

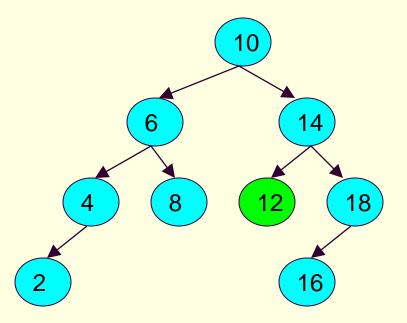
COP 3502 – Computer Science I



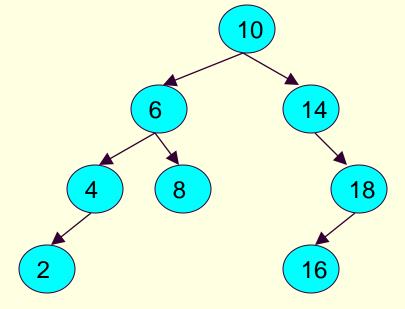
- Deletion From a Binary Search Tree
 - Deleting nodes from a BST requires some thought
 - There are 3 possible cases
 - And we deal with each in a different fashion
 - The 3 cases are:
 - 1) Deleting of a leaf node
 - Deleting a node with one child
 - 3) Deleting a node with two children
 - We examine each case separately



- Deleting a Leaf Node
 - This one is pretty easy



Initial BST with node 12 marked for deletion



BST after deletion of node 12

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- Deleting a Leaf Node
 - This one is pretty easy
 - We start by identifying the parent of the node we wish to delete
 - Which we actually do in ALL three cases
 - We then free that node, accessing it via parent:
 - free(parent->left) or
 - free(parent->right)
 - Now we need to simply update the parent's left or right pointer, signifying that the parent no longer has that child



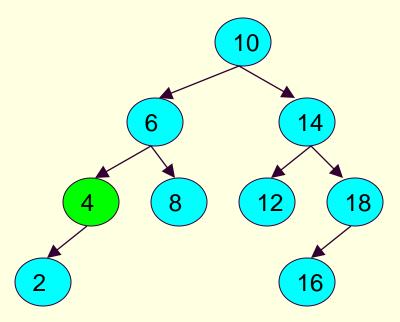
- Deleting a Leaf Node
 - This one is pretty easy
 - We start by identifying the parent of the node we wish to delete
 - Which we actually do in ALL three cases
 - Just set the appropriate node to NULL:
 - parent->left = NULL or
 - parent->right = NULL
 - So now instead of pointing to the toBeDeleted node
 - The parent simply points to NULL



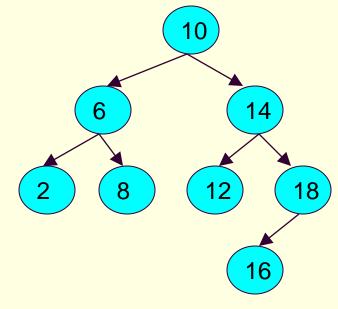
- Deleting a Node with One Child
 - This one is also not too complicated
 - But does require more thought than deleting a leaf node
 - Again, we start by finding the parent
 - meaning, the parent of the node we want to delete
 - The parent's pointer to the node is changed to now point to the deleted node's child
 - This has the effect of lifting up the deleted nodes child by one level in the tree
 - All great-great-...-great-grandchildren will lose one 'great' from their kinship designations



- Deleting a Node with One Child
 - parent->left = parent->left->left;



