

units

Starting node called source

Ending node call sink

Definition of a flow

weighted graph (most of the time directed).
typically edge weights are referred to capacities.

A flow from s to t is as follow

- 1) Each edge is assigned a non-negative value, flow of edge, that each flow is no greater than the edge capacity

$$\forall (u, v) \in E \quad (f_{uv} \leq w_{uv})$$

↑ flow ↑ capacity

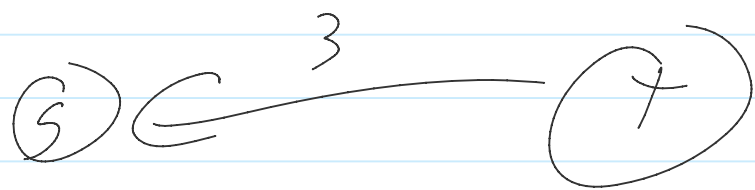
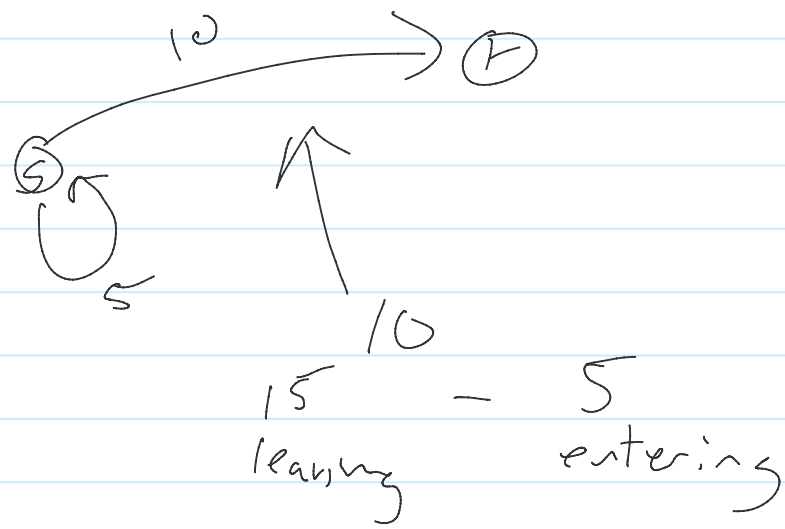
- 2) Aside from source (s) and sink (t) each node need to conserve the flow. The in flow is out flow

$$\forall (v \in V \setminus \{s, t\}) \quad \left(\sum_{(u,v) \in E} f_{uv} = \sum_{(v,u) \in E} f_{vu} \right)$$

