

The Protection of Information in Computer Systems

Written by
Jerome H. Saltzer
Michael D. Schroeder

- ◆ Presented by KeeHong Pang



Organization

- ◆ Section I
 - Desired functions
 - Design principles
 - Examples of elementary protection and authentication mechanisms
- ◆ Section II
 - Principles of modern protection architectures
 - The relation between capability systems and access control list systems
 - Protected subsystems and protected objects
- ◆ Section III
 - Review of the state of the art and current research projects



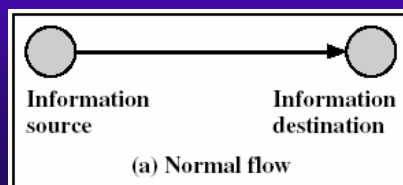
The Beginning


- ◆ Goal
 - Explores the mechanics of protecting computer information from unauthorized use or modification.
- ◆ Motive
 - To control sharing of information among multiple users.
- ◆ This paper concentrates on
 - Protection
 - Authentication



Security violation categories

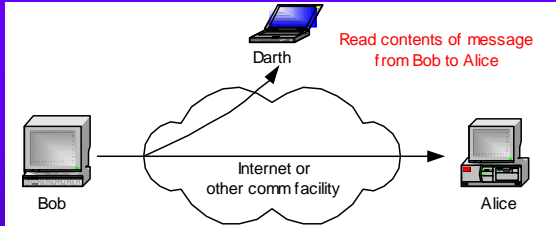
- ◆ Passive attack
 - Release of message contents
 - Traffic analysis
- ◆ Active attack
 - Masquerade
 - Replay
 - Modification of message
 - Denial of service





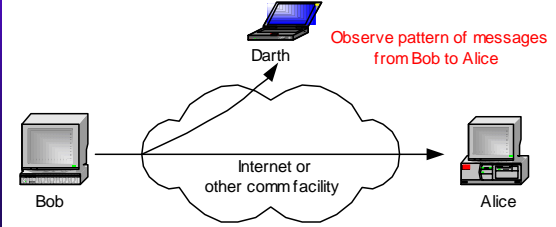
Passive Attacks

- ◆ Release of message contents




Read contents of message from Bob to Alice

- ◆ Traffic analysis

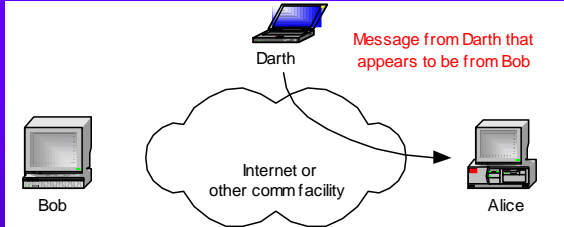


Observe pattern of messages from Bob to Alice



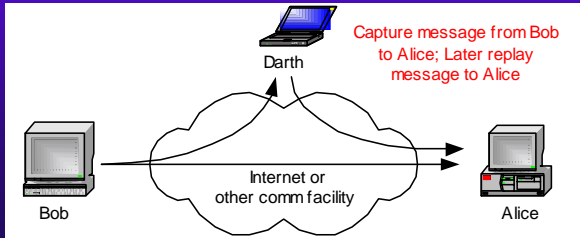
Active Attacks

- ◆ Masquerade



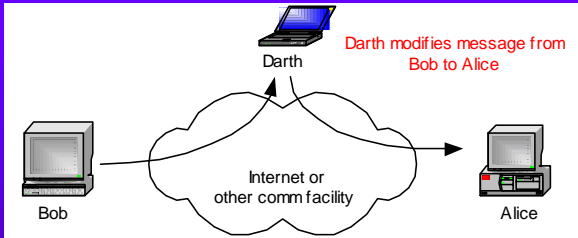
Message from Darth that appears to be from Bob