Initial BST with node 4 marked for deletion



BST after deletion of node 4. Node 2 has taken the place of the deleted node

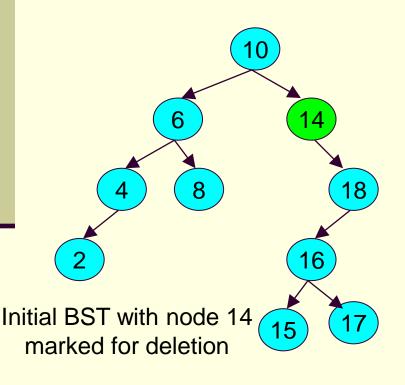
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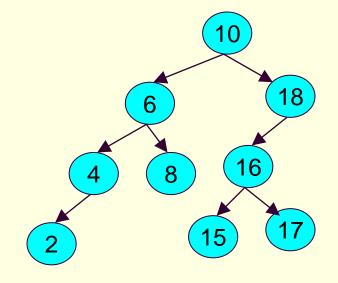


- Deleting a Node with One Child
 - The previous example illustrated when we delete a <u>left node that has a left child</u>
 - Notice that it makes no difference whether the only child is a left child or a right child
 - The next example illustrated when we delete a right node that has a right child
 - Again, the deletion simply lifts up the subtree of the deleted node by one level
 - The other possibilities is a left node that has a right child or a right node that has a left child



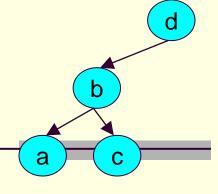
- Deleting a Node with One Child
 - parent->right = parent->right->right;





BST after deletion of node 14. Node 18 has taken the place of the deleted node and the entire subtree moved up one level.

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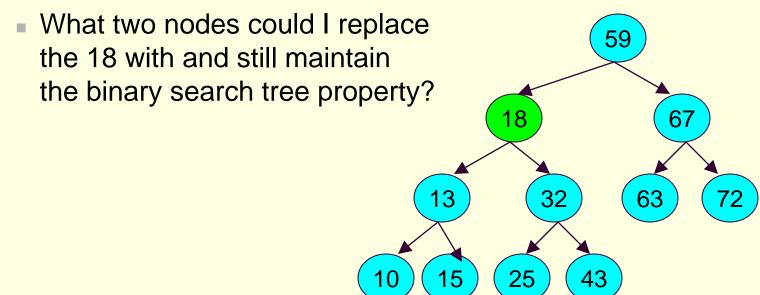


- Deleting a Node with two children
 - This is the scenario that requires a bit of thought
 - If we wish to delete node b (example above)
 - We can't just raise up b's children
 - since node d can't use its left pointer to point to more than one child
 - d's left can't point to both node a and node c

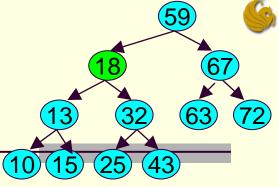
 So think about what we need to do in order to maintain the structure (and ordering property)



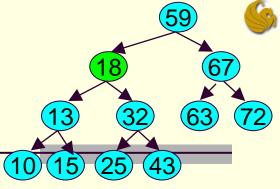
- Deleting a Node with two children
 - Consider this example tree
 - We want to delete node 18
 - Ask yourself:



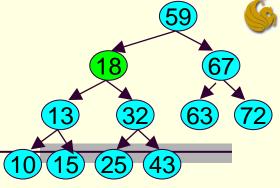
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- Deleting a Node with two children
 - Remember:
 - All the nodes to the left of 18 MUST be smaller than 18
 - All the nodes right of 18 MUST be greater than 18
 - Thus, if we delete 18
 - there are only two nodes we could put at 18's position without causing serious repercussions:
 - 1) The **maximum value in the left subtree** of node 18
 - Which is 15
 - 2) The minimum value in the right subtree of node 18
 - Which is 25



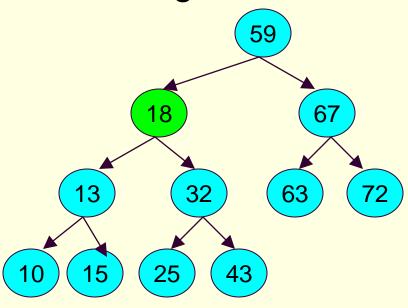
- Deleting a Node with two children
 - Thus, if we delete 18
 - There are two possible nodes that could go into 18's position:
 - 1) Node 15 (greatest value in left subtree)
 - 2) Node 25 (smallest value in right subtree)
 - We simply pick one of these to put at 18's position
 - We essentially <u>copy</u> the node to 18's position
 - Then we have to delete the actual node that we just copied
 - Meaning, if we copy node 15 into 18's positon
 - We will have two 15s
 - So we now need to delete the leaf node 15



