

AVL TREES

COP 3502

- We know that the search time for a node in a balanced binary search tree is O(log n)
 - We're dividing our search space in half each time we search the left or right branch.





But if trees get out of balance, or have deep search paths

- Their search performance deteriorates
- In the worst case instead of having an O(log n) search time
- The search time is O(n)



- So what we want is a tree that stays relatively balanced so that we can maintain the O(log n) search time,
 - BUT doesn't require too much work in maintaining the balance so that we can still have O(log n) insertion time.
- 2 Russian mathematicians, Adelson-Velski and Landis, created this type of almost balanced trees
 - known as AVL trees



The <u>AVL tree property</u> is that for any node N, the height of N's left and right subtrees must be equal or differ by 1.



height of a binary tree:

the length of the longest path from the root to a leaf. (the height of an empty tree is -1) (the height of a leaf is 0)



The <u>AVL tree property</u> is that for any node N, the height of N's left and right subtrees must be equal or differ by 1.

The <u>Balance Factor</u> is the difference in heights of the left and right subtrees at any node.

height of a binary tree:

the length of the longest path from the root to a leaf. (the height of an empty tree is -1) (the height of a leaf is 0)



Non-AVL Trees





- Now that we know what an AVL tree is,
 - now the question is how do we maintain this AVL tree property when new nodes are inserted or deleted?
 - When an imbalance is introduced to a tree, it is localized to 3 nodes and their 4 subtrees.
 - > Denote the 3 nodes as **A**, **B**, **C** in their inorder listing.
 - Here are the 4 possibilites of the imbalances that could occur:



All 4 imbalance cases can be solved by converting to the following tree:





AVL Tree Insert

- So now the question is, how can we use these rotations to actually perform an insert on an AVL tree?
 - Here are basic steps:
 - Do a normal binary search tree insert

Restore the tree based on this new leaf node, steps for restoration:

- Calculate the heights of the left and right subtrees, use this to set the potentially new height of the node.
- 2) If they are within one of each other, recursively restore the parent node.
- 3) If not, then perform the appropriate rotations on that particular node, THEN recursively restore the heights of the parent node.
- Note: No recursive call is made if the node in question is the root node and has no parents.
- Note: one rebalancing will always do the trick, though we must make the recursive calls to move up the tree so that the heights stored at each node are properly recalculated.

- The most simple insert into an AVL Tree that causes a rebalance is inserting a 3rd node into an AVL tree that creates a tree of height 2.
 - In this example, consider inserting the value 5:





In this example, consider inserting the value 20:

In this situation, the nodes 27 and 15 are balanced and we don't discover an imbalance until we trace up to 30. At this point, we label the nodes A, B and C based on our trace up the tree. The three values we passed were 27, 15 and 30, respectively. Thus, our labels are A = 15, B = 27, and C = 30.



In this example, consider inserting the value 46:





In this example, consider inserting the value 46:





In this example, consider inserting the value 46:





In this example, consider inserting the value 61:





In this example, consider inserting the value 61:





AVL Tree example

Show the resulting tree after inserting 15 into the tree below:





AVL Tree example - ANSWER

Show the resulting tree after inserting 15 into the tree below: Imbalanced, Balance Factor = 2 30 С H=1 H=3 36 20 B 43 31 10 25 18 29 15



AVL Tree example - ANSWER

Show the resulting tree after inserting 15 into the tree below:



