Free and Bound Variable Identifiers, Desugaring, in Oz

A Grammar for Oz with Sugars

Consider the grammar in Figure 1, which describe data structures (records called abstract syntax trees). These are to be thought of as representing in Oz data the statements and expressions of Oz itself. Thus the grammar in the figure represents the abstract syntax of a sugared version of Oz. For example the following Oz statement:

```
A = B
is represented by the record structure:
assignStmt("A" varIdExp("B"))
```

Note that the grammar uses Oz strings to represent variable identifiers. A more complex example that illustrates more features of the representation is given in Figure 2 on the next page.

```
⟨Statement⟩ ::= skipStmt
                       seqStmt(\langle List \langle Statement \rangle \rangle)
                       localStmt((String) (Statement))
                       assignStmt(\langle String \rangle \langle Expression \rangle)
                       ifStmt (\langle Expression \rangle \langle Statement \rangle \langle Statement \rangle \rangle \langle Statement \rangle 
                      caseStmt ((Expression) (Pattern) (Statement))
                      applyStmt((Expression) (List (Expression)))
                       \texttt{namedFunStmt} \; (\langle String \rangle \; \; \langle List \; \langle Pattern \rangle \rangle \; \; \langle Expression \rangle)
                      inStmt ((Pattern) \langle (Expression) \langle (Statement))
\langle Expression \rangle ::= varIdExp(\langle String \rangle)
                       atomExp(\langle Atom \rangle) | boolExp(\langle Bool \rangle)
                      numExp(\langle Number \rangle)
                     recordExp(\langle Expression \rangle \langle List \langle Field \rangle \rangle)
                     procExp((List (Pattern)) (Statement))
                     ifExp((Expression) \langle (Expression))
                      caseExp((Expression) (Pattern) (Expression))
                     applyExp((Expression) (List (Expression)))
⟨Pattern⟩ ::= varIdPat(⟨String⟩)
                | atomPat (\langle Atom \rangle) | boolPat (\langle Bool \rangle)
               | recordPat((Atom) (List(Field)))
\langle \text{Field} \rangle ::= \text{colonFld}(\langle \text{Atom} \rangle \langle \text{Expression} \rangle)
               \mid posFld(\langle Exp \rangle)
```

Figure 1: Grammar for a simplified subset of the practical version of the declarative language of chapter 3, based on the textbook's section 2.6 [VH04]. The type $\langle String \rangle$ is the type of character strings in Oz, $\langle Atom \rangle$ is the type of atoms in Oz, etc.

```
The following Oz statement:
fun {AddToEach A#B Ls}
   case Ls of
     (X#Y) | T then (\{Plus A X\} # \{Plus B Y\}) | \{AddToEach A#B T\}
   else nil
   end
end
is represented by the following record structure:
namedFunStmt ("AddToEach"
              [recordPat('#'
                          [posFld(varIdExp("A")) posFld(varIdExp("B"))])
               varIdPat("Ls")]
   caseExp(varIdExp("Ls")
            recordPat('|'
                       [posFld(recordExp(atomExp('#')
                                         [posFld(varIdExp("X"))
                                          posFld(varIdExp("Y"))]))
                       posFld(varIdExp("T"))])
            recordExp(atomExp('|')
                       [posFld(recordExp(atomExp('#')
                                          [posFld(applyExp(varIdExp("Plus")
                                                             [varIdExp("A")
                                                             varIdExp("X")]))
                                           posFld(applyExp(varIdExp("Plus")
                                                             [varIdExp("B")
                                                              varIdExp("Y")]))))
                       posFld(applyExp(varIdExp("AddToEach")
                                         [recordExp(atomExp('#')
                                                     [posFld(varIdExp("A"))
                                                      posFld(varIdExp("B"))])
                                          varIdExp("T")]))])
            atomExp(nil)))
```

Figure 2: Example showing how Oz is parsed into the record structures (abstract syntax trees) from Figure 1 on the previous page.

The Functions (or Problems to be Solved)

1. The function

```
FreeVarIds: <fun {$ <Statement>}: <Set <String>>
```

takes a $\langle Statement \rangle$ in the grammar of Figure 1 on page 1 and returns a set of all the variable identifiers that occur free in its argument. Its implementation is in FreeVarIds.oz. There are tests in the file FreeVarIdsTest.oz.

See the file TestOutput.txt for results of testing.

2. The function

```
BoundVarIds: <fun {$ <Statement>}: <Set <String>>
```

takes a $\langle Statement \rangle$ in the grammar of Figure 1 on page 1 and returns a set of all the variable identifiers that occur bound in its argument. Its implementation is in BoundVarIds.oz. There are tests in the file BoundVarIdsTest.oz.

See the file TestOutput.txt for results of testing.

3. The function

```
Desugar: <fun {$ <Statement>}: <Statement>
```

takes a (Statement) in the grammar of Figure 1 on page 1 and returns a (Statement) in the subset of that grammar that represents Oz statements in the kernel language of the textbook [VH04, Section 2.3].

An implementation is in Desugar.oz. There are tests in the file DesugarTest.oz.

See the file TestOutput.txt for results of testing.

References

[VH04] Peter Van Roy and Seif Haridi. *Concepts, Techniques, and Models of Computer Programming*. The MIT Press, Cambridge, Mass., 2004.