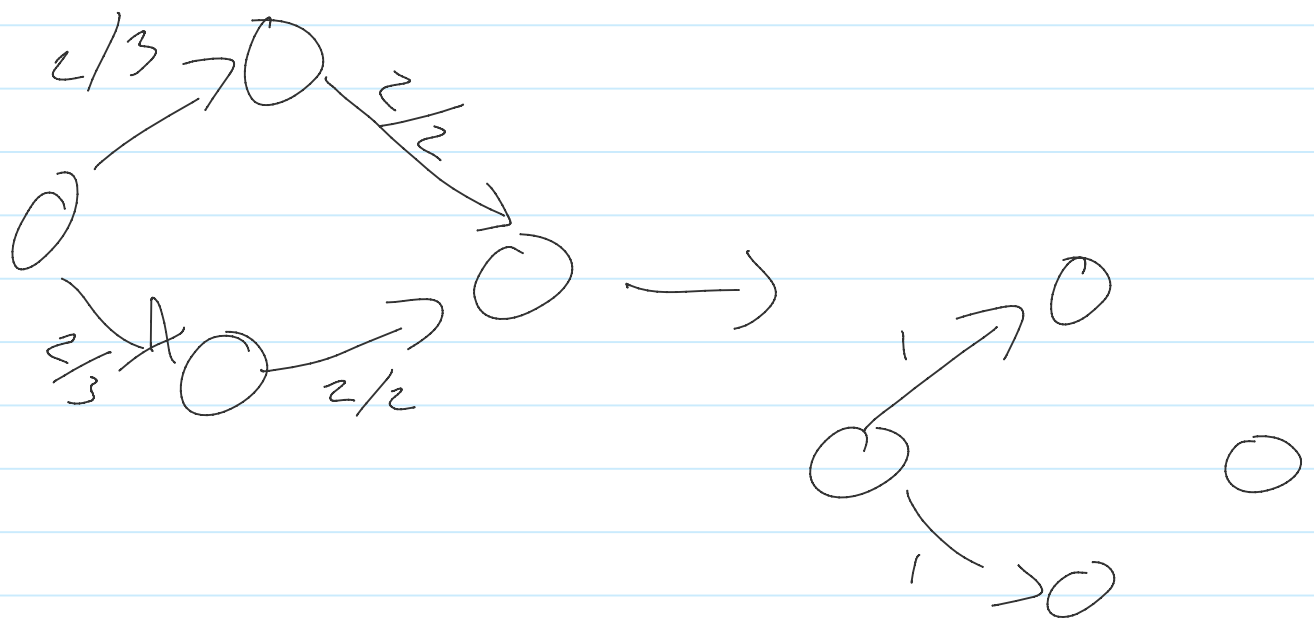
↑ set difference ↑ out flow ↑ in flow

The flow of the graph is

out flow of the source node minus the in flow of the source node

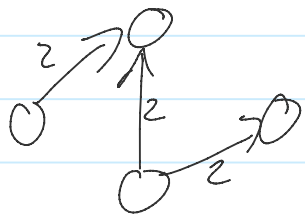
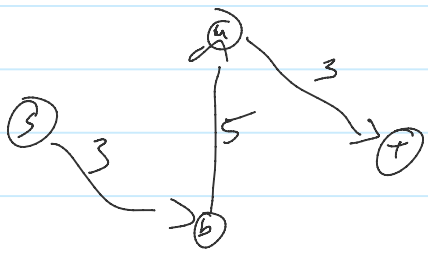
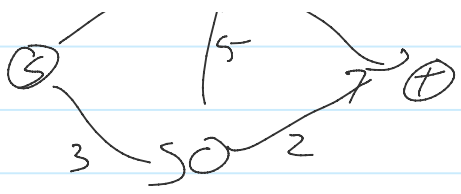


max flow for this graph is 0
 but you could have -3 flow



Find a path send flow through the path (the flow of which is the min capacity along path) update graph and remove full edges

