Intro to Discrete Structures Lecture 13

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The Euclidean Algorithm

Lemma 1: Let
$$a=bq+r$$
, where a,b,q , and r are integers. Then

$$\gcd(a,b) = \gcd(b,r).$$

$$d \mid a \wedge d \mid b \Rightarrow d \mid (ma+nb)$$

$$r = a - bq$$

$$\gcd(a,b) \mid \gcd(b,r)$$

$$q \in A \mid d \mid r$$

$$\gcd(a,b) \mid \gcd(b,r)$$

$$q \in A \mid d \mid r$$

$$\gcd(a,b) \mid \gcd(b,r)$$

The Euclidean Algorithm

$$\frac{f \mid g \land g \mid f}{g \bowtie (a,b) \mid g \bowtie (b,r) \land}$$

$$g \bowtie (a,b) \mid g \bowtie (a,b) \Rightarrow g \bowtie (a,b) = g \bowtie (b,r)$$

$$g \bowtie (a,b) \mid a \land g \bowtie (a,b) \mid b \Rightarrow g \bowtie (a,b) \mid b \Rightarrow g \bowtie (a,b) \mid b \Rightarrow g \bowtie (a,b) \mid r \Rightarrow g \bowtie (a,b) \mid r \Rightarrow g \bowtie (a,b) \mid r \Rightarrow g \bowtie (a,b) \mid g \bowtie (b,r)$$

$$\Rightarrow g \bowtie (a,b) \mid r \Rightarrow g \bowtie (a,b) \mid g \bowtie (b,r)$$

The Euclidean Algorithm

Example 12: Find $gcd(414,\underline{662})$ using the Euclidean Algorithm.

$$\frac{662}{a} = \frac{414 \cdot 1}{b} + \frac{248}{7} \qquad g\omega(662,414) = g\omega(444,248)$$

$$414 = 248 \cdot 1 + 166$$

$$248 = 166 \cdot 1 + 82$$

$$166 = 82 \cdot 2 + 21$$

$$82 = 2 \cdot 41 + 0 \implies g\omega(662,414) = 21$$

Some Useful Facts

Theorem 1:

$$\forall a \, \forall b \, \exists s \, \exists t \, \gcd(a, b) = sa + tb$$
.

The pair (s,t) can be efficiently computed with the extended Euclidean algorithm.

The Extended Euclidean Algorithm

Example 1: Express gcd(252, 198) = 18 as a linear combination of 252 and 198.

$$36 = 198 - 3.54$$

$$198 = 3.54 + 36$$

$$18 = 54 - 36$$

$$36 = 1.36 + 18$$

$$36 = 2.18 + 0$$

$$9 \text{ and } (36, 18) = 9 \text{ and } (18, 0) = 18$$

Some Useful Facts

Lemma 2: If p is a prime and $p \mid a_1 a_2 \cdots a_n$, where each a_i is an integer, then $p \mid a_i$ for some i.

$$6 | 12$$

$$12 = 24 4 \cdot 3$$

$$6 = 2 \cdot 3$$

$$2 \cdot 3 | 4 \cdot 3$$

Some Useful Facts

Lemma 1:

$$\gcd(a,b) = 1 \land a \mid bc \Rightarrow a \mid c$$

Proof of the uniqueness of the prime factorization of a positive integer:

$$gcd(3,5)=1$$

Mathematical Induction

Mathematical induction is based on the rule of inference

1.
$$P(1)$$

2. $\forall k (P(k) \rightarrow P(k+1))$
3. $\therefore \forall n P(n)$

which is true for the domain of natural numbers \mathbb{N} .

Climbing an Infinite Ladder

- 1. We can reach the first rung of the ladder.
- 2. If we can reach a particular rung of the ladder, then we can reach the next rung of the ladder.
- 3. Therefore, we can reach any rung of the ladder.

Principle Mathematical Induction

To prove that P(n) is true for all natural numbers n, where P(n) is a propositional function, we complete two steps:

- **Description** Basis step: We verify that P(1) is true.
- We show that the conditional statement

$$P(k) \rightarrow P(k+1)$$

is true for all natural numbers k.

Warning

- In a proof of mathematical induction is is **not** assumed that P(k) is true for all k.
- It is only shown that **if it is assumed** that P(k) is true, then P(k+1) is also true.
- Thus, a proof of mathematical induction is not a case of begging the question, or circular reasoning.

