

**COT 3100 Recitation #2: Logic Practice**  
**1/17-20/2017**

**Warm-Up Problems**

1) Al gets the disease algebritis and must take one green pill and one pink pill each day for two weeks. A green pill costs \$1 more than a pink pill and Al's pills cost a total of \$546 for the two weeks. How much does one green pill cost?

**Solution**

Let the green pill cost  $g$  dollars. Then the pink pill costs  $g - 1$  dollars. Thus we have:

$$\begin{aligned}14(g + g - 1) &= 546 \\2g - 1 &= 39 \\g &= 20\end{aligned}$$

One green pill costs **\$20**.

2) The second and fourth terms of a geometric sequence are 2 and 6, respectively. What are the possible values of the first term in the sequence?

**Solution**

Let  $r$  be the common ratio of the sequence. The given information indicates that  $r^2 = \frac{6}{2} = 3$ . It follows that  $r = \pm\sqrt{3}$ . The two possible values of the first term are simply two (the second term) divided by the two possible common ratios:

$$\frac{2}{\pm\sqrt{3}} = \pm \frac{2\sqrt{3}}{3}$$

3) Let  $s(x)$  denote the sum of the digits of the positive integer  $x$ . For example  $s(8) = 8$  and  $s(123) = 1 + 2 + 3 = 6$ . For how many two-digit values of  $x$  is  $s(s(x)) = 3$ ?

**Solution**

The question is simply asking for how many two-digit numbers,  $x$ , is the sum of digits  $y$ , where the sum of digits of  $y$  is 3. The possible values for  $y$  are 3, 12, 21 and 30. Thus, we must list all two-digit numbers that have a sum of digits of 3, 12, 21 or 30. The maximum sum of two digits is 18, so we must simply find all two-digit numbers whose sum of digits is either 3 or 12. These are 12, 21, 30, 39, 48, 57, 66, 75, 84, and 93. Thus, there are **10** such values.

4) Cassandra sets her watch to the correct time at noon. At the actual time of 1:00 PM, she notices that her watch reads 12:57 and 36 seconds. Assuming that her watch loses time at a constant rate, what will be the actual time when her watch first reads 10:00 PM?

**Solution**

There are 3600 seconds in an hour. Cassandra's watch ticks  $57 \times 60 + 36 = 3456$  seconds in this time. Thus, the actual time moves at a rate of  $\frac{3600}{3456} = \frac{225}{216}$  compared to Cassandra's clock. Thus, the number of actual seconds that elapse by the time Cassandra's clock reads 10:00 PM (when it had ticked  $10 \text{ hr} \times \frac{3600 \text{ ticks}}{\text{hr}} = 36000 \text{ ticks}$ ), can be calculated as follows:

$$\frac{225 \text{ sec}}{216 \text{ ticks}} \times 36000 \text{ ticks} = 225 \times \frac{1000}{6} \text{ sec} \times \frac{1 \text{ hr}}{3600 \text{ sec}} = 10 \frac{5}{12} \text{ hr}$$

Since there are 60 minutes in an hour, this equates to **10:25 PM** for the actual time.

5) If  $\log(xy^3) = 1$  and  $\log(x^2y) = 1$  what is  $\log(xy)$ ?

**Solution**

Let  $a = \log(x)$  and  $b = \log(y)$ . We seek  $\log(xy) = \log(x) + \log(y) = a + b$ .

Write both equations out and simplify:

$$\begin{aligned} \log(xy^3) &= 1 \\ \log(x) + \log(y^3) &= 1 \\ \log(x) + 3\log(y) &= 1 \\ a + 3b &= 1 \end{aligned}$$

$$\begin{aligned} \log(x^2y) &= 1 \\ \log(x^2) + \log(y) &= 1 \\ 2\log(x) + \log(y) &= 1 \\ 2a + b &= 1 \end{aligned}$$

Solve the system for  $b$ :

$$\begin{array}{r} 2a + 6b = 2 \\ 2a + b = 1 \\ \hline 5b = 1, b = \frac{1}{5} \end{array}$$

Since  $a + 3b = 1$ ,  $a + 3(\frac{1}{5}) = 1$ , so  $a = \frac{2}{5}$ . Finally, the sum of  $a$  and  $b$ , our desired quantity, is  $\frac{1}{5} + \frac{2}{5} = \frac{3}{5}$ .

## Logic Problems

6) Using the following premises:

$$(p \wedge t) \rightarrow (r \vee s)$$

$$q \rightarrow (u \wedge t)$$

$$u \rightarrow p$$

$$\bar{s}$$

Derive the conclusion  $q \rightarrow r$ .