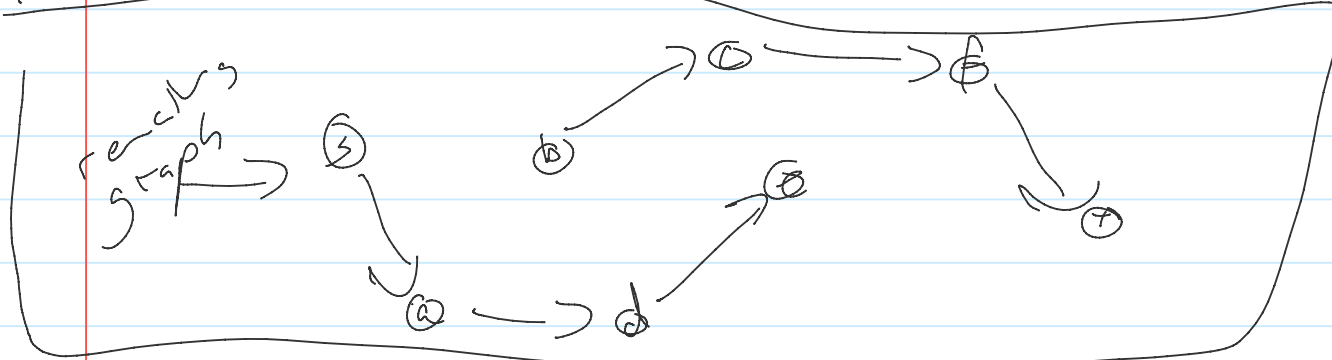
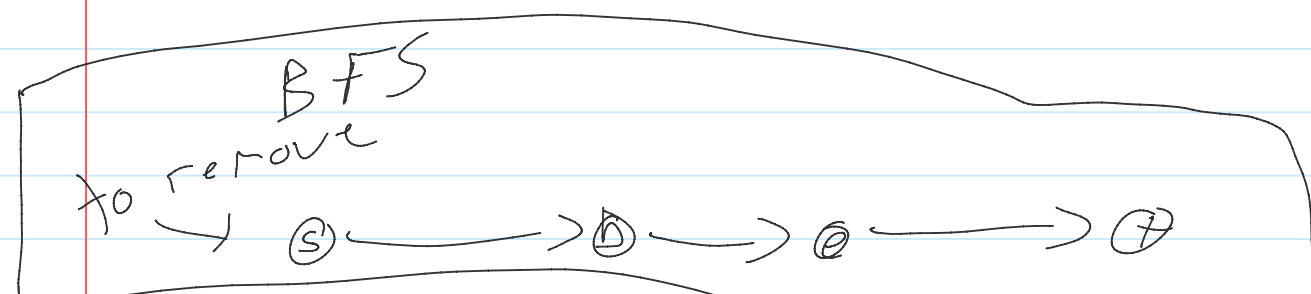
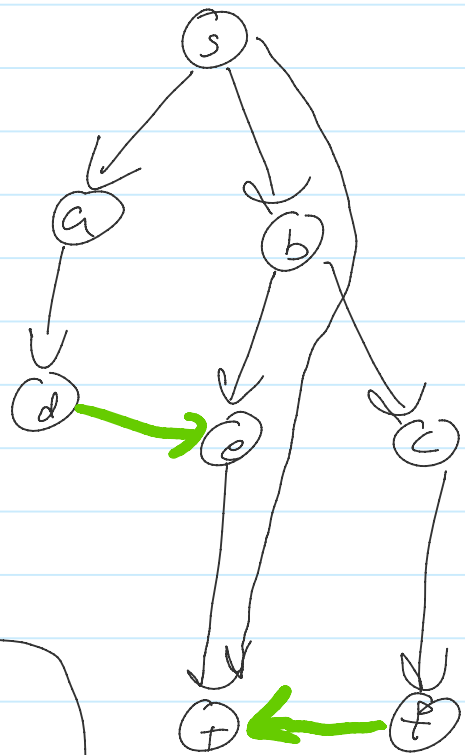
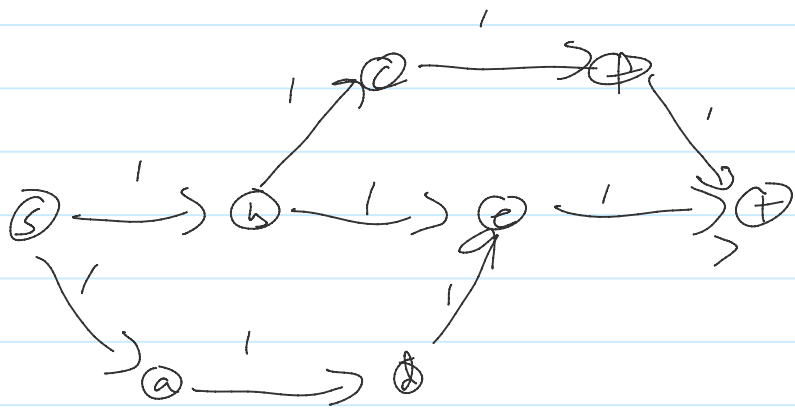




found flow is 3

actual max flow is 5

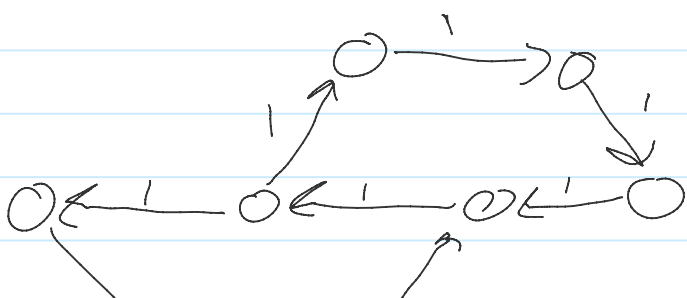
try a BFS; this did not work
Other tricky case



