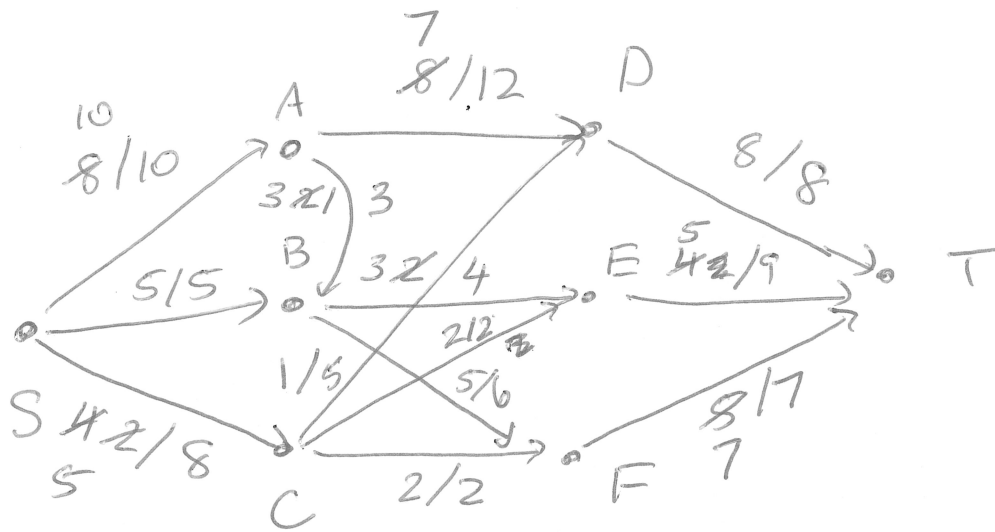


# Network Flow



Directed  
Weighted  
S - source  
only has outgoing  
edges  
T - sink  
only has incoming  
edges

edge weights = capacities (max amt stuff flow through per unit of time)

Max Flow network = max units stuff/unit of time that can go from S to T.

How to find?

MOST ALGS LOOK LIKE THIS

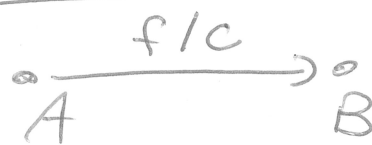
While (we can send more flow)

This is an augmenting path

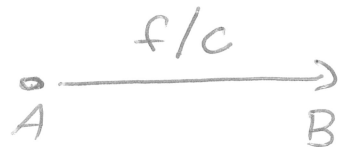
⇒ Send max flow down that path.

## AUGMENTING PATHS

forward edge



if  $f < c$   
send  $c - f$



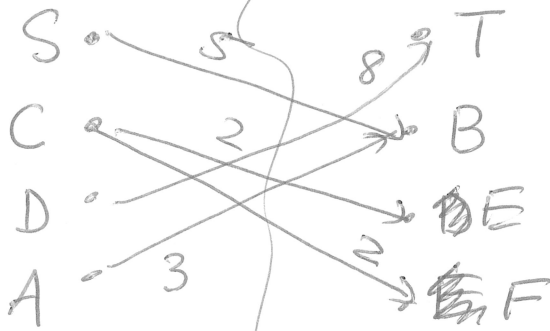
if  $f > 0$   
←  $f$  backwards

- SADT → 8
- SBIT → 5
- SABET → 2
- SCFT → 2
- SCET → 2
- SCDABET → 1

20

SEARCH S, C, D, A, STOPS

At most 20 units can flow from  $\{S, C, D, A\} \rightarrow \{B, E, F, T\}$



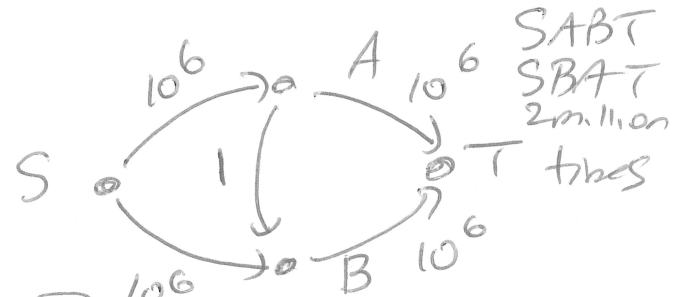
max flow = min cut

~~A cut~~

A cut is a partition of the vertices w.r.t in one set, T in the other. The value of a cut is the sum of capacities of all edges from S's set to T's set.

