

# Sample Assignment # 5.1

For each of the following, prove it is not regular by using the Pumping Lemma or Myhill-Nerode. You must do one of these using the Pumping Lemma and one using Myhill-Nerode.

- a.  $\{ a^{k!} \mid k > 0 \}$  This is set  $\{ a^1, a^2, a^6, a^{24}, a^{120}, \dots \}$
- b.  $\{ a^i b^j c^k \mid i \geq 0, j \geq 0, k \geq 0, j = i + k \}$

# Assignment # 5.1 Answer

1a.  $\{ a^k \mid k > 0 \}$  using P.L.

1. Assume that L is regular
2. Let N be the positive integer given by the Pumping Lemma
3. Let s be a string  $s = a^{(N+1)!} \in L$
4. Since  $s \in L$  and  $|s| \geq N$ , s is split by PL into xyz, where  $|xy| \leq N$  and  $|y| > 0$  and for all  $i \geq 0$ ,  $xy^iz \in L$
5. We choose  $i = 2$ ; by PL:  $xy^2z = xyyz \in L$
6. Thus,  $a^{(N+1)!+|y|}$  would be  $\in L$ . This means that there is a factorial between  $(N+1)!$  and  $(N+1)!+N$ , but the smallest factorial after  $(N+1)!$  is  
 $(N+2)! = (N+2)(N+1)! = N(N+1)! + 2(N+1)! > (N+1)! + 2N > (N+1)!+N$
7. This is a contradiction, therefore L is not regular ■
8. Note: Using N is dangerous because N could be 1 and  $2!$  is within N (1) of  $1!$

1b.  $\{ a^i b^j c^k \mid i \geq 0, j \geq 0, k \geq 0, j = i + k \}$  using P.L.

1. Assume that L is regular
2. Let N be the positive integer given by the Pumping Lemma
3. Let s be the string  $s = a^N b^N \in L$
4. Since  $s \in L$  and  $|s| \geq N$ , s is split by PL into xyz, where  $|xy| \leq N$  and  $|y| > 0$  and for all  $i \geq 0$ ,  $xy^iz \in L$
5. We choose  $i = 0$ ; by PL:  $xz = xz \in L$
6. Thus,  $a^{N-|y|} b^N$  would be  $\in L$ , but it's not since  $N-|y| + 0 < N$ . Note: The 0 is because there are 0 c's
7. This is a contradiction, therefore L is not regular ■

# Assignment # 5.1 Answer

1a.  $\{ a^k \mid k > 0 \}$  using M.N.

We consider the collection of right invariant equivalence classes  $[a^{j-j}]$ ,  $j \geq 0$ .

It's clear that  $a^{j-j}a^j$  is in the language, but  $a^{k-j}a^j$  is not when  $j < k$

This shows that there is a separate equivalence class  $[a^{j-j}]$  induced by  $R_L$ , for each  $j \geq 0$ .

Thus, the index of  $R_L$  is infinite and Myhill-Nerode states that  $L$  cannot be Regular. ■

1b.  $\{ a^i b^j c^k \mid i \geq 0, j \geq 0, k \geq 0, j = i + k \}$  using M.N.

We consider the collection of right invariant equivalence classes  $[a^j]$ ,  $j \geq 0$ .