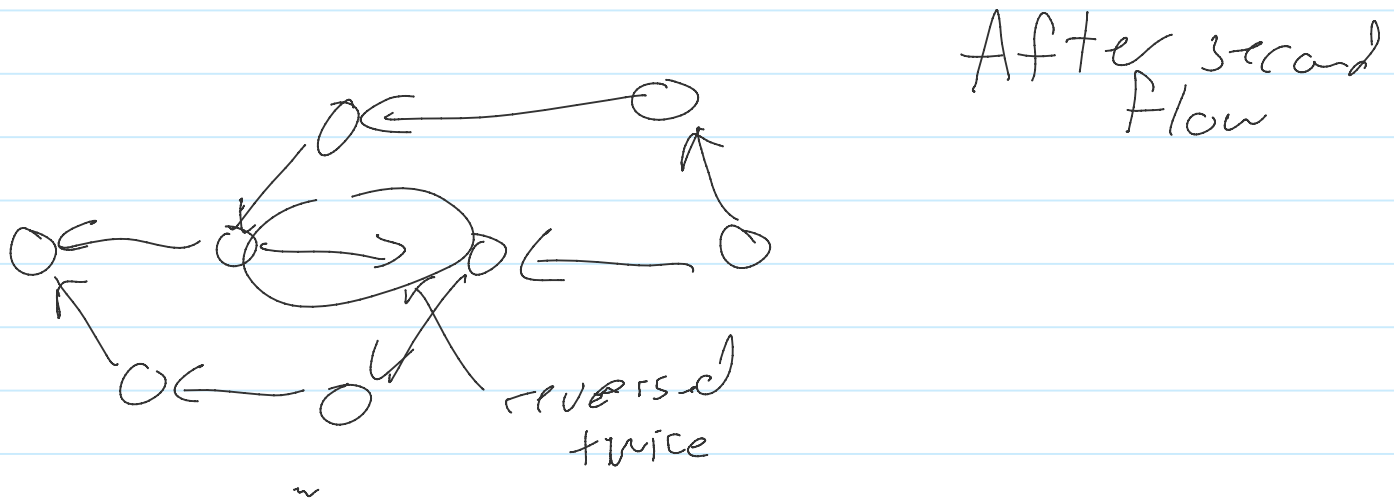
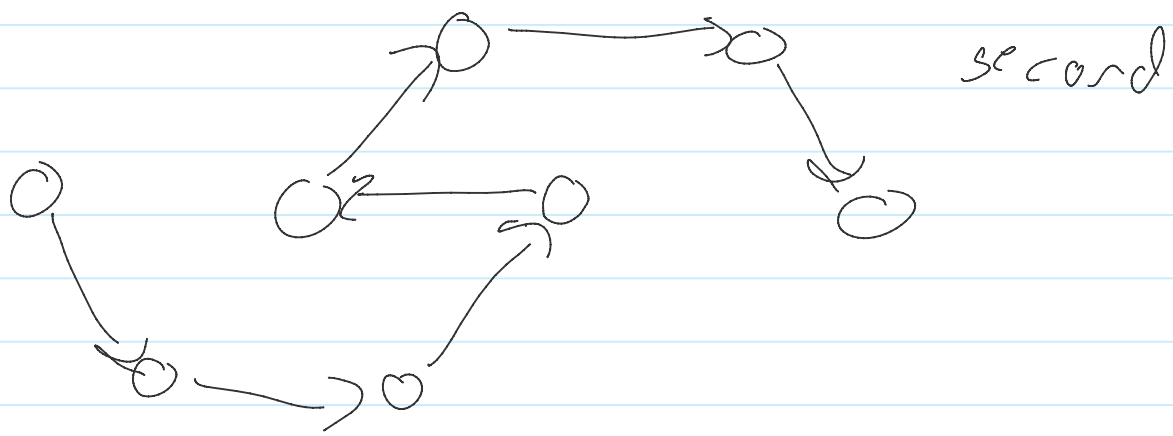
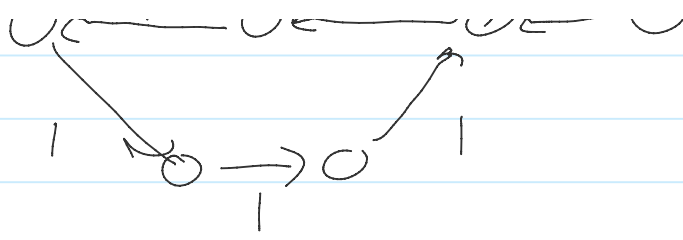
found flow is 1