### 3 ALGS

- 1) ~~Folk~~  
Ford-Fulkerson  
- DFS



SABT  
SBAT  
2 million times

} not poly time

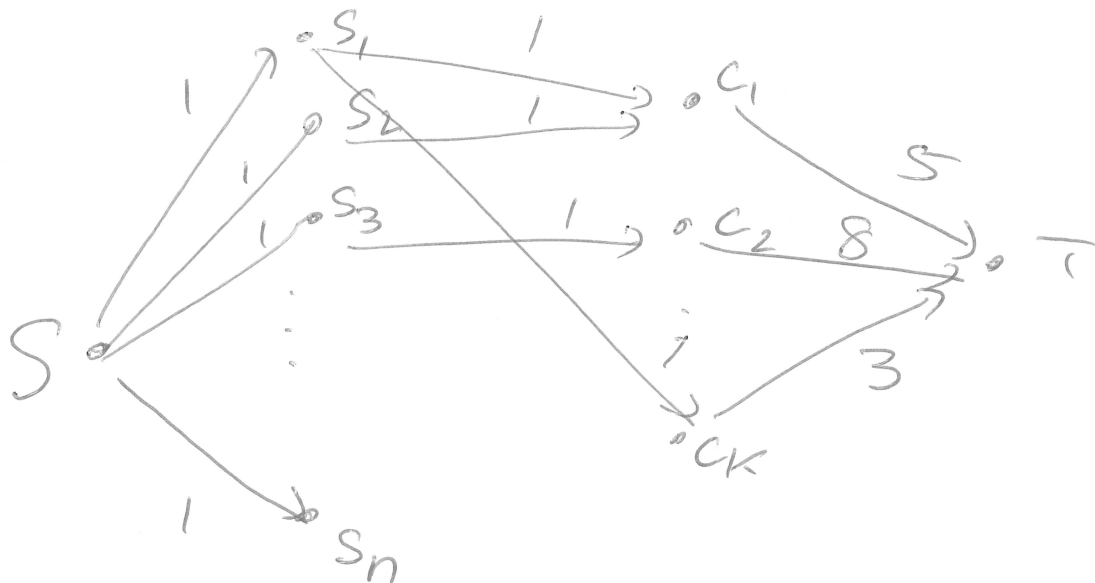
- 2) Edmonds-Karp  
- BFS } finish in poly times

- 3) Dinic's (Dinitz)