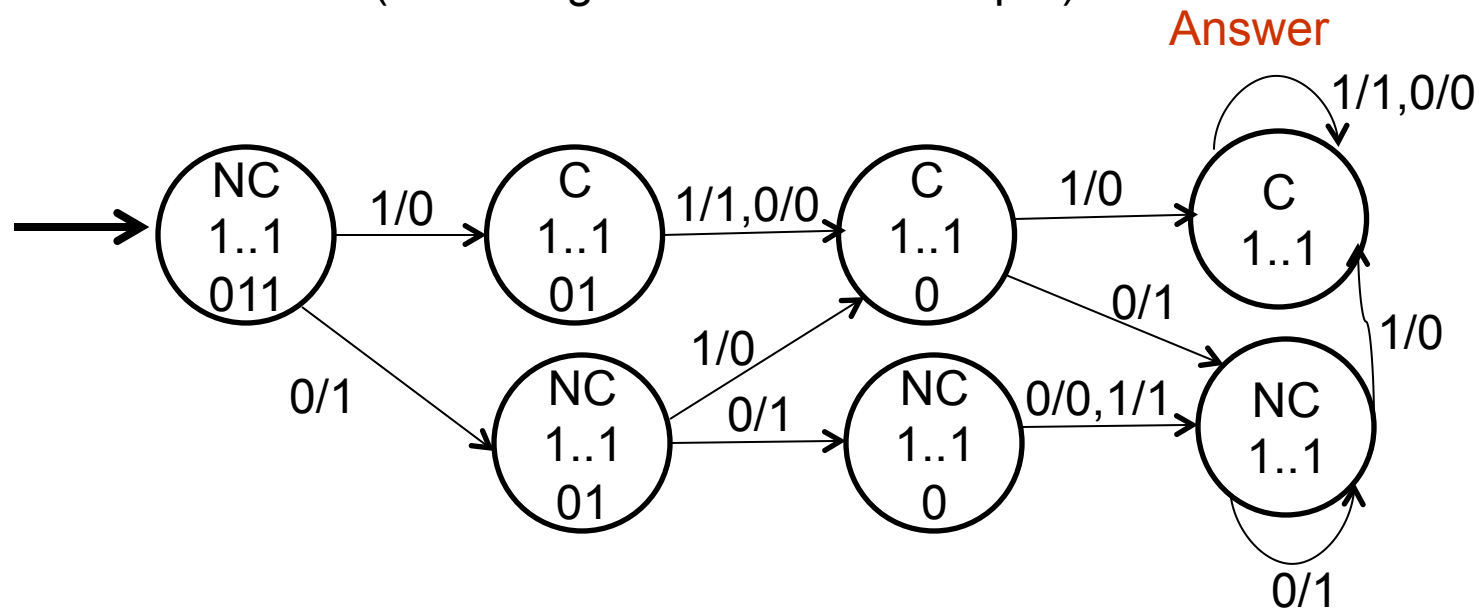
It's clear that  $a^i b^j$  is in the language, but  $a^k b^j$  is not when  $j \neq k$

This shows that there is a separate equivalence class  $[a^j]$  induced by  $R_L$ , for each  $j \geq 0$ .

Thus, the index of  $R_L$  is infinite and Myhill-Nerode states that  $L$  cannot be Regular. ■

# Assignment # 5.2

Write a Mealy finite state machine that produces the 2's complement result of subtracting 101 from a binary input stream (assuming at least 3 bits of input)



# Assignment # 5.3

Write a regular (right linear) grammar that generates the set of strings denoted by the regular expression  $(0+11)^* (101 (00 + 1)^*)^*$

$G = (\{<0+11>, <1\_1>, <1\_01>, <10\_1>, <00+11>, <0\_0>\}, \{0,1\}, R, <0+11>)$

$<0+11> \rightarrow 0<0+11> \mid 1<1\_1> \mid 1<1\_01> \mid \lambda$

$<1\_1> \rightarrow 1<0+11>$

$<1\_01> \rightarrow 0<10\_1>$

$<10\_1> \rightarrow 1<00+1>$

$<00+11> \rightarrow 0<0\_0> \mid 1<00+11> \mid \lambda \mid 1<1\_01>$

$<0\_0> \rightarrow 0<00+11>$

Or

$G = (\{S,T,U,V,W,X\}, \{0,1\}, R, S)$

$S \rightarrow 0S \mid 1T \mid 1U \mid \lambda$

$T \rightarrow 1S$

$U \rightarrow 0V$

$V \rightarrow 1W$

$W \rightarrow 0X \mid 1W \mid \lambda \mid 1U$

$X \rightarrow 0W$