actual max flow is 2

Residual Graph



After first flow



Run time of Algorithm n The Algorithm is Ford-Fulkerson Algorithm

$$O(F \times \text{time to find path})$$

$$\uparrow \text{DFS} = O(|V| + |E|)$$

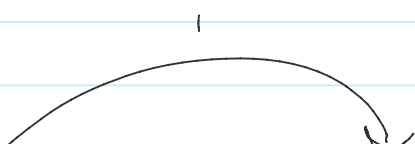
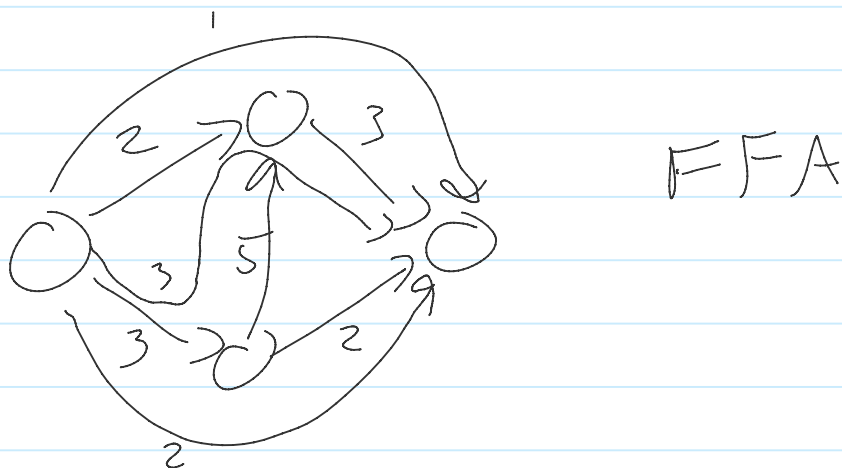
$$\text{DFS} = O(|V| + |E|)$$

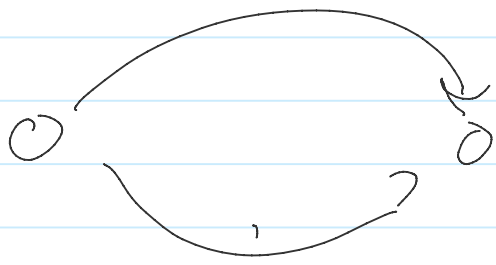
$$O(F(|V| + |E|))$$

Dinitz's (define Dinitz)

Dinic's

Better algorithm $O(|V|^2 |E|)$





Correctness of FFA

Suppose FFA did not find an optimum solution. Let F' be the flow graph of the optimum ~~graph~~ flow.

Sum of flow for FFA (F) nodes is zero
and sum of flow for F' nodes is zero

Consider $F' - F$

We know that a positive path exists in

$F' - F$ we should have found that path

in FFA we reached a contradiction.