

Assignment#4 Key

1. NotDominating(ND) = { f | for some x, f(x) < x }.

a.) Show some minimal quantification of some known primitive recursive predicate that provides an upper bound for the complexity of **ND**.

$\exists \langle x, t \rangle [STP(f, x, t) \ \& \ (VALUE(f, x, t) < x)]$

b.) Use Rice's Theorem to prove that **ND** is undecidable. Be Complete.

ND is non-trivial as $C0(x) = 0 \in \mathbf{ND}$ and $S(x) = x+1 \notin \mathbf{ND}$

Let **f, g** be two arbitrary indices of procedures such that $\forall x \ f(x) = g(x)$

$f \in \mathbf{ND}$ iff $\exists x \ f(x) < x$ iff $f(x_0) < x_0$ for some x_0 iff $g(x_0) < x_0$ as $\forall x \ f(x) = g(x)$ implies $\exists x \ g(x) < x$ iff $g \in \mathbf{ND}$

$f \notin \mathbf{ND}$ iff $\forall x \ [f(x) \downarrow \text{ implies } f(x) > x]$ iff $\forall x \ [g(x) \downarrow \text{ implies } g(x) > x]$ as $\forall x \ f(x) = g(x)$ iff $g \notin \mathbf{ND}$

c.) Show that $K = \{ f \mid f(f) \text{ converges} \}$ is many-one reducible to **ND**.

Let **f** be an arbitrary index. From **f**, define $\forall x \ F_f(x) = f(f) - f(x)$.

$f \in K$ implies $\forall x \ F_f(x) = 0$ implies $F_f \in \mathbf{ND}$.

$f \notin K$ implies $\forall x \ F_f(x)$ diverges implies $F_f \notin \mathbf{ND}$.

Thus, $K \leq_m \mathbf{ND}$

d.) Show that **ND** is many-one reducible to $K = \{ f \mid f(f) \text{ converges} \}$

Let **f** be an arbitrary index. From **f**, define $\forall y \ F_f(y) = \exists \langle x, t \rangle \ STP(f, x, t) \ \& \ (VALUE(f, x, t) < x)$

$f \in \mathbf{ND}$ implies $\forall y \ F_f(y)$ converges implies $F_f(F_f)$ converges implies $F_f \in K$

$f \notin \mathbf{ND}$ implies $\forall y \ F_f(y)$ diverges implies $F_f \notin K$.

Thus, $\mathbf{ND} \leq_m K$

2. AlwaysDominates(AD) = { f | for all x, f(x) > x }

a.) Show some minimal quantification of some known primitive recursive predicate that provides an upper bound for the complexity of AD.

$\forall x \exists t [\text{STP}(f,x,t) \ \& \ (\text{VALUE}(f,x,t) > x)]$

b.) Use Rice's Theorem to prove that AD is undecidable. Be Complete.

AD is non-trivial as $S(x) = x+1 \in \text{AD}$ and $C0(x) = 0 \notin \text{AD}$

Let f,g be two arbitrary indices of procedures such that $\forall x f(x) = g(x)$

$f \in \text{AD}$ iff $\forall x f(x) < x$ iff $\forall x g(x) < x$ iff $g \in \text{AD}$

c.) Show that $\text{TOT} = \{ f \mid \text{for all } x, f(x) \text{ converges} \}$ is many-one reducible to AD.

Let f be an arbitrary index. From f, define $\forall x F_f(x) = f(x) - f(x) + x + 1$.

$f \in \text{TOT}$ implies $\forall x F_f(x) = x+1$ implies $F_f \in \text{AD}$.

$f \notin \text{TOT}$ implies $\exists x F_f(x)$ diverges implies $F_f \notin \text{AD}$.

Thus, $\text{TOT} \leq_m \text{AD}$

d.) Show that AD is many-one reducible to $\text{TOT} = \{ f \mid \text{for all } x, f(x) \text{ converges} \}$

Let f be an arbitrary index. From f, define $\forall x F_f(x) = \mu y [f(x) > y]$

$f \in \text{AD}$ implies $\forall x F_f(x)$ converges implies $F_f \in \text{TOT}$

$f \notin \text{AD}$ implies $\exists x F_f(x)$ diverges implies $F_f \notin \text{TOT}$.

Thus, $\text{AD} \leq_m \text{TOT}$