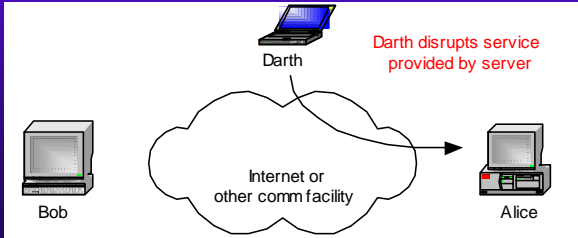
- ◆ Replay



Capture message from Bob to Alice; Later replay message to Alice

Active Attacks

- ◆ Modification of messages
 

The diagram shows Bob on the left and Alice on the right, connected by a cloud labeled 'Internet or other comm facility'. Above the cloud is a laptop labeled 'Darth'. An arrow points from Bob to the cloud, and another from the cloud to Alice. A red arrow points from Darth to the communication path with the text 'Darth modifies message from Bob to Alice'.
- ◆ Denial of service
 

The diagram shows Bob on the left and Alice on the right, connected by a cloud labeled 'Internet or other comm facility'. Above the cloud is a laptop labeled 'Darth'. An arrow points from Bob to the cloud, and another from the cloud to Alice. A red arrow points from Darth to the communication path with the text 'Darth disrupts service provided by server'.

Protection schemes

- ◆ Unprotected systems
 - No provision for protection.
- ◆ All-or-nothing systems
 - Provide isolation of users or total sharing of some info.
- ◆ Controlled sharing
 - Control who may access each data item stored in the system.
- ◆ User-programmed sharing controls
 - Restrict access to a file in a way not provided in the standard.
- ◆ Putting strings on information
 - Maintain control over the user of the information even after releasing.

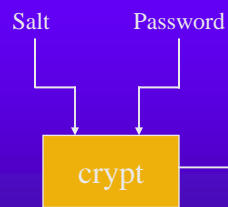


Design Principles

- ◆ Economy of mechanism
- ◆ Fail-safe defaults
- ◆ Complete mediation
- ◆ Open design
- ◆ Separation of privilege
- ◆ Least privilege
- ◆ Least common mechanism
- ◆ User friendly interface

