

### COT 3100 — Recitation 3 — Logic makes the world go 'round

- Below are the truth tables for the operators NAND, NOR, and XOR. Write definitions for these operators using only AND, OR, and NOT. Try to explain in words what each one means.

| A | B | A NAND B | A | B | A NOR B | A | B | A XOR B |
|---|---|----------|---|---|---------|---|---|---------|
| T | T | F        | T | T | F       | T | T | F       |
| T | F | T        | T | F | F       | T | F | T       |
| F | T | T        | F | T | F       | F | T | T       |
| F | F | T        | F | F | T       | F | F | F       |

- The NAND operator, described above, is a *sole sufficient operator*, which means that all the other operators can be rewritten using only NAND. This is an important result in chip design, for example, because it can make chips cheaper to build if they use all NAND gates. How would you translate “not P” using only combinations of the NAND operator?
- Use the laws of logic to prove the following. Remember, for each line of your proof, you should include the rule and lines that support it.

$$\begin{array}{l}
 P \wedge \neg Q \\
 R \vee S \\
 S \rightarrow Q \\
 \hline
 \therefore R \vee U
 \end{array}$$

Translate the following three stories into formal proofs. For each one, decide what the primitive statements should be, then describe the situation in logical symbols by writing a system of propositions using those primitives. *Next*, pick one story and use a truth table to determine its validity. *Finally*, pick a different story, and determine its validity by applying the laws of logic to its proposition system.

- If I don't practice regularly, I won't make it onto *American Idol*. But I am practicing regularly, so I will make it.
- Heads I win, tails you lose. We cannot both win and we will not both lose. Therefore I will certainly win.
- Distributed database  $D$  has two client processes,  $P_1$  and  $P_2$ . In order to complete transaction  $T$  in this database, a process must have an *exclusive lock* on each of the two data rows  $R_1$  and  $R_2$ . If one process has an exclusive lock on anything, then the other process cannot use it.  $P_1$  and  $P_2$  both need to complete transaction  $T$ . Process  $P_1$  starts by taking an exclusive lock on  $R_1$ . At the same time, process  $P_2$  gets an exclusive lock on  $R_2$ . Now neither  $P_1$  nor  $P_2$  can complete the transaction.

Finished early? Why not translate AND and OR using only NAND...