### Solution

1. $q \rightarrow (u \wedge t)$ ,	Given
2. $\bar{q} \vee (u \wedge t)$ ,	Definition of Implication
3. $(\bar{q} \vee u) \wedge (\bar{q} \vee t)$ ,	Distributive Law
4. $(\bar{q} \vee u)$ ,	Conjunctive Simplification with step 3
5. $q \rightarrow u$ ,	Definition of Implication with step 4
6. $u \rightarrow p$ ,	Given
7. $q \rightarrow p$ ,	Law of Syllogism with steps 5, 6
8. $(\bar{q} \vee t)$ ,	Conjunctive Simplification with step 3
9. $q \rightarrow t$ ,	Definition of Implication with step 8
10. $(\bar{q} \vee p)$ ,	Definition of Implication with step 7
11. $(\bar{q} \vee t) \wedge (\bar{q} \vee p)$ ,	Conjunction with steps 8 and 10
12. $\bar{q} \vee (p \wedge t)$ ,	Distributive Property with step 11
13. $q \rightarrow (p \wedge t)$ ,	Definition of Implication with step 12
14. $(p \wedge t) \rightarrow (r \vee s)$	Given
15. $q \rightarrow (r \vee s)$ ,	Law of Syllogism with steps 13, 14
16. $\bar{q} \vee (r \vee s)$	Definition of Implication with step 15
17. $(\bar{q} \vee r) \vee s$ ,	Associative Law step 16
18. $(q \rightarrow r) \vee s$	Definition of Implication with step 17
19. $\bar{s}$	Premise
20. $q \rightarrow r$	Rule of Disjunctive Syllogism with steps 18, 19

7) In class, Modus Ponens was proved using just the laws of logic. Prove Modus Tollens in the same manner.

**Solution**

We must show that  $(\bar{q} \wedge (p \rightarrow q)) \rightarrow \bar{p}$  is a tautology via the Laws of Logic.

$(\bar{q} \wedge (p \rightarrow q)) \rightarrow \bar{p},$	Given
$(\bar{q} \wedge (p \rightarrow q)) \vee \bar{p},$	Definition of Implication
$(\bar{q} \vee \overline{p \rightarrow q}) \vee \bar{p},$	De Morgan's Law
$(q \vee \overline{p \vee q}) \vee \bar{p},$	Double Negation, Definition of Implication
$(q \vee (\bar{p} \wedge \bar{q})) \vee \bar{p},$	De Morgan's Law
$(q \vee (p \wedge \bar{q})) \vee \bar{p},$	Double Negation
$((q \vee p) \wedge (q \vee \bar{q})) \vee \bar{p},$	Distributive Law
$((q \vee p) \wedge T) \vee \bar{p},$	Negation Laws
$(q \vee p) \vee \bar{p},$	Identity Law
$q \vee (p \vee \bar{p}),$	Associative Law
$q \vee T,$	Negation Law
$T,$	Domination Law

8) Simplify the following logical expression as much as possible using the laws of logic only. Show each step and state which rule is being used. (Note: You may combine both associative and commutative in a single step, so long as you do so properly.)

$$p \vee [p \wedge [\neg(\neg r \vee \neg q) \vee (\neg r \wedge q)]]$$

**Solution**

$p \vee [p \wedge [\neg(\neg r \vee \neg q) \vee (\neg r \wedge q)]]$	$\leftrightarrow p \vee [p \wedge [(\neg \neg r \wedge \neg \neg q) \vee (\neg r \wedge q)]]$	DeMorgan's
	$\leftrightarrow p \vee [p \wedge [(r \wedge q) \vee (\neg r \wedge q)]]$	Double Negation
	$\leftrightarrow p \vee [p \wedge [(r \vee \neg r) \wedge q]]$	Distributive
	$\leftrightarrow p \vee [p \wedge [T \wedge q]]$	Inverse Law
	$\leftrightarrow p \vee [p \wedge q]$	Identity Law
	$\leftrightarrow p$	Absorption Law

9) Show that  $(p \rightarrow r) \wedge (q \rightarrow r)$  and  $(p \vee q) \rightarrow r$  are logically equivalent using the laws of logic equivalence and the definition of the conditional statement only. Show each step and state which rule is being used.

**Solution**

$(p \rightarrow r) \wedge (q \rightarrow r)$	
$= (\neg p \vee r) \wedge (\neg q \vee r)$	Definition of Implication
$= ((\neg p \vee r) \wedge \neg q) \vee ((\neg p \vee r) \wedge r)$	Distributive Law
$= ((\neg p \vee r) \wedge \neg q) \vee r$	Absorption Law
$= (\neg p \wedge \neg q) \vee (r \wedge \neg q) \vee r$	Distributive Law
$= \neg(p \vee q) \vee (r \wedge \neg q) \vee r$	DeMorgan's Law
$= \neg(p \vee q) \vee r$	Absorption Law
$= (p \vee q) \rightarrow r$	Definition of Implication

Here is an alternate and quicker solution:

$(p \rightarrow r) \wedge (q \rightarrow r)$	
$= (\neg p \vee r) \wedge (\neg q \vee r)$	Definition of Implication
$= (\neg p \wedge \neg q) \vee r$	Distributive Law
$= \neg(p \vee q) \vee r$	DeMorgan's Law
$= (p \vee q) \rightarrow r$	Definition of Implication

10) Use the Rules of Inference and the Law of Contraposition to validate the conclusion drawn below. (Each of the items above the dotted line is a premise, while the conclusion to draw is below the dotted line.) Show each step and state which rule is being used.

$q \rightarrow (u \wedge t)$   
 $u \rightarrow p$   
 $q$   
 -----  
 $p \wedge t$

**Solution**

Step	Rule
1. $q \rightarrow (u \wedge t)$	Premise
2. $q$	Premise
3. $(u \wedge t)$	Modus Ponens (1, 2)
4. $u$	Simplification (3)
5. $u \rightarrow p$	Premise
6. $p$	Modus ponens (4, 5)
7. $t$	Simplification (4)
$p \wedge t$	Conjunction (6, 7)