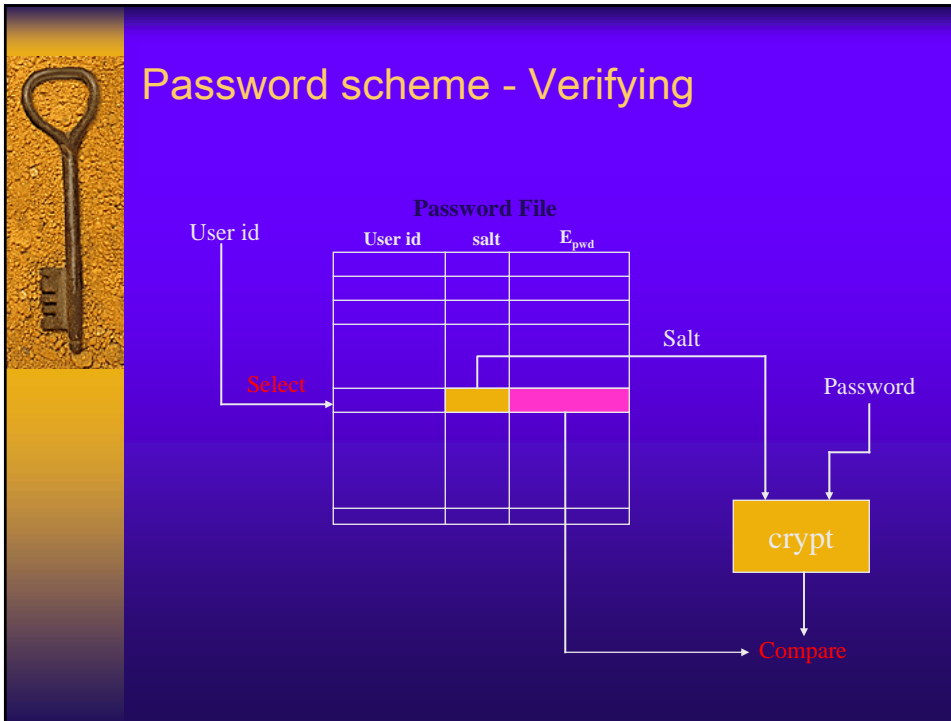


Password scheme - Loading

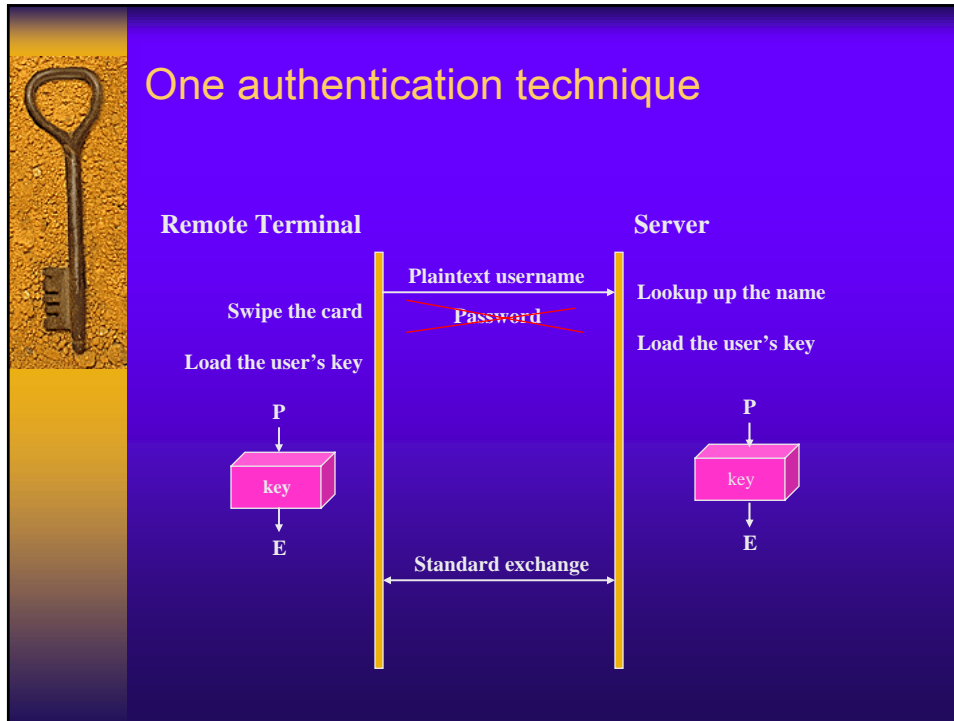


Password File

User id	salt	E _{pwd}
⋮	⋮	⋮



- ## Defects in password systems
- ◆ Choice of password
 - Limit of length and combination
 - Password aging
 - System-generated password
 - ◆ Plaintext transfer
 - Encryption
 - One-time password
 - ◆ One-way authentication
 - Use LUCIFER system



Access Control


◆ Authentication → Authorization

◆ Terminology

- Objects
 - An entity to which access must be controlled.
 - EX) process, file, database, semaphore, printer, memory segment
 - Type: the set of operations
- Subjects
 - An entity whose access to objects must be controlled.
 - EX) process, user
- Protection rules
 - Definition in which subjects can allowed to access objects.
 - Access right (subject, object)

◆ Models

- Access matrix model
- Information flow control model
- Security kernel model



Protection Domains

- ◆ An abstract definition of a set of access rights.
 - Not disjoint
 - Existence in multiple domains

D1

(File1, {Read, Write, Execute})
(File2, {Read})


{Semaphore1, {Up, Down}}

D3

(File3, {Read, Write, Execute})
(DapeDrive1, {Read, Write, Rewind})

D2

(File1, {Read, Write})
(File2, {Read, Write, Execute})
(TapeDrive1, {Read})



