**Generally useful information.**

* The notation **z =** **<x,y>** denotes the pairing function with inverses **x =** **<z>1** and **y =** **<z>2**.
* The minimization notation **μ y [P(…,y)]** means the least **y** (starting at **0**) such that **P(…,y)** is true. The bounded minimization (acceptable in primitive recursive functions) notation
**μ y (u≤y≤v) [P(…,y)]** means the least **y** (starting at **u** and ending at **v**) such that **P(…,y)** is true. Unlike the text, I find it convenient to define **μ y (u≤y≤v) [P(…,y)]** to be **v+1**, when no **y** satisfies this bounded minimization.
* The tilde symbol, **~,** means the complement. Thus, set **~S** is the set complement of set **S**, and predicate **~P(x)** is the logical complement of predicate **P(x).**
* The minus symbol, –, when applied to sets is set difference, so **S – T = {x | x∈S && x∉T}.**
* The absolute value, **|z|**, is the magnitude of **z**. Thus, **|x-y|** is the difference between **x** and **y**, when **x** and **y** are both non-negative.
* A function **P** is a predicate if it is a logical function that returns either **1** (**true**) or **0** (**false**). Thus, **P(x)** means **P** evaluates to **true** on **x**, but we can also take advantage of the fact that **true** is **1** and **false** is **0** in formulas like **y × P(x)**, which would evaluate to either **y** (if **P(x)**) or **0** (if **~P(x)**).
* A set **S** is recursive if **S** has a total recursive characteristic function **χS**, such that **x ∈ S ⇔ χS(x)**. Note **χS** is a predicate. Thus, it evaluates to **0** (**false**), if **x ∉ S**.
* When I say a set **S** is re, unless I explicitly say otherwise, you may assume any of the following equivalent characterizations:
1. **S** is either empty or the range of a total recursive function **fS**.
2. **S** is the domain of a partial recursive function **gS**.
3. **S** is recognizable by a Turing Machine.
* If I say a function **g** is partially computable, then there is an index **g** (I know that’s overloading, but that’s okay as long as we understand each other), such that **Φg(x) = Φ(g, x) = g(x)**. Here **Φ** is a universal partially recursive function.
Moreover, there is a total recursive function **STP**, such that
**STP(g. x, t)** is **1** (true), just in case **g**, started on **x**, halts in **t** or fewer steps.
**STP(g. x, t)** is **0** (false), otherwise.
Finally, there is another total recursive function **VALUE**, such that
**VALUE(g. x, t)** is **g(x)**, whenever **STP(g. x, t)**.
**VALUE(g. x, t)** is defined but meaningless if **~STP(g. x, t)**.
* The notation **f(x)↓** means that **f** converges when computing with input **x**, but we don’t care about the value produced. In effect, this just means that **x** is in the domain of **f**.
* The notation **f(x)↑** means **f** diverges when computing with input **x**. In effect, this just means that **x** is **not** in the domain of **f**.
* The **Halting Problem** for any effective computational system is the problem to determine of an arbitrary effective procedure **f** and input **x**, whether or not **f(x)↓**. The set of all such pairs is a classic re non-recursive one. The set of all such **<f,x>** is denoted **K0**. A related set **K** is the set of all **f** that halt on their own indices. Thus, **K = { f | Φf(f) ↓ }** and **K0 = {<f,x> |Φf(x)↓ }**
* The **Uniform Halting Problem** is the problem to determine of an arbitrary effective procedure **f**, whether or not **f** is an algorithm (halts on all input). The set of all such function indices is a classic non re one and is often called **TOTAL**.

**COT 4210 Fall 2014 Final Exam Sample Questions**

 **1**. Let set **A** be recursive, **B** be re non-recursive and **C** be non-re. Choosing from among **(REC)** **recursive**, **(RE)** **re non-recursive**, **(NR)** **non-re**, categorize the set **D** in each of a) through d) by listing **all** possible categories. No justification is required.

**a.) D = ~C**

**b.) D ⊆ (A∪ C)**

**c.) D = ~B**

**d.) D = B − A**

 **2.** Prove that the **Halting Problem** (the set **K0 )** is not recursive (decidable) within any formal model of computation. (Hint: A diagonalization proof is required here.)

