



UCF

**MORE RECURSION:**

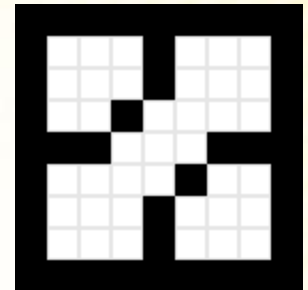
**FLOOD FILL &**

**EXPONENTIATION**

COP 3502

# Recursive Flood Fill Algorithm

- A Flood Fill is a name given to the following basic idea:
  - In a space (typically 2-D, or 3-D) with an initial starting square, fill in all the adjacent squares with some value or item.
    - Until some boundary is hit.
    - For example, the paint bucket in MS Paint is an example of flood fill.

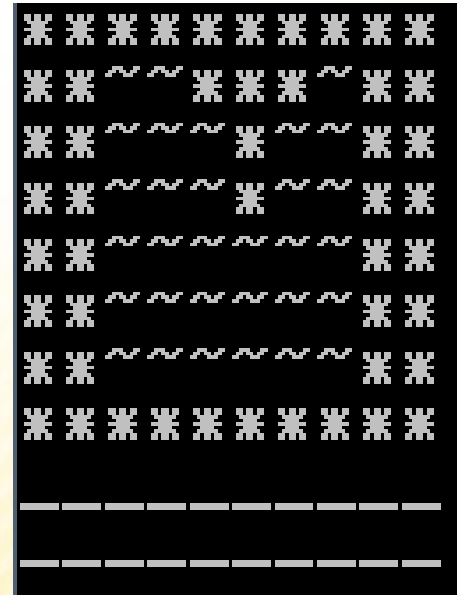
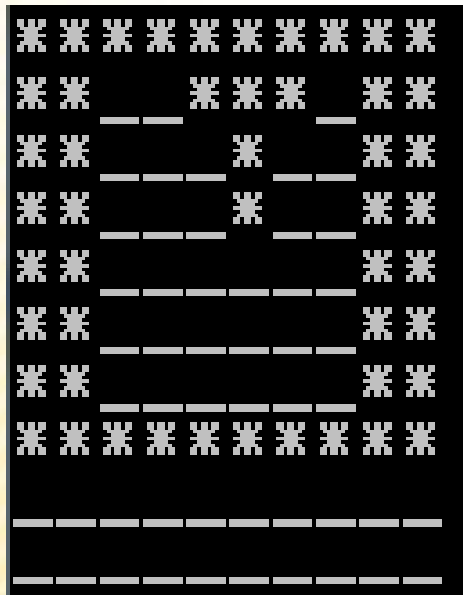


Example of a Recursive Flood Fill with 4 directions



# Recursive Flood Fill Algorithm

- Imagine you want to fill in a “lake” with the ~ character.
  - We’d like to write a function that takes in one spot in the lake (the coordinates to that spot in the grid)
  - In the example, you can see we don’t want to just replace all “\_” with “~”, because we just want to fill the contiguous area.



# Recursive Flood Fill Algorithm

- Depending on how the floodfill should occur
  - Do we just fill in each square above, below, left, and right
  - OR do we ALSO fill in the diagonals
- The basic idea behind a recursive function, is shown in pseudocode:

```
Void FloodFill(char grid[][SIZE], int x, int y) {  
    grid[x][y] = FILL_CHARACTER;  
  
    for (each adjacent location i,j to x,y) {  
        if (i,j is inbounds and not filled)  
            FloodFill(grid, i, j);  
    }  
}
```



# Recursive Flood Fill Algorithm

- When we actually write the code,
  - We may either choose a loop to go through the adjacent locations, or simply spell them out.
  - If there are 8 locations (using the diagonal) a loop is better.
  - If there are 4 or fewer (North, South, East, West)
    - It might make more sense to write each recursive call separately.

```
Void FloodFill(char grid[][SIZE], int x, int y) {  
    grid[x][y] = FILL_CHARACTER;  
  
    for (each adjacent location i,j to x,y) {  
        if (i,j is inbounds and not filled)  
            FloodFill(grid, i, j);  
    }  
}
```



# General Structure of Recursive Functions

- Here are 2 general constructs of recursive functions

```
if (termination condition) {  
    DO FINAL ACTION  
}  
else {  
    Take 1 step closer to  
    terminating condition  
  
    Call function RECURSIVELY  
    on smaller sub-problem  
}
```

Typically, functions that return values use this construct.

```
if (!termination condition) {  
    Take 1 step closer to  
    terminating condition  
  
    Call function RECURSIVELY  
    on smaller sub-problem  
}
```

While void recursive function use the this construct.

