

0-1 Knapsack

- ① 0-1 Knapsack
- ② 2nd Change Prob

Given: list of items, each has weight, value

Item	Weight	Value
1	4	10
2	2	3
3	5	11
4	7	13

4/5/18 ①

Knapsack Weight Limit = 9 units

Problem what's the max value of any subset of items that weighs less than or equal to the weight limit?

Alg #1

Divide + Conquer

- try all subsets w/ item 1
- try all subsets w/o item 1

runtime $O(2^n)$, $n = \#$ items

Alg #2 - Dynamic Programming Alg

max Weight = W

0	0	0	0	...	0
0	1	2			W

$dp[i]$ = max value of any subset of items that weighs i of fewer units.

$\begin{matrix} & W & V \\ I_1 & \rightarrow & 4, 10 \\ I_2 & & 2, 3 \\ I_3 & & 5, 11 \\ I_4 & & 7, 13 \end{matrix}$

0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0

4/5/18 ②

Can I_1 help us?
 Currently I can hold 7 lbs.
 and my best value is 0.

If I took I_1 , I only 3 lbs left
 \hookrightarrow new value $dp[3] + 10 = 10$

$j = \text{item \#}$

$$dp[i] = \max(dp[i], dp[i - \text{weights}[j]] + \text{values}[j])$$

We run inner loop backwards to prevent stealing > 1 copy of an item!

Don't take item j .

Do take item j .

for ($j = 0; j < n; j++$) {

if forwards, allow > 1 copy

for ($i = \text{max}; i \geq \text{weights}[j]; i--$)

$$dp[i] = \text{Math.max}(dp[i], dp[i - \text{weights}[j]] + \text{values}[j]);$$

}

$\begin{matrix} W & V \\ I_1 & 4, 10 \\ I_2 & 2, 3 \\ I_3 & 5, 11 \\ I_4 & 7, 13 \end{matrix}$

	0	1	2	3	4	5	6	7	8	9
	0	0	0	0	0	0	0	0	0	0
I_1	0	0	0	0	10	10	10	10	10	10
I_2	0	0	3	3	10	10	13	13	13	13
I_3	0	0	3	3	10	11	13	14	14	21
I_4	0	0	3	3	10	11	13	14	14	21

4/5/18 ③

<u>Items</u>	<u>W</u>	<u>Value</u>
CC granola	4	6
Mandarin	2	4
Cheese Stick	3	5
Chocolate	1	3
Bananas	6	9
CCPB granola	4	7

$$W = 10$$

	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>	<u>7</u>	<u>8</u>	<u>9</u>	<u>10</u>
	0	0	0	0	0	0	0	0	0	0
CCg	0	0	0	6	6	6	6	6	6	6
mand	0	4	4	6	6	10	10	10	10	10
Cheese	0	4	5	6	9	10	11	11	15	15
Choc	3	4	7	8	9	12	13	14	15	18
Bananas	3	4	7	8	9	12	13	14	16	18
CCPB granola	3	4	7	8	10	12	14	15	16	19

Change Problem

Given : Denominations

1, 2, 4, 7, 12

What is the fewest # of coins necessary to make change for n cents?

$dp[i]$ = fewest # of coins to make change for i cents

but $dp[i] = i$ (use pennies)

rest of the coins

$$dp[i] = \min(dp[i],$$

$$dp[i - denom[j]] + 1)$$

$$dp[74] = \min(dp[74], dp[74 - 25] + 1)$$

↑ fewest coins to make change for $i - denom[j]$ cents.
↑ new coin j