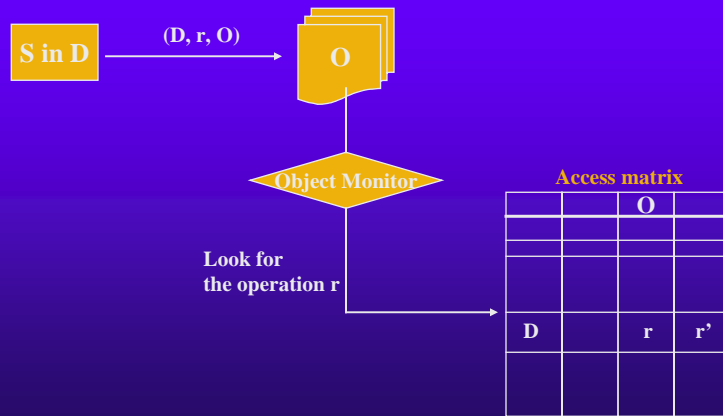
Access Matrix

- ◆ A matrix representing which rights on which objects belong to a particular domain.

Object Domain	F1 (File1)	F2 (File2)	F3 (File3)	S1 (Semaphore1)	T1 (Tape drive1)
D1	Read Write Execute	Read		Up Down	
D2	Read Write	Read Write Execute		Up Down	Read
D3			Read Write Execute		Read Write Rewind

Validation of access


- ◆ Object monitor with each type of object



Domain Switching

- ◆ Guarantee the principle of least privilege.
- ◆ Operation - *switch*


Object \ Domain	F1	F2	F3	S1	T1	D1	D2	D3
D1	Read Write Execute	Read		Up Down			Switch	Switch
D2	Read Write	Read Write Execute		Up Down	Read			Switch
D3			Read Write Execute		Read Write Rewind			



Change to the Protection State (1)

- ◆ *Copy* right
 - Copy an access right from one domain to another
 - Transfer / Copy with propagation not allowed /Copy with propagation allowed

Object Domain	F1	F2	F3
D1	Read* Write* Execute	Read	
D2	Read* Write	Read* Write Execute*	
D3			Read Write Execute



Change to the Protection State (2)

- ◆ *Owner* right
 - Adding/deleting of rights to column entries

Object Domain	F1	F2	F3
D1	Read* Write* Execute Owner	Read	
D2	Read* Write	Read* Write Execute* Owner	
D3			Read Write Execute Owner



Change to the Protection State (3)

◆ *Control* right

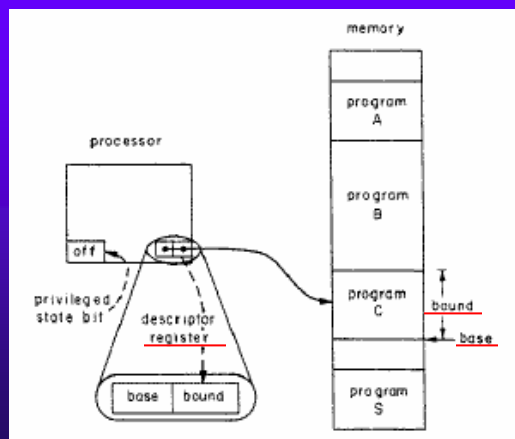
- Only applicable to domain objects
- A process can change the entries in a row

Object Domain	F1	F2	F3	S1	T1	D1	D2	D3
D1	Read Write Execute	Read		Up Down			Switch Control	Switch Control
D2	Read Write	Read Write Execute		Up Down	Read			Switch Control
D3			Read Write Execute		Read Write Rewind			



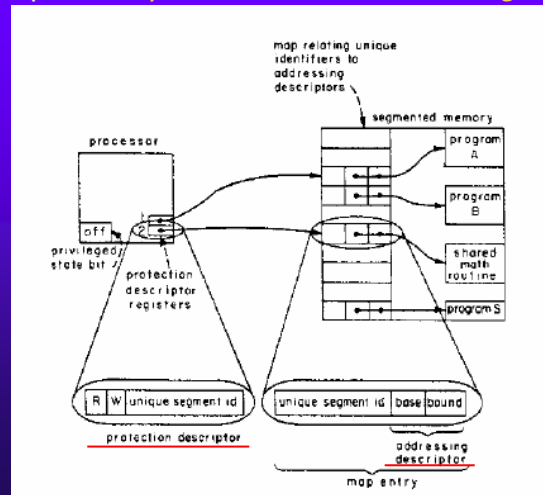
Descriptor

- ◆ The value of the Descriptor Register to protect information.