while aug path [ 1 BFS  $\Rightarrow$  DFS stops at length of BFS

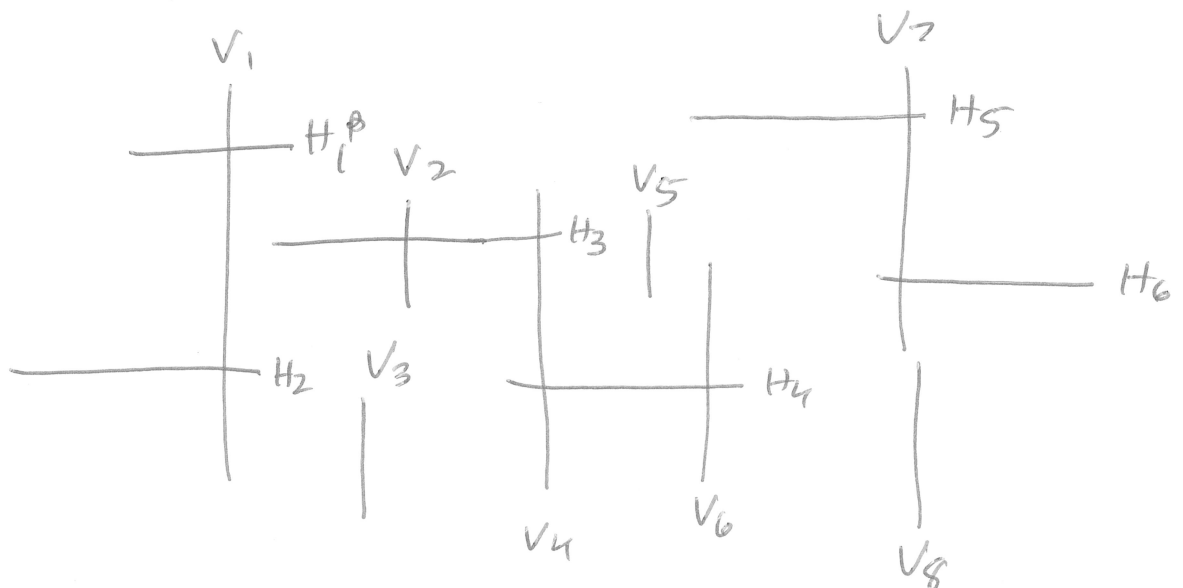
① Few Applications

② Walk Through Posted Code

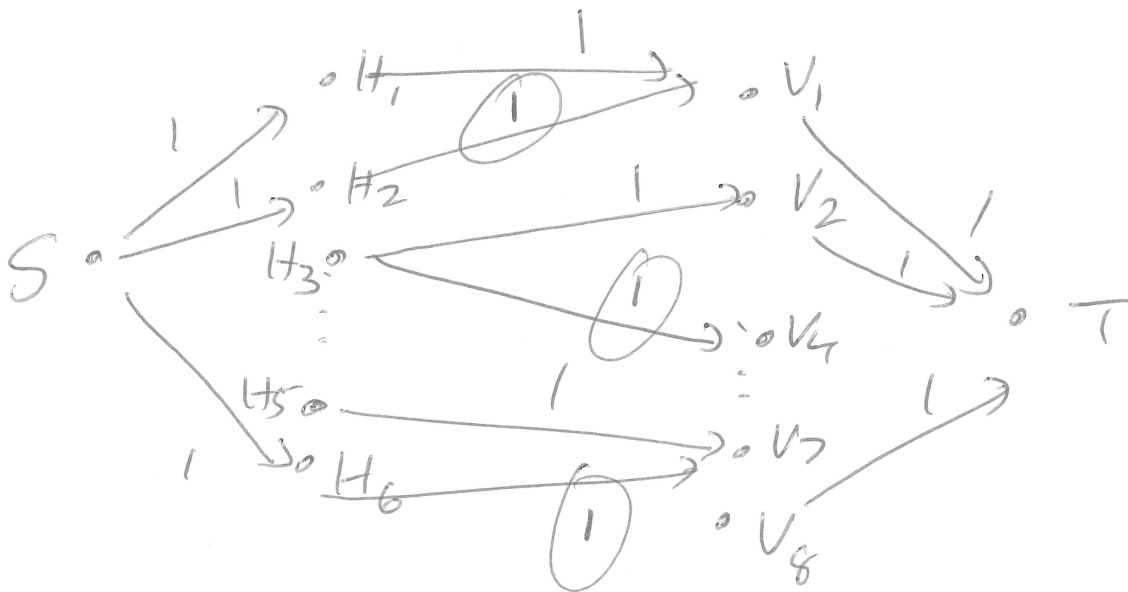


max flow = max # students who could get internships

Cow Steeplechase

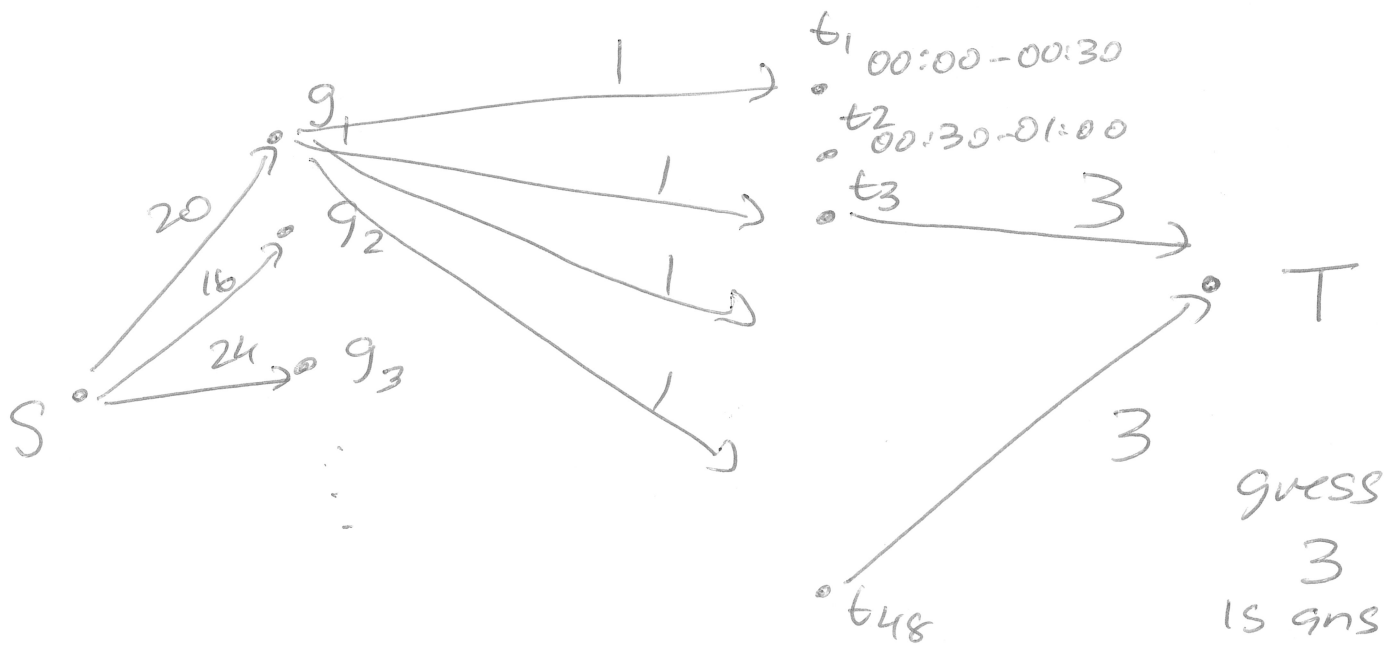


What is the fewest # of line segments I need to remove so no 2 line segment intersect?



max flow represents the # of ordered pairs where we are forced to remove one line segment.

### Museum Guards



run max flow  $\Rightarrow$  if ans =  $3 \times 48$   
 then it's possible to use 3 guards

Binary Search the correct ans if  
 flow works ans  $\geq$  tried, else ans  $<$  tried