already defined in string theory (requirement of well-developed theories).
- Should be able to use any proof assistant with the theory library.

Math  
Proofs

Math  
Definitions,  
Results

# Lessons Learned

- It is possible to use an interactive proof assistant as an automated VC prover.
- Interactive proof assistants can be used off-the-shelf by importing the mathematical theories of interest.

# Questions?