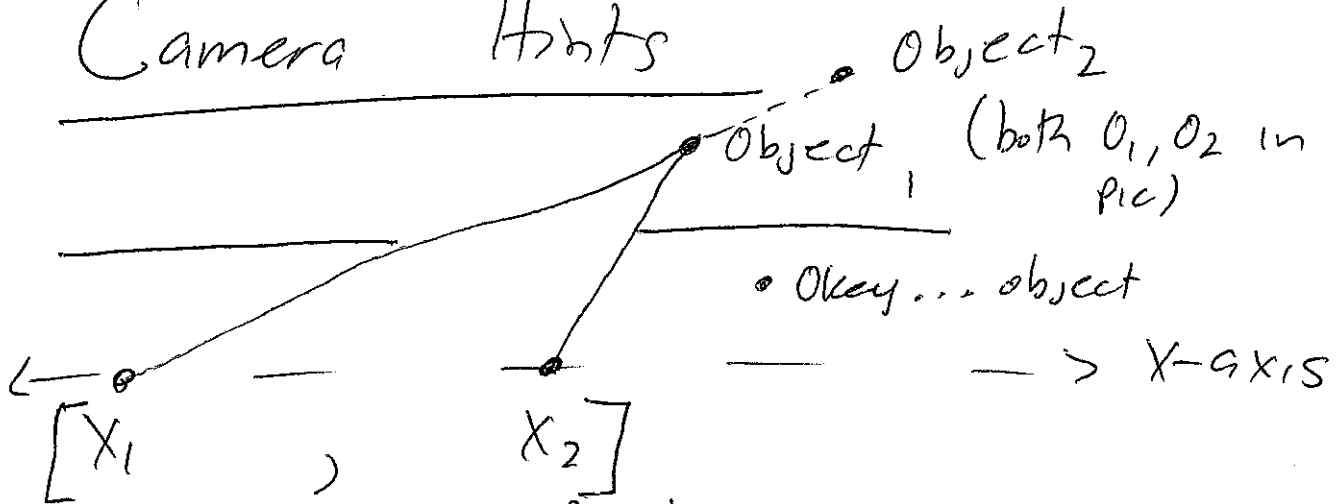


3/1/18 ①

Camera Hints



Client #1 $[-15, 12]$

$\frac{Pic = 4}{\checkmark}$

Scheduling problem

Client #2 $[-10, 8]$

\checkmark

Client #3 $[3, 16]$

\checkmark

To compare equality btw double a, b

$$\text{Math.abs}(a-b) <= \underline{\underline{1e-9}}$$

BFS Applications

3/1/18 (2)

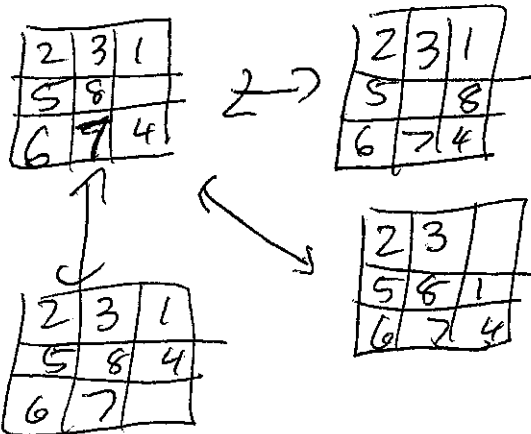
Typical use: find shortest distance (min # edges to traverse) from a starting vertex to all others in an undirected unweighted graph.

Two Different Application

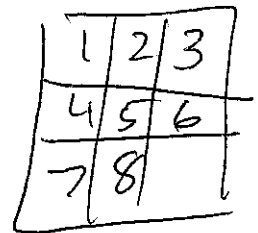
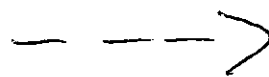
① Eight Puzzle

② Rush Hour

Eight Puzzle



GOAL



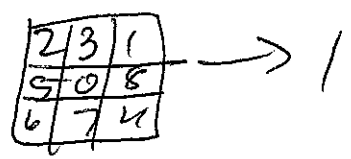
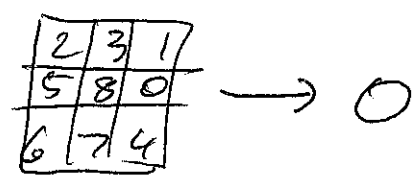
What is the fewest # of moves necessary to solve a given puzzle?