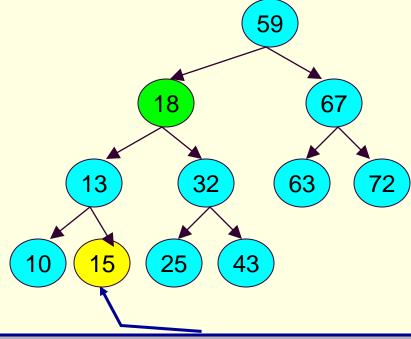
- Deleting a Node with two children
 - We are guaranteed that this node
 - Node 15 in this example
 - Has AT MOST only one child
 - Meaning it will be easy to delete!
 - Why is that? How is this guarantee true?
 - The greatest node in a left subtree cannot have two children
 - If it did, its right child would be greater than it
 - Similarly, the <u>smallest node in a right subtree</u> cannot have two children
 - If it did, its left child would be smaller than it



Deleting a Node with two children (example)



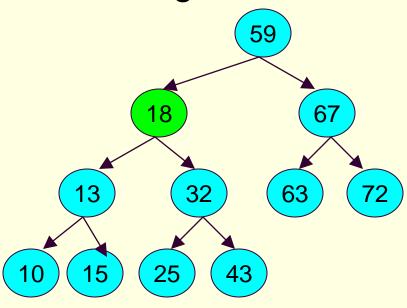
Initial BST with node 18 marked for deletion. Note that this node has two children with values 13 and 32.



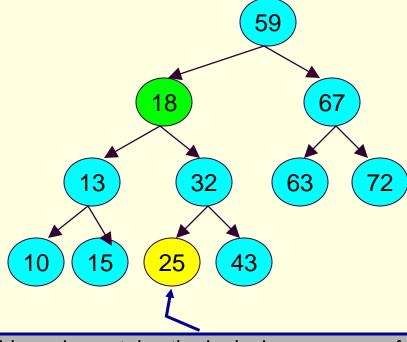
This node contains the logical <u>predecessor</u> of the node to be deleted. Note that it is the <u>greatest node</u> in the <u>left subtree</u> of the node to be deleted.



Deleting a Node with two children (example)



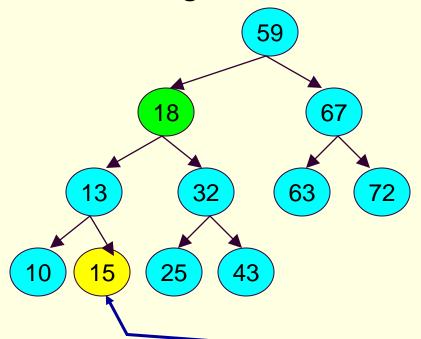
Initial BST with node 18 marked for deletion. Note that this node has two children with values 13 and 32.



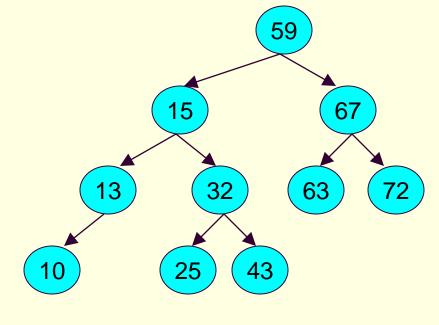
This node contains the logical <u>successor</u> of the node to be deleted. Note that it is the <u>smallest node</u> in the <u>right subtree</u> of the node to be deleted.