 **3.** Using reduction from the known undecidable **HasZero, HZ = { f | ∃x f(x) = 0 }**, show the non-recursiveness (undecidability) of the problem to decide if an arbitrary primitive recursive function **g** has the property **IsZero, Z = { f | ∀x f(x) = 0 }**.

 **4**. Choosing from among **(D)** **decidable**, **(U)** **undecidable**, **(?)** **unknown**, categorize each of the following decision problems. No proofs are required.

|  |  |  |
| --- | --- | --- |
| **Problem / Language Class** | **Regular** | **Context Free** |
| L = Σ\* ? |  |  |
| **L = φ ?** |  |  |
| **x ∈ L2, for arbitrary x ?** |  |  |

 **5**. Choosing from among **(Y)** **yes**, **(N)** **No**, **(?)** **unknown**, categorize each of the following closure properties. No proofs are required.

|  |  |  |
| --- | --- | --- |
| **Problem / Language Class** | **Regular** | **Context Free** |
| Closed under intersection? |  |  |
| **Closed under quotient?** |  |  |
| **Closed under quotient with Regular languages?** |  |  |
| **Closed under complement?** |  |  |

 **6**. Prove that any class of languages, ***C***, closed under union, concatenation, intersection with regular languages, homomorphism and substitution (e.g., the Context-Free Languages) is closed under **MissingMiddle**, where, assuming L is over the alphabet **Σ**,
**MissingMiddle(L) = { xz | ∃y ∈ Σ\* such that xyz ∈ L }**You must be very explicit, describing what is produced by each transformation you apply.

 **7.** Use **PCP** to show the undecidability of the problem to determine if the intersection of two context free languages is non-empty. That is, show how to create two grammars **GA** and **GB** based on some instance **P = <<x1,x2,…,xn>, <y1,y2,…,yn>>** of **PCP**, such that **L(GA) ∩ L(GB) ≠ φ** iff **P** has a solution. Assume that **P** is over the alphabet **Σ**.You should discuss what languages your grammars produce and why this is relevant, but no formal proof is required.

 **8.** Consider the set of indices **CONSTANT = { f | ∃K ∀y [ ϕf(y) = K ] }**. Use Rice’s Theorem to show that **CONSTANT** is not recursive. Hint: There are two properties that must be demonstrated.

 **9.** Show that **CONSTANT ≡m TOT**, where **TOT = { f | ∀y ϕf(y)↓ }**.

 **10.** Why does Rice’s Theorem have nothing to say about each of the following? Explain by showing some condition of Rice’s Theorem that is not met by the stated property.

 ***a.***) **AT-LEAST-LINEAR = { f | ∀y ϕf(y) converges in no fewer than y steps }**.

 ***b.***) **HAS-IMPOSTER = { f | ∃ g [ g≠f** and **∀y [ ϕg(y) = ϕf(y) ] ] }**.

 **11.** We described the proof that **3SAT** is polynomial reducible to Subset-Sum.

a.) Describe **Subset-Sum**

b.) Show that **Subset-Sum** is in **NP**

c.) Assuming a **3SAT** expression **(a + ~b + c) (b + b + ~c)**, fill in the upper right part of the reduction from **3SAT** to **Subset-Sum**.

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
|  | **a** | **b** | **c** | **a + ~b + c** | **b + b + ~c** |
| **a** | **1** |  |  |  |  |
| **~a** | **1** |  |  |  |  |
| **b** |  | **1** |  |  |  |
| **~b** |  | **1** |  |  |  |
| **c** |  |  | **1** |  |  |
| **~c** |  |  | **1** |  |  |
| **C1** |  |  |  | **1** |  |
| **C1’** |  |  |  | **1** |  |
| **C2** |  |  |  |  | **1** |
| **C2’** |  |  |  |  | **1** |
|  | **1** | **1** | **1** | **3** | **3** |

**12.** Use the appropriate Pumping Lemmas to show:

a.) **{ ww | w is over {a,b} }** is not Regular

b.) **{ ww | w is over {a,b} }** is not Context Free

**13.** Write a context-free grammar for the complement of the language **{ ww | w is in {a,b}\* }**