Assignment#2 Key

1a. EveryOther(L) = { $a_1\ a_3\ \dots\ a_{2n-1}$ | $a_1\ a_2\ a_3\ \dots\ a_{2n-1}a_{2n}$ is in L }

- Approach 1: Let **L** be a Regular language over the finite alphabet **Σ**. For each **a∈Σ**, define **f(a) = {a,a'}**, **g(a) = a'** and **h(a) = a**, **h(a') = λ**,
  f is a substitution, **g** and **h** are homomorphisms.
  EveryOther(L) = h(f(L) ∩ **(Σ ▪ g(Σ))*)**
- Why this works:
  f(L) gets us every possible random priming of letters of strings in L.
  (Σ ▪ g(Σ))* gets every word composed of pairs of unprimed and primed letters from **Σ**. Intersecting this with **f(L)** gets strings of the form
  $\mathbf{a_1\ a_2'\ a_3\ \ a_4' \dots\ a_{2n-1}a_{2n}'}$ where $\mathbf{a_1\ a_2\ a_3\ a_4 \dots\ a_{2n-1}a_{2n}}$ is in **L**
  Applying the homomorphism **h** erases all primed letters resulting in every string $\mathbf{a_1\ a_3\ \dots\ a_{2n-1}}$ where $\mathbf{a_1\ a_2\ a_3\ a_4\ \dots\ a_{2n-1}a_{2n}}$ is in **L**, precisely the language **EveryOther(L)** that we sought. This works as Regular Languages are closed under intersection, concatenation, *, substitution and homomorphism.

1a. EveryOther(L) = { $a_1\ a_3 \dots a_{2n-1}$ | $a_1\ a_2\ a_3 \dots a_{2n-1}a_{2n}$ is in L }

- Approach 2: Let **L** be a Regular language over the finite alphabet **Σ**. Assume **L** is recognized by the DFA $\mathbf{A_1 = (Q, \Sigma, \delta_1, q_1, F)}$. Define NFA $\mathbf{A_2 = (Q, \Sigma, \delta_2, q_1, F)}$, where $\mathbf{\delta_2(q,a) = union(b{\in}\Sigma)\ \{\ \delta_1(\delta_1(q,a),b)\ \}}$
- Why this works:
  Every transition that $\mathbf{A_2}$ takes is one that $\mathbf{A_1}$ would have taken when reading a pair that starts with the character read by $\mathbf{A_1}$ followed by any arbitrary character.

1b. Half(L) = { x | there exists a y, |x| = |y| and xy is in L }

- Let **L** be a Regular language over the finite alphabet $\mathbf{\Sigma}$. Assume **L** is recognized by the DFA $\mathbf{A_1 = (Q, \Sigma, \delta_1, q_1, F)}$. Define the NFA
  $\mathbf{A_2 = ((Q\times Q \times Q)\cup\{q_0\}, \Sigma, \delta_2, q_0, F')}$, where
  $\mathbf{\delta_2(q_0,\lambda) = union(q\in Q)\ \{<q_1, q, q>\}}$ **and**
  $\mathbf{\delta_2(< q, r, s >, a) = union(b\in\Sigma)\ \{ < \delta_1(q,a), \delta_1(r,b), s > \},\ \ q,r,s \in Q}$
  $\mathbf{F' = union(q\in Q)\ \{<q, f, q>\}, f\in F}$
- Why this works:
  The first part of a state **< q, r, s >** tracks $\mathbf{A_1}$**.**
  The second part of a state **< q, r, s >** tracks $\mathbf{A_1}$ for precisely all possible strings that are the same length as what $\mathbf{A_1}$ is reading in parallel. This component starts with a guess as to what state $\mathbf{A_1}$ might end up in.
  The third part of a state **< q, r, s >** remembers the initial guess.
  Thus, $\mathbf{\delta_2{}^*(< q_1, q, q >, x) = \{\delta_1{}^*(q_0,x), \delta_1{}^*(q,y), q >\}}$ for arbitrary **y**, **|x|=|y|**
  We accept if the initial guess was right and the second component is final, meaning **xy** is in **L**..

2. $L = \{ a^m b^n c^t \mid t = \min(m,n) \}$

a.) Use the **Myhill-Nerode Theorem** to show **L <u>is not</u>** Regular.

Define the equivalence classes **$[a^i b^i]$, $i \geq 0$**

Clearly **$a^i b^i c^i$** is in **L**, but **$a^j b^j c^i$** is not in **L** when **$j \neq i$**

Thus, **$[a^i b^i] \neq [a^j b^j]$** when **$j \neq i$** and so the index of **R_L** is infinite.

By Myhill-Nerode, **L** is not Regular.

2. L = { a^m b^n c^t | t = min(m,n) }

b.) Use the **Pumping Lemma for CFLs** to show **L** <u>**is not**</u> a CFL

Me: **L** is a CFL

PL: Provides **N>0**

Me: **z = a^N b^N c^N**

PL: **z = uvwxy, |vwx| ≤ N**, **|vx| > 0**, and **∀i≥0 uv^iwx^iy ∈ L**

Me: Since **|vwx| ≤ N**, it can consist of **a**'s and/or **b**'s or **b**'s and/or **c**'s but never all three.
Assume it contains no **c**'s then **i=0** decreases the number of **a**'s and/or the number of **b**'s, but not the **c**'s and so there are more **c**'s than the minimum of **a**'s and **b**'s.
Assume it contains **c**'s then **i=2** increases the number of **c**'s and maybe number of **b**'s, but not the **a**'s and so there are more than **N c**'s but just N a's.

2. $L = \{ a^m b^n c^t \mid t = \min(m,n) \}$

c.) Present a CSG for **L** to show it **is** context sensitive

G = ({ A, B, C, <a>, <b> }, { a, b, c }, R, A)

A → aBbc | abc | a | b | λ

B → aBbC | a<a>bC | ab<b>C // allow more a's or more b's

Cb → bC // Shuttle C over to a c

Cc → cc // Change C to a c

<a> → a<a> | λ

<b> → b<b> | λ