(1) Each board position is a vertex

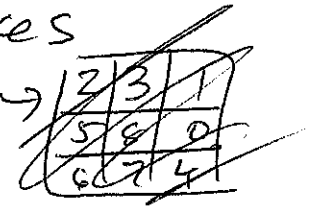
(2) We connect 2 board positions (vertices) if we can reach one from the other in a single move.

DON'T NEED TO STORE THE WHOLE GRAPH!!!

HashMap < BOARD, Integer > distances



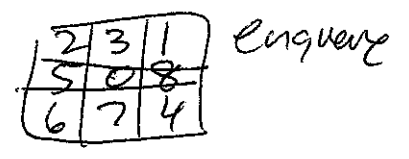
Queue < BOARD > q



BOARD POSITION

9! (300,000 / 400,000)

edges ≤ 2 × 9!



Two ideas to store board for HM, Q:

- (1) String "231580674" } either is fine
- (2) Integer 231580674 }

2 TYPES OF STORAGE

- (1) 2D array → good making moves
- (2) String/Integer → good HM, Q

We need 2 functions

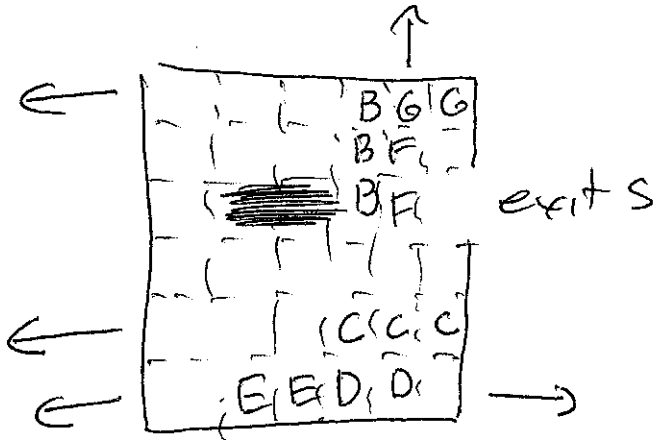
- (1) 2D array → String/Integer compress/encrypt
- (2) String/Integer → 2D array expand/decrypt

Solve 1 puzzle ~ 1,000,000 moves. What about solving 10,000 puzzles? 10^4 × 10^6 = 10^10 BAD.

TRICK: RUN ONE BFS FROM END BOARD TO ALL OTHERS SINCE MOVES ARE SYMMETRIC!

3/1/18 (4)

Rush Hour



1 MOVE -

SLIDING A CAR
AS MANY SQUARES
AS YOU CAN IN A
SINGLE DIRECTION

= A ↓

CALC FEWEST # MOVES
TO FREE YOUR CAR

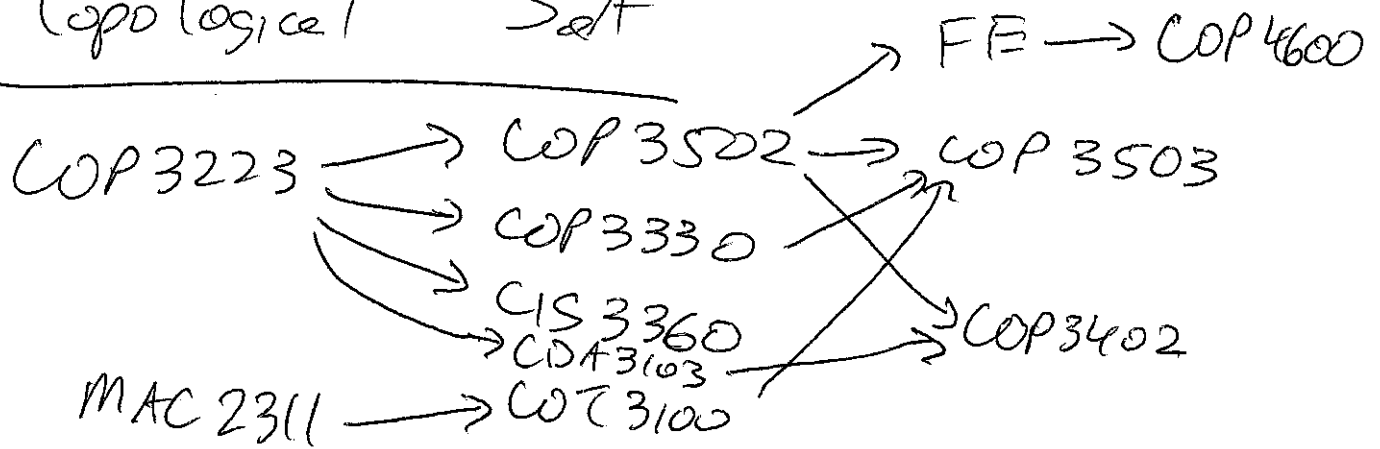
- (1) each board position is a vertex
- (2) we connect 2 board positions if we can obtain one from the other in a single move.