**Note: These are not the ONLY layouts of recursive programs, just common ones.**



# Recursive Flood Fill Algorithm

- Implementation shown in class...





# **FAST EXPONENTIATION**

COP 3502



# Fast Exponentiation

- On the first lecture on recursion we discussed the Power function:
  - But this is slow for very large exponents.

```
// Pre-conditions:  exponent is >= to 0
// Post-conditions: returns baseexponent

int Power(int base, int exponent) {

    if (exponent == 0)
        return 1;
    else
        return (base*Power(base, exponent - 1));
}
```



# Fast Exponentiation

- An example of an application that uses very large exponents is data encryption
  - One method for encryption of data (such as credit card numbers) involves modular exponentiation, with very large exponents.
    - Using the original recursive Power, it would take thousands of years just to do a single calculation.
    - Luckily, with one very simple observation, the algorithm can take a second or two with these large numbers.



# Fast Exponentiation

- The key idea is that IF the exponent is even, we can exploit the following formula:
  - $b^e = (b^{e/2}) \times (b^{e/2})$
  - For example,  $2^8 = 2^4 * 2^4$ 
    - Now, if we know  $2^4$  we can calculate  $2^8$  with a single multiplication.
    - $2^4 = 2^2 * 2^2$
    - And  $2^2 = 2 * 2$
  - Now we can return:
    - $2 * 2 = 4, 4 * 4 = 16, 16 * 16 = 256$
    - This required only 3 multiplications, instead of 7



# Fast Exponentiation

- The key idea is that IF the exponent is even, we can exploit the following formula:
  - $b^e = (b^{e/2}) \times (b^{e/2})$
  - So, In order to find,  $b^n$  we find  $b^{n/2}$ 
    - Half of the original amount
  - And then to find  $b^{n/2}$ , we find  $b^{n/4}$ 
    - Again, Half of  $b^{n/2}$
- So if we are reducing the number of multiplications we have to make in half each time, what might the run time be?
  - Log n multiplications
  - Which is much better than the original n multiplications.
- But this only works if n is even...



# Fast Exponentiation

- The key idea is that IF the exponent is even, we can exploit the following formula:
  - $b^e = (b^{e/2}) \times (b^{e/2})$
  - Since  $n$  is an integer, we have to rely on integer division which rounds down to the closest integer.
  - What if  $n$  is odd?
    - $b^n = b^{n/2} * b^{n/2} * b$
    - So  $2^9 = 2^4 * 2^4 * 2$
  - Which gives us the following formula to base our recursive algorithm on:
    - $b^n = \begin{cases} b^{n/2} * b^{n/2} & \text{if } n \text{ is even} \\ b^{n/2} * b^{n/2} * b & \text{if } n \text{ is odd} \end{cases}$



# Fast Exponentiation

- Here is the code, notice it uses the same base case as the previous Power function:

```
int PowerNew(int base, int exp) {
    if (exp == 0)
        return 1;
    else if (exp == 1)
        return base;
    else if (exp%2 == 0)
        return PowerNew(base*base, exp/2);
    else
        return base*PowerNew(base, exp-1);
}
```




# Fast Exponentiation

- Here is the code for Fast Exponentiation using Mod:

```
int modPow(int base, int exp, int n) {
    base = base%n;

    if (exp == 0)
        return 1;
    else if (exp == 1)
        return base;
    else if (exp%2 == 0)
        return modPow(base*base%n, exp/2, n);
    else
        return base*modPow(base, exp-1, n)%n;
}
```

- Even using mod, the stack is overflowed quickly, so this solution needs to be translated to an iterative solution. 

# Practice Problem

- Print a String in reverse order:
- For example, if we want to print “HELLO” backwards,
  - we first print: “O”, then we print “HELL” backwards... this is where the recursion comes in!
- See if you can come up with a solution for this





# Practice Problem

- Write a recursive function that:
  - Takes in 2 non-negative integers
  - Returns the product
    - Does NOT use multiplication to get the answer
- So if the parameters are 6 and 4
  - We get 24
  - Not using multiplication, we would have to do  $6+6+6+6$