Deleting a Node with two children (example)



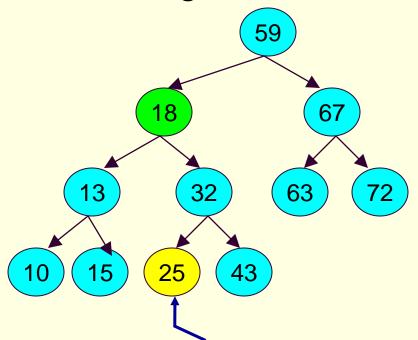
This node contains the logical <u>predecessor</u> of the node to be deleted. Note that it is the <u>greatest node</u> in the <u>left subtree</u> of the node to be deleted.



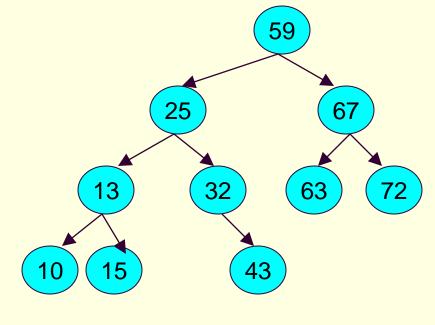
The BST after the deletion of node 18 using the replacement by the logical predecessor node.



Deleting a Node with two children (example)



This node contains the logical <u>successor</u> of the node to be deleted. Note that it is the <u>smallest node</u> in the <u>right subtree</u> of the node to be deleted.



The BST after the deletion of node 18 using the replacement by the logical successor node.



- Deleting a Node with two children
 - For the previous example,
 - The nodes that we copied to 18's position were both leaf nodes
 - How did this help?
 - Once copied over, we need to delete those nodes
 - Since they are leaves, this process is easy
 - But what if they are not leaf nodes?
 - Meaning, they have one child
 - Remember, we are guaranteed that they have AT MOST one kid
 - It is still easy!
 - We would simply be deleting a node with one child
 - Which simply "lifts" up that subtree one level



Brief Interlude: FAIL Picture



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Daily UCF Bike Fail



Courtesy of Kristina Lister



- Deleting a Node with two children
 - There's a lot going on for deletion
 - So when you examine the code,
 - You will see many auxiliary functions used, such as:
 - 1) **findNode**: returns a pointer to a node in a given tree that stores a particular value
 - 2) parent: finds the parent of a given node in a given binary tree
 - 3) minval: finds the minimum value in a given binary tree
 - 4) maxVal: finds the maximum value in a given binary tree
 - 5) isLeaf: determines if a node is a leaf node or not
 - 6) hasOnlyLeftChild: determines if a node ONLY has a left child
 - 7) hasOnlyRightChild: determines if a node ONLY has right kid



- Auxiliary functions:
 - findNode: returns a pointer to a node in a given tree that stores a particular value
 - The first step in deletion is finding the node in the tree
 - This is basically the search function from last time
 - The arguments to the function are:
 - A pointer to the root of some tree (or subtree)
 - The value we are searching for
 - IF found, the function returns a pointer to the node
 - Else, NULL is returned



- Auxiliary functions:
 - findNode: returns a pointer to a node in a given tree that stores a particular value



- Auxiliary functions:
 - parent: finds the parent of a given node in a given binary tree
 - Remember: we need to know the parent of the node we wish to delete
 - This allows us to modify the left/right pointers of the parent, effectively removing the child node
 - The arguments to the function are:
 - The root of the tree
 - 2) A pointer to the node we want to find the parents of
 - If found, a pointer to the parent node is returned



- Auxiliary functions:
 - parent: finds the parent of a given node in a given binary tree



- Auxiliary functions:
 - minVal: finds the minimum value in a given binary tree
 - Remember: when we delete a node with two children
 - We need to replace that node with either:
 - The minimum value in the right subtree, or
 - The maximum value in the left subtree
 - The argument to the function is the root of the tree
 - The function simply returns a pointer to the node containing the minimum value



