## COT 6410 Assignment 6

Solution

For the 3SAT instance: $\left(x_{1} \vee x_{2} \vee x_{3}\right) \wedge\left(\neg x_{1} \vee \neg x_{2} \vee x_{3}\right) \wedge\left(x_{1} \vee x_{4} \vee x_{4}\right) \wedge\left(x_{2} \vee x_{3} \vee x_{3}\right)$ :
(1) The equivalent SubsetSum instance:

|  |  | $x_{1}$ | $x_{2}$ | $x_{3}$ | $x_{4}$ | $x_{1} \vee x_{2} \vee x_{3}$ | $-x_{1} \vee-x_{2} \vee x_{3}$ | $x_{1} \vee x_{4} \vee x_{4}$ | $x_{2} \vee x_{3} \vee x_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $\mathrm{x}_{1}$ | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 2 | $-x_{1}$ | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 3 | $\mathrm{x}_{2}$ | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| 4 | $-x_{2}$ | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 5 | $\mathrm{x}_{3}$ | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 2 |
| 6 | $-x_{3}$ | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 7 | $\mathrm{x}_{4}$ | 0 | 0 | 0 | 1 | 0 | 0 | 2 | 0 |
| 8 | $-\mathrm{x}_{4}$ | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 9 | $\mathrm{C}_{1}$ | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 10 | $\mathrm{C}_{1}{ }^{\prime}$ | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 11 | $\mathrm{C}_{2}$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 12 | $\mathrm{C}_{2}{ }^{\prime}$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 13 | $\mathrm{C}_{3}$ | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| 14 | $\mathrm{C}_{3}{ }^{\prime}$ | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| 15 | $\mathrm{C}_{4}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 16 | $\mathrm{C}_{4}{ }^{\prime}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
|  | $\mathrm{Goal}^{2}$ | 1 | 1 | 1 | 1 | 3 | 3 | 3 | 3 |

To achieve the desired sum, below rows can be selected: 1 ( $x_{1}=$ True), 3 ( $x_{2}=$ True), 5 ( $x_{3}=$ True), 7 ( $\mathrm{x}_{4}=$ True), 11 ( $\mathrm{C}_{2}=$ True), and $12\left(\mathrm{C}_{2}{ }^{\prime}=\right.$ True $)$. In this case, all clauses are satisfied.
(2) The equivalent Vertex Cover instance:

Variable gadgets:


Clause gadgets:


Combined gadgets:


The number of vertices needed to be selected is $k=n+2 m=4$ (the number of variables) $+2 \times 4$ (the number of clauses) $=12$. Since the graph above has a vertex cover with exact 12 vertices (the circled ones), all clauses are satisfied.
(3) The equivalent Independent Set instance:

## Clause gadgets:



Combined gadgets:


The number of vertices needed to be selected in the independent set is $\mathrm{k}=\mathrm{m}=4$ (the number of clauses). Since the graph above has an independent set with exact 4 vertices (the circled ones), all clauses are satisfied.
(4) The equivalent Hamiltonian Circuit instance:

Assume for each path $i$ has $3 m+3$ vertices (i.e. vertex 1 , vertex $2, \ldots$, vertex $3 m+3$ ), where $m$ is the number of clauses. If variable $x_{i}$ is True, the direction of passing the path $i$ is left to right. If variable $-x_{i}$ is True, the direction of passing the path $i$ is right to left. For clause $C_{j}$, if $x_{i}$ is in $C_{j}, C_{j}$, has edge from vertex $3 j$ to vertex $3 j+1$; if $\neg x_{i}$ is in $C_{j}, C_{j}$, has edge from vertex $3 j+1$ to vertex $3 j$.

Variable gadgets:
$\mathrm{x}_{1}$


Combined gadgets:


Below is the graph with a Hamiltonian Circuit highlighted, indicating all clauses are satisfied:


