

COT 4210 Quiz #3 Part B: Undecidability, Reducibility 3/25/2021

Regular Start Time: 2:10 pm (EST)

Regular End Time: 2:45 pm (EST)

Regular Late Time: 2:55 pm (EST)

4) (10 pts) Prove that the following language is undecidable:

$L_4 = \{ \langle M \rangle, k \mid \langle M \rangle \text{ is the encoding of a Turing Machine which accepts all strings of length } k \}$

Solution

We will use proof by contradiction to prove the claim. Assume to the contrary that L_4 is decidable and let a Turing Machine X decide membership in L_4 . We describe how to create a Turing Machine Y which decides membership in A_{TM} by using X :

$Y(\text{Turing Machine } M, \text{String } w) \{$

1. Create a Turing Machine M' which operates as follows: if its input is length $|w|$, it erases its input, replaces it with w , and then follows all of the steps of M . If its input is not length $|w|$, M' rejects.
2. We call TM X with input $(M', |w|)$ and accept if X accepts $(M', |w|)$ and reject if X rejects $(M', |w|)$

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Since we've created a decider for A_{TM} , which is known to be undecidable, we have arrived at a contradiction. It follows that our initial assumption, that L_4 was decidable is incorrect, which means that L_4 is undecidable, as desired.

Grading Criteria: +2 for assuming decider for L_4
+2 for attempting to write a decider for a known decidable language
+5 for the actual proof
+1 for the conclusion

5) (12 pts) Consider the two following problems:

$\text{SAFE-CLASSES} = \{ (S, E) \mid S \text{ is a list of students, and } E \text{ is a list of pair of students who are not allowed to be in the same class, and there exists a way to split all students in } S \text{ into two non-empty classes such that the requirements given by } E \text{ are satisfied} \}$

$\text{TWO-COLORABLE} = \{ G \mid G \text{ is an undirected graph such that each vertex in } G \text{ can be assigned one of two colors, Red or Blue, such that no edge in } G \text{ connects two vertices that have the same color.} \}$

SAFE-CLASSES is mapping reducible to TWO-COLORABLE . To do the reduction, given an input of students and pairs of students who can't be in the same class, you can create a graph G

that is two-colorable if and only if there's a way to split the students in S into two classes satisfying the restriction.

(a) Give an unambiguous algorithmic description of how to take (S, E) and convert it to a corresponding G that completes the mapping reduction successfully.

(b) Use your mapping reduction to create a graph for the following input for the problem SAFE-CLASSES:

S = { Ariel, Binh, Carthik, Deanna, Eduardo, Freddy }

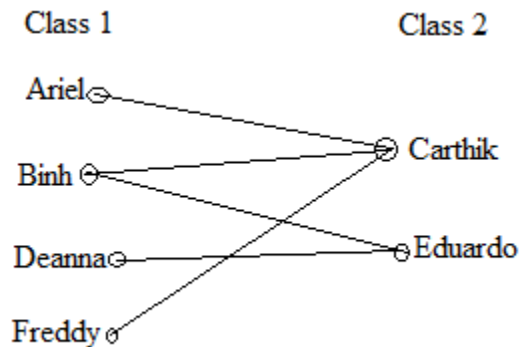
E = { (Ariel, Carthik), (Binh, Carthik), (Binh, Eduardo), (Carthik, Freddy), (Deanna, Eduardo) }

(c) Using your graph or otherwise, give a valid listing of two classes of students which proves that the input above belongs to the language SAFE-CLASSES

Solution

(a) For each student in S, create a vertex in the graph G. For each pair of students in E, add an edge in G between the pair of corresponding vertices.

(b) Here is the corresponding drawing for G created by the mapping reduction described in (a):



(c) From the drawing above, we can see that one valid way to split the classes is to have {Ariel, Binh, Deanna and Freddy} in one class and {Carthik, Eduardo} in the other.

Grading Criteria: (a) +2 for ANY valid graph construction

+2 for the correct number of vertices, 1 per student

+2 for adding edges for each pair that can't be together.

(b) +4 for any drawing consistent with the description, give partial as needed

(c) +2 for a valid class split, but I think this might be the only one that works.

7) (3 pts) Who invented the idea of the Turing Machine?

Alan Turing, give to all