A = left 1, B = down 1, C = left 1, 2, 3

D = right 1, E = left 1, F = down 1

G = no moves

Topological Sort



Provide a valid ordering of classes

Iterative Algorithm

(1) Calculate the in degree of each vertex

COP 3223	MAC 2311	COP 3502	3330	3360	3103	3100	FE	3503	3402	4600
0	0	1	1	1	1	1	1	3	2	1

(2) Create a queue of vertices. Add to that queue all vertices of degree 0.

Q → 3223 → 2311

(3) Create an array to store result

(4) for ($i = 0; i < n; i++$) // loop through vertices ^{num} times

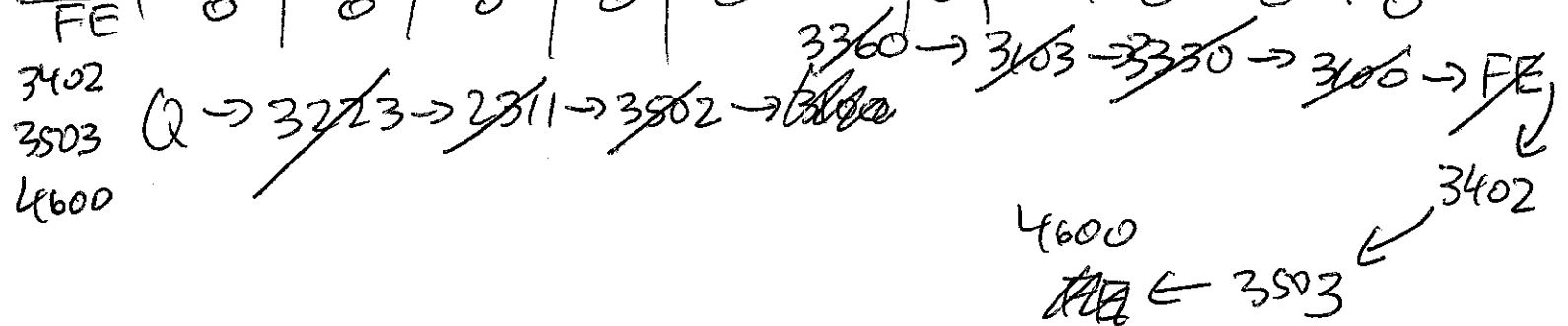
(a) if the queue empty, return false/null

(b) $res[i] = queue.poll()$ (Add the next class from queue to your schedule)

(c) Update all indegrees of edges leaving $res[i]$. If any of $3/1/18$ (6) these hit 0, add to the queue.

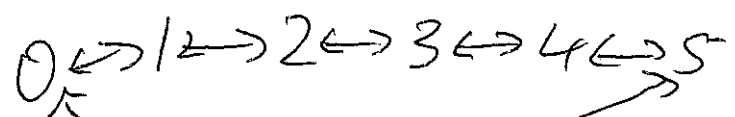
(5) If you get here return res .

	COP3223	2311	3502	3376	3360	3103	3100	FE	3503	3402	4600
Sched	0	0	1	1	1	1	1	1	3	2	1
3223	0	0	0	0	0	0	1	1	3	2	1
2311	0	0	0	0	0	0	0	1	3	2	1
3502	0	0	0	0	0	0	0	0	2	1	1
3360	0	0	0	0	0	0	0	0	2	1	1
3103	0	0	0	0	0	0	0	0	2	0	1
3330	0	0	0	0	0	0	0	0	1	0	1
3100	0	0	0	0	0	0	0	0	0	0	1
FE	0	0	0	0	0	0	0	0	0	0	0



3/1/18 (7)

DFS doesn't do shortest distance



DFS would mark 0 to 5 distance of 5

NOTES