1/24 Minimum Spanning Tree

Wednesday, January 24, 2018 5:57 PM

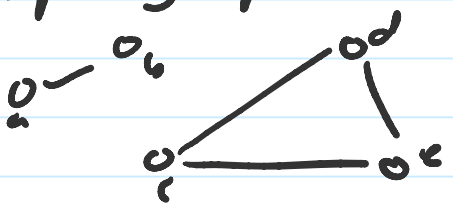
MST or Minimum Spanning Tree

Graph 2 sets
Vertex Set V or V(G)
Edge set E or E(G)

Vertex (Points or Nodes)
Edge (connection)

e = (a, b) where
a, b ∈ V

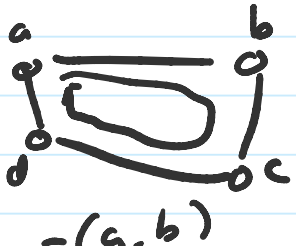
Example Graph



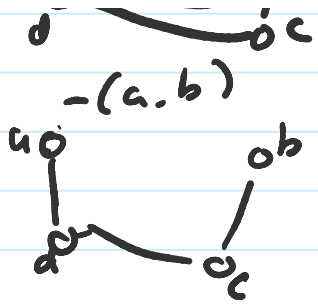
V = {a, b, c, d, e}

Tree a connect graph with no self loops or cycles

example cycle



e(a, b)



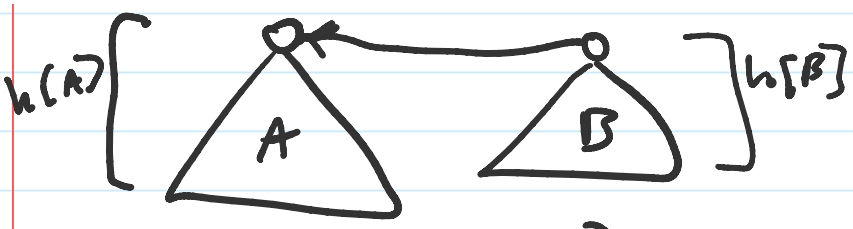
w(a, d)
w(d, c)
w(b, c)

Disjoint set revisited

V = {a}, {e}, ...
{b}, {d}, ...,
for our purposes
= {1}, {2}, {3}, ... {n}

find(int n)
returns a representative for
the set containing n;

union(int a, int b)
int rep a =
int rep b
adjust the disjoint set
so repa will represent
set containing b or vice
versa



if (h(A] > h(B))
then
adding B to A
won't change the height
of A
otherwise
The height of A going
into B will be at most
h[B] + 1

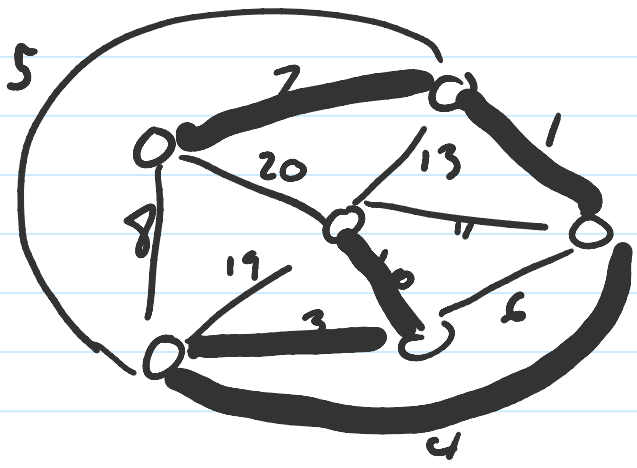
Suppose we have
some connected graph G

if we add an edge to
G we induce a cycle

add edge (u, v)

Sort all the edges by weight

By greedily selecting the least weight edge and adding it to our graph if the two vertexes are disconnected we can build a tree.



A tree with n Nodes contains n-1 Edges

A forest is a graph (potentially disconnected) with no cycles.

Kruskal's works by greedily adding edges until all edges have been checked

Kruskal's step by step

1. Initialize chosen Edges to $\emptyset$ (Empty Set)

Kruskal's step by step

1. Initialize chosen Edges to $\emptyset$ (Empty Set)
2. Initialize a Disjoint Set to track connected nodes
3. Loop through edges least to greatest
   - 3a. Check if current edge connects two separate components.
   - 3b. Add edge if components are disconnected
   - 3c. Update Disjoint Set if necessary
4. Return chosen Edge Set.