- Auxiliary functions:
 - minVal: finds the minimum value in a given binary tree
 - Remember:
 - The minimum value in a given binary tree will either be the root OR it will be found in the left subtree

```
struct tree_node* minVal(struct tree_node *root) {
      // Root stores the minimal value.
      if (root->left == NULL)
            return root;

      // The left subtree of the root stores the minimal value.
      else
            return minVal(root->left);
}
```



- Auxiliary functions:
 - maxVal: finds the maximum value in a given binary tree
 - Remember:
 - The maximum value in a given binary tree will either be the root OR it will be found in the right subtree

```
struct tree_node* maxVal(struct tree_node *root) {
    // Root stores the maximal value.
    if (root->right == NULL)
        return root;

    // The right subtree of the root stores the maximal value.
    else
        return maxVal(root->right);
}
```



- Auxiliary functions:
 - isLeaf: determines if a node is a leaf node or not
 - Remember:
 - The easiest form of deletion is when the node to be deleted is a leaf node
 - That's why we need this function
 - How can you tell if a node is a leaf?
 - Leaves don't have kids!
 - BOTH the left and right pointers will be NULL

```
// Returns 1 if node is a leaf node, 0 otherwise.
int isLeaf(struct tree_node *node) {
    return (node->left == NULL && node->right == NULL);
}
```



- Auxiliary functions:
 - hasOnlyLeftChild: determines if a node ONLY has a left child
 - Remember:
 - The second easiest form of deletion is when the node to be deleted has only one child
 - We simply delete that node and "lift" the child subtree 1 level
 - How would you determine if a node has only a left kid?

```
// Returns 1 if node has a left child and no right child.
int hasOnlyLeftChild(struct tree_node *node) {
    return (node->left!= NULL && node->right == NULL);
}
```



- Auxiliary functions:
 - hasOnlyRightChild: determines if a node ONLY has a right child
 - Same as with left child
 - Just we're now checking for an only right child

```
// Returns 1 if node has a right child and no left child.
int hasOnlyRightChild(struct tree_node *node) {
    return (node->left== NULL && node->right != NULL);
}
```



- Deletion From a Binary Search Tree
 - So now we can examine the full delete function
 - Too large to put on these slides
 - Here's the link to the binary tree code:
 - http://www.cs.ucf.edu/courses/cop3502/spr2011/programs/trees/bintree.c

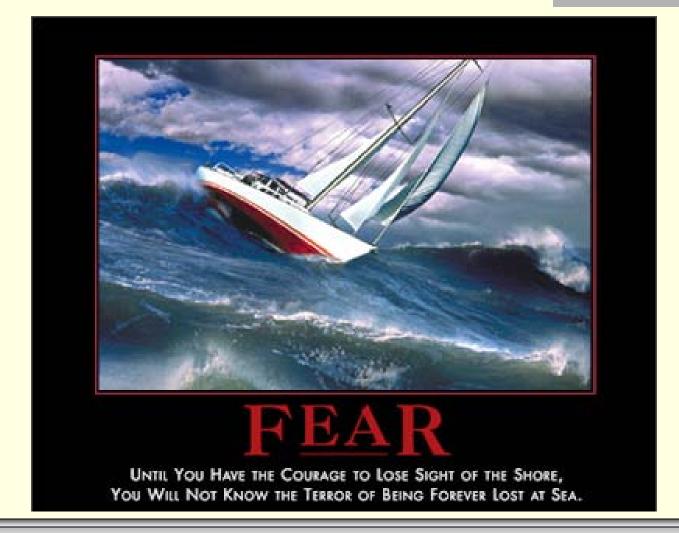


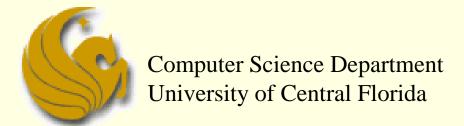
WASN'T THAT STUNNING!

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Daily Demotivator





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