Assignment # 8.1 Sample Key

1. Use reduction from Halt to show that one cannot decide REPEATS, where REPEATS = { f | for some x and y, x ≠ y, $\varphi_f(x)\downarrow$, $\varphi_f(y)\downarrow$ and $\varphi_f(x) == \varphi_f(y)$ }

Let f,x be an arbitrary pair of natural numbers. <f,x> is in Halt iff $\varphi_f(x)\downarrow$

Define g by $\varphi_g(y) = \varphi_f(x) - \varphi_f(x)$, for all y.

Clearly, $\varphi_g(y) = 0$, for all y, iff $\varphi_f(x)\downarrow$, and $\varphi_g(y)\uparrow$, for all y, otherwise.

Summarizing, <f,x> is in Halt implies g is in REPEATS and <f,x> is not in Halt implies g is not in REPEATS

Halt $\leq_m$ REPEATS as we were to show.

Note: I have not overloaded the index of a function with the function in my proof, but I do not mind if you do such overloading.

Assignment # 8.2 Sample Key

2. Show that REPEATS reduces to Halt. (1 plus 2 show they are equally hard)

Let f be an arbitrary natural number. f is in REPEATS iff for some x and y, $x \neq y$, $\varphi_f(x)\downarrow$, $\varphi_f(y)\downarrow$ and $\varphi_f(x) == \varphi_f(y)$

Define g by $\varphi_g(z) = \exists \langle x,y,t\rangle$ [STP(f,x,t) & STP(f,y,t) & ($x\neq y$) & (VALUE(f,x,t) = (VALUE(f,y,t))], for all z.

Clearly, $\varphi_g(z) = 1$, for all z, iff there is some pair, x,y, such that $\varphi_f(x)\downarrow$ and $\varphi_f(y)\downarrow$ and $\varphi_f(x) = \varphi_f(y)$, and $\varphi_g(z)\uparrow$, for all z, otherwise.

Summarizing, f is in REPEATS iff <g,0> is in Halt and so

REPEATS $\leq_m$ Halt as we were to show.

Assignment # 8.3 Sample Key

3. Use Reduction from Total to show that DOUBLES is not even re, where
   DOUBLES = { f | for all x, $\varphi_f(x)\downarrow$, $\varphi_f(x+1)\downarrow$ and $\varphi_f(x+1)=2*\varphi_f(x)$ }

Let f be an arbitrary natural number. f is in Total iff $\forall x\ \varphi_f(x)\downarrow$

Define g by $\varphi_g(x) = \varphi_f(x) - \varphi_f(x)$, for all x.

Clearly, $\varphi_g(x) = 0$, and so $\varphi_g(x+1) = 2*\varphi_g(x) = 0$ for all x, iff $\forall x\ \varphi_f(x)\downarrow$; otherwise $\varphi_g(x)\uparrow$ for some x.

Summarizing, f is in Total iff g is in DOUBLES and so

TOTAL $\leq_m$ DOUBLES as we were to show.

Assignment # 8.3 Alternate Key

3. Use Reduction from Total to show that DOUBLES is not even re, where
   DOUBLES = { f | for all x, $\varphi_f(x)\downarrow$, $\varphi_f(x+1)\downarrow$ and $\varphi_f(x+1)=2*\varphi_f(x)$ }

Let f be an arbitrary natural number. f is in Total iff $\forall x\ \varphi_f(x)\downarrow$

Define g by $\varphi_g(x) = \varphi_f(x) - \varphi_f(x)+2$^x for all x.

Clearly, $\varphi_g(x) = 2$^x, and so $\varphi_g(x+1) = 2*\varphi_g(x) = 2$^(x+1) for all x, iff $\forall x\ \varphi_f(x)\downarrow$; otherwise $\varphi_g(x)\uparrow$ for some x.

Summarizing, f is in Total iff g is in DOUBLES and so

TOTAL $\leq_m$ DOUBLES as we were to show.

Assignment # 8.4 Sample Key

4. Show DOUBLES reduces to Total. (3 plus 4 show they are equally hard)

Let f be an arbitrary natural number. f is in DOUBLES iff $\forall x\ \varphi_f(x)\downarrow$, $\varphi_f(x+1)\downarrow$ and $\varphi_f(x+1)=2*\varphi_f(x)$.

Define g by $\varphi_g(x) = \mu y[\varphi_f(x+1) = 2*\varphi_f(x)]$, for all x.

Clearly, $\varphi_g(x)\downarrow$, for all x, iff $\forall x\ \varphi_f(x)\downarrow$, $\varphi_f(x+1)\downarrow$ and $\varphi_f(x+1)=2*\varphi_f(x)$; otherwise $\varphi_g(x)\uparrow$ for some x.

Summarizing, f is in DOUBLES iff g is in Total and so

DOUBLES $\leq_m$ TOTAL as we were to show.

Assignment # 8.5 Sample Key

5. Use Rice's Theorem to show that REPEATS is undecidable

First, REPEATS is non-trivial as C0(x) = 0 is in REPEATS and S(x) = x+1 is not.

Second, REPEATS is an I/O property.

To see this, let f and g are two arbitrary indices such that
$\forall x\, [\varphi_f(x) = \varphi_g(x)]$

f $\in$ REPEATS iff $\exists$ y,z, y $\neq$ z, such that $\varphi_f(y)\downarrow$, $\varphi_f(z)\downarrow$ and $\varphi_f(y) = \varphi_f(z)$ iff, since $\forall x\, [\varphi_f(x) = \forall x\, \varphi_g(x)]$, $\exists$ y,z, y $\neq$ z, (same y,z as above) such that $\varphi_g(y)\downarrow$, $\varphi_g(z)\downarrow$ and $\varphi_g(y) = \varphi_g(z)$ iff g $\in$ REPEATS.

Thus, f $\in$ REPEATS iff g $\in$ REPEATS.

Assignment # 8.6 Sample Key

6. Use Rice's Theorem to show that DOUBLES is undecidable

First, DOUBLES is non-trivial as C0(x) = 0 (2*0 = 0) is in DOUBLES and S(x) = x+1 is not.

Second, DOUBLES is an I/O property.

To see this, let f and g are two arbitrary indices such that $\forall x\ [\varphi_f(x) = \varphi_g(x)]$.

f $\in$ DOUBLES iff for all x, $\varphi_f(x)\downarrow$, $\varphi_f(x+1)\downarrow$ and $\varphi_f(x+1)=2*\varphi_f(x)$ iff, since $\forall x\ [\varphi_f(x) = \varphi_g(x)]$, for all x, $\varphi_g(x)\downarrow$, $\varphi_g(x+1)\downarrow$ and $\varphi_g(x+1)=2*\varphi_g(x)$ iff g $\in$ DOUBLES.

Thus, f $\in$ DOUBLES iff g $\in$ DOUBLES.