Separation of Addressing and Protection

- ◆ All memory accesses were divided into two levels of descriptors → *protection* and *addressing*



Approaches

- ◆ Concept
- ◆ Validation
- ◆ Sharing
- ◆ Revocation

Capability

- ◆ Decompose access matrix by rows
- ◆ Maintain (object, rights) pairs – *capability*

capability

Object id.	Rights info.
program B	
database Y	
shared math routine	
database X	
program A	

segment name	capability for segment
X	→ database X
A	→ program A
Math	→ shared math routine

Catalog for UserA

Simple Capability System

processor

protection descriptor registers

1
2
3
4

segmented memory

program A

shared math routine

C2

C1

private data base X

private data base C5

private data base Z

C4

C3

catalog for Doe

catalog for Smith

Access Control Lists

- ◆ Decompose access matrix by columns
- ◆ Maintains (domain, rights) pairs for each object

access controller


base	bound
D1	read write
⋮	⋮
D2	read
⋮	⋮

addressing descriptor for this segment

principal identifier
permission

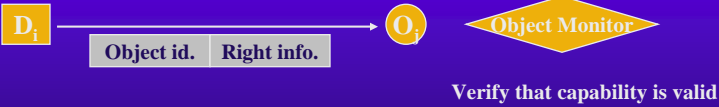
Access Control List System


The diagram illustrates the Access Control List System architecture. On the left, a **processor** contains a **principal identifier register** and **pointer registers**. The **principal identifier register** outputs a **unique identifier**. The **pointer registers** output a **unique identifier**, **base**, and **bound**. These three components are combined into a single structure. This structure is used to access **segmented memory**, which contains **segment X**. An **access controller AC_i for segment X** is associated with segment X. A **map relating unique identifiers to addressing descriptors** is also shown, which maps the unique identifier to the addressing descriptor (base and bound).



Access Validation - Capability

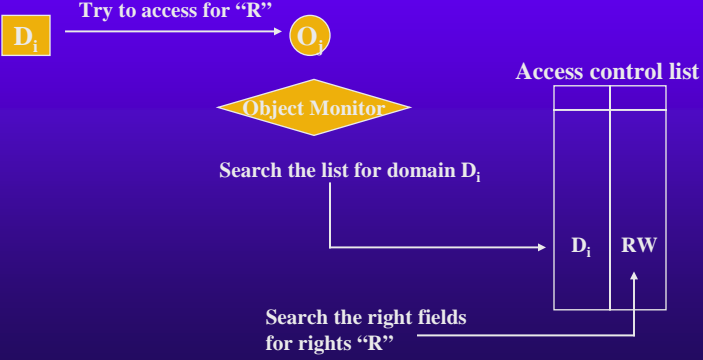
- ◆ A capability = An unforgeable ticket
- ◆ No need to search a list → Only verify that capability is valid
- ◆ Access the object without any further check.





Access Validation – ACLs (1)

- ◆ The access list for object O is first searched for D.
- ◆ The rights field of this element is searched.
- ◆ Check the access list on every access.
 - More security, but not efficiency



D _i	RW

Access Validation – ACLs (2)

- ◆ Use of “shadow” capability registers
 - Invisible to the virtual processor.
 - The shadow register is loaded with directly access to the segment.
 - EX) file open and close in UNIX system

The diagram illustrates the hardware mechanism for access validation. On the left, a 'processor' contains 'pointer registers' and 'principal identifiers'. On the right, 'segmented memory' contains 'access controller' and 'segment'. A dashed arrow points from 'pointer registers' to 'shadow registers' (located between the processor and memory). Solid arrows show the flow of information: from 'principal identifiers' to 'access controller', from 'access controller' to 'segment', and from 'shadow registers' to 'access controller'.