Example: Show that

$$P(\Pi): \sum_{i=1}^{n} i = \frac{n(n+1)}{2}. \text{ is true.}$$
 Basis slep
$$P(I) \quad I = \sum_{i=1}^{l} i = \frac{I(I+1)}{2} = I$$

Inductive Step
$$\forall k P(k) \longrightarrow P(k+1)$$

assume that $P(k)$ is time
$$\sum_{k=1}^{K} i = \frac{k(k+1)}{2!} \qquad P(k) \longrightarrow P(k+1)$$

$$\sum_{k=1}^{K} i = \sum_{k=1}^{K} i + (k+1)$$

$$= \frac{k(k+1)}{2!} + (k+1) = \frac{k(k+1)+2(k+1)}{2!}$$

$$= \frac{(k+1)(k+2)}{2!}$$

Example 2: Conjecture a formula for the sum of the first n positive integers. Then prove your conjecture using mathematical induction.

$$\sum_{i=1}^{k+1} i = \frac{(k+1)(k+2)}{2i} \iff P(k+1) \text{ holds}$$

The Number of Subsets of a Finite Set

The Number of Subsets of a Finite Set

DeMorgan for Intersection

Example 10: Prove

$$\bigcap_{j=1}^{n} A_j = \bigcup_{j=1}^{n} \overline{A_j}$$

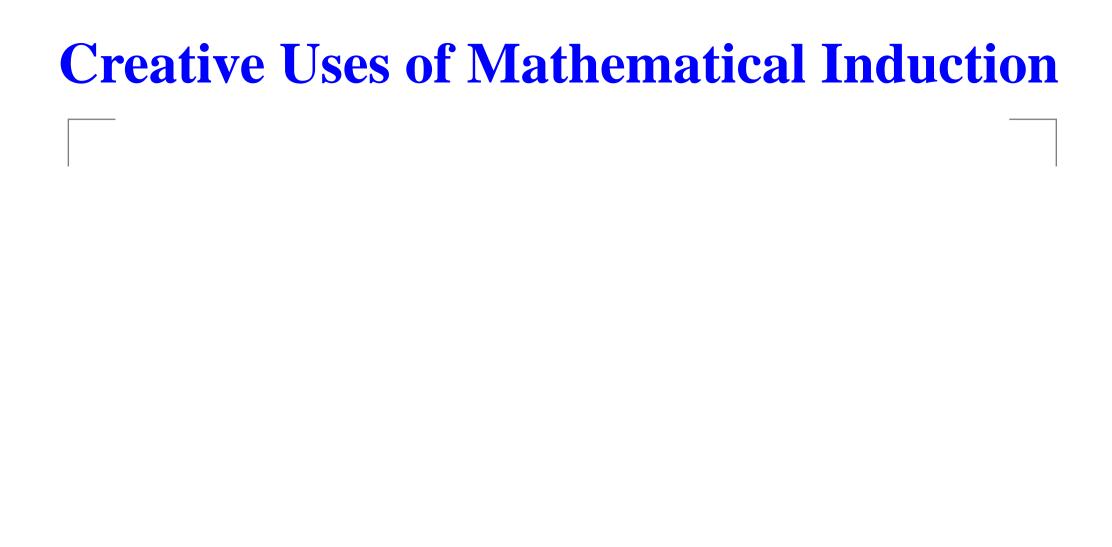
whenever A_1, A_2, \ldots, A_n are subset of a universal U and $n \geq 2$.

DeMorgan for Intersection

DeMorgan for Intersection

Creative Uses of Mathematical Induction

Example: Show that every $2^n \times 2^n$ checkerboard with one square removed can be tiled using L-shaped triominoes.



Strong Induction

Strong induction is based on the rule of inference

1.
$$P(1)$$

2. $\forall k (\wedge_{j=1}^k P(j) \rightarrow P(k+1))$
3. $\therefore \forall n P(n)$

which is true for the domain of natural numbers \mathbb{N} .

Strong Induction

To prove that P(n) is true for all natural numbers n, where P(n) is a propositional function, we complete two steps:

- **Description** Basis step: We verify that P(1) is true.
- We show that the conditional statement

$$\bigwedge_{j=1}^{k} P(j) \to P(k+1)$$

is true for all natural numbers k.

Existence of Prime Factorization

Example: Show that if n is an integer greater than 1, then n can be written as the product of primes.

Existence of Prime Factorization

Existence of Prime Factorization

Example 2: Show that if n is an integer greater than 1, then n can be written as the product of primes.

Winning Strategy

Example 3:

Winning Strategy

Winning Strategy

$$d = \gcd(a,b)$$

$$d \mid a \iff \exists k \quad a = k \cdot d \text{ multiples}$$

$$d \mid b \iff \exists l \quad b = l \cdot d \text{ of } d$$

$$let m, n \in \mathbb{Z} \text{ be arbihary.}$$

$$ma + \pi b = mkd + \pi ld$$

$$= (mk + \pi l) \cdot d$$

$$\iff d \mid (ma + \pi b)$$

 $\Rightarrow d | e \wedge d | f$ $\Rightarrow d | g \omega d (e, f)$