## Assignment # 2.1 Key

Let L be a language over  $\{a,b\}$  where every string is of even length and is of the form WX, where |W|=|X| but  $W\neq X$ . Design and present an algorithm that recognized strings in L using no unbounded amount of storage (no stacks, no queues). This means that any memory required must be of a fixed size independent of the length of an input string. Note: You cannot play the game of using unbounded recursion, as each call consumes stack space.

You can attack this deterministically or non-deterministically. I will do deterministically. Consider any string z=WX, |W|=|X| but  $W\neq X$ . Such a string need only have one transcription error when copying W as X to be in L. check fits the bill.

```
int check(const char z){ int p = -1; int odd = 0; int mid;
```

```
while (z[++p]) { odd = 1-odd; }
if ( odd ) return 0;
mid = p/2;
for (p=0; p<mid; p++) if (z[p] != z[mid+p]) return 1;
return 0;</pre>
```

}

## Assignment # 2.2 Key

Present a language L over  $\Sigma = \{a\}$  where  $L^4 = L^5$  but  $L \neq L^2$ ,  $L^2 \neq L^3$  and  $L^3 \neq L^{34}$ Note:  $L^k = \{x_1x_2...x_k \mid x_1,x_2,...,x_k \in L\}$ . This is basically a giveaway, since I showed exactly how to do it.

Proof:

Consider L =  $\{a\}^*$  -  $\{aa, aaa, aaaa\}$ 

 $L^2 = \{a\}^* - \{aaa, aaaa\}$  since the presence of the empty string in  $\{a\}^*$  means all strings in L are in L<sup>2</sup>. Additionally, aa = a ° a and so aa is in L<sup>2</sup> but aaa and aaaa are not since they cannot be formed from any pair of members in L  $L^3 = \{a\}^* - \{aaaa\}$  since the presence of the empty string in  $\{a\}^*$  means all strings in L are in L<sup>3</sup>. Additionally, aaa = aa ° a and so aaa is in L<sup>3</sup> but aaaa is not since it cannot be formed from any triple of members in L  $L^4 = \{a\}^*$  since the presence of the empty string in  $\{a\}^*$  means all strings in L<sup>3</sup> are in L<sup>3</sup>. Additionally, aaaa = aaa ° a and so aaaa is in L<sup>4</sup>  $L^4 = L^5$  since L<sup>4</sup> is already  $\{a\}^*$  and so nothing new can be created and the

presence of the empty string in  $\{a\}^*$  means all in L<sup>4</sup> are in L<sup>5</sup>