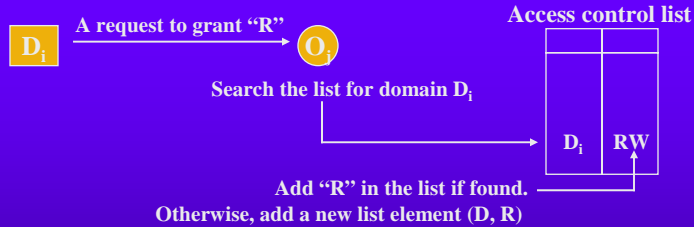
- ◆ Limit the number of entries on each access control list.

Dynamic Sharing - Capability

- ◆ One or more object managers for each type of object.

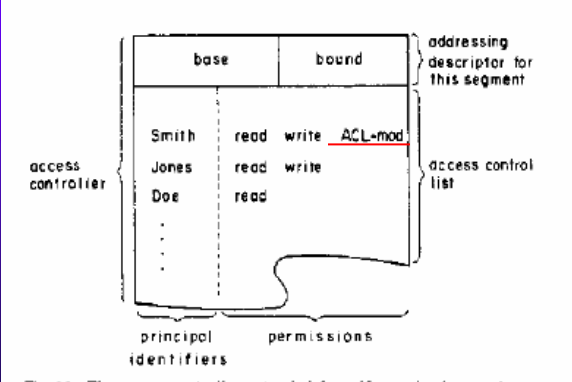
The diagram shows two scenarios of dynamic sharing. In scenario 1, 'Object Manager(a)' and 'Object Manager(j)' are shown. 'Object Manager(a)' sends 'A request to create an object' to 'Object Manager(j)'. 'Object Manager(j)' then sends 'Capability' back to 'D_i' and also performs the action 'Generate a capability with all rights'. In scenario 2, 'Object Manager(j)' sends 'A request to perform some operation' to 'D_i'. 'Object Manager(j)' then sends 'New capability' to both 'D_i' and 'D_k', and performs the action 'Generate a new capability'.

Dynamic Sharing – ACLs (1)

- ◆ To grant access right r for object O to domain D

- ◆ To pass access right r from a domain $D1$ to another domain $D2$
 - Check if $D1$ possesses either *owner* right or *copy* right for access right r .

Dynamic Sharing – ACLs (2)

- ◆ Self control
 - Permission to modify the access control list.
 - Too absolute – no provision for another way to control.



Dynamic Sharing – ACLs (3)

- ◆ Hierarchical control
 - The creator specifies some previously existing access controller whenever a new object is created.
 - Too powerful authority in higher level.

Revocation (1) - Capability

- ◆ The capabilities for an object are stored in several capability lists.
 - Difficulty to determine which subjects have what rights for the object.
- ◆ Method for implementing revocation
 - Back Pointers
 - Keep track of all the capabilities for an object.
 - Change/Delete the capabilities selectively.
 - Maintain a list of pointers with the object.

Revocation (2) - Capability

- Indirection
 - Each capability points to an indirect object.
 - Non-selective revocation.

- Use of keys
 - *Key* - A field that contains a unique bit pattern in each capability.
 - Each object has a *master key*.

Revocation - ACLs

- ◆ Simply delete access right *r* from the rights set of domain *D* in the access list for *O*.

D _i	RW



Conclusions

- ◆ Understanding of password and authentication mechanism.
- ◆ Comparison between capability systems and access control list systems.
- ◆ Need to develop various mechanisms for supporting security services.



References

- ◆ The Protection of Information in Computer Systems, J.H.Saltzer and M.D.Schroeder.
- ◆ Distributed Operating Systems.
- ◆ Operating System Concepts 6th ed..
- ◆ Cryptography and Network Security.
- ◆ Network Security Essentials