Important COP 3502 Section 2 Final Exam Information

Date: Thursday, April 27, 2017
Time: 4 - 6:50 pm, Room: CB1-121

Exam Aids: 4 sheets of regular 8.5"x11" paper, front and back

Test Format:

Part A: 4 - 6 PM (Two Hour Time Limit)

EVERYONE TAKES THIS PART
If you come late, you still have to turn it in at 6 pm.
100 points

Part B: 6 - 6:45 PM

ONLY THOSE WHO DIDN'T DO COMMUNITY SERVICE DO THIS PART.
If you finish Part A early, you can use more time on this part. It gets collected at 6:45 pm regardless of when you started it.

For both sections the format will be free response (coding, tracing, short answer, math questions, etc.)

Areas of Focus: Bitwise Operators, Recurrence Relations, Linked List and Binary Tree Coding (but all areas are fair game.)

SARC REVIEW SESSION - KRYSTAL SILER
DATE: SUNDAY, APRIL 23rd
TIME: 2:30 - 4:30 pm
LOCATION: KEY WEST AB

Outline of Topics for the Exam

I. Basics of C – if, loops, functions, array, strings, files

II. Mathematical Background

- **a. base conversion (2, 10, 16, other)**
- **b. logs and exponents**
- **c. sums**
- **d. Big-Oh Timing problems**
- **e. Recurrence Relations**

III. Pointers and Dynamic Arrays

- **a. how to allocate space dynamically**
  (array, 2d array, array of struct, array of ptr to struct, linked list node, bin tree node, etc.)
- **b. how to free space**
- **c. how to "resize" an existing array**

IV. Structs

- **a. how to declare structs**
- **b. how to use pointers to structs**
- **c. how to use arrays of structs**
- **d. how to pass structs into functions**

V. Recursion

- **a. Fibonacci**
- **b. Factorial**
- **c. Towers of Hanoi**
- **d. Binomial Coefficients**
- **e. Binary Search**
- **f. Generating Permutations**
- **g. Fast Modular Exponentiation**
- **h. Floodfill**

VI. Algorithm Analysis

- **a. Average case vs. Worst case**
- **b. Determining a Big-Oh bound**
- **c. Use of sums, etc.**

VII. Linked Lists

- **A. Creating Nodes**
- **B. Insertion, Searching**
- **C. Deletion**
- **D. Circularly linked**
- **E. Doubly linked**

VIII. Stacks

- **A. Array Implementation**
- **B. Dynamically Sized Array Implementation**
- **C. Linked List Implementation**
- **D. Efficiency of push, pop**
- **E. Determining the Value of Postfix Expressions**
- **F. Converting Infix to Postfix**

IX. Queues

- **A. Array Implementation**
- **B. Dynamically Sized Array Implementation**
- **C. Linked List Implementation**
- **D. Efficiency of Enqueue and Dequeue**

X. Binary Search Trees

- **A. Creating Nodes**
- **B. Tree Traversals (preorder, inorder, postorder)**
- **C. Insertion**
- **D. Searching**
- **E. Deletion**
- **F. Code Tracing**
- **G. Writing Code (recursive)**

XI. AVL Trees

- **A. AVL Tree Property**
- **B. Identifying nodes A, B and C for both insert and delete**
- **C. Restructuring for both insert and delete**
- **D. Delete may have multiple restructures**

XII. Binary Heaps

- **A. percolateUp**
- **B. percolateDown**
- **C. Insert**
- **D. deleteMin**
- **E. makeHeap**
- **F. Heap Sort**

XIII. Hash Tables

- **A. linear probing replacement technique**
- **B. quadratic probing replacement technique**
- **C. linear chaining hashing**

XIV. Sorting and Selection

- **A. Bubble Sort**
- **B. Insertion Sort**
- **C. Selection Sort**
- **D. Merge Sort**
- **E. Quick Sort**
- **F. Quickselect**

XV. Bit-Wise Operators

- **A. left shift, right shift, and, or, xor, complement**
- **B. How to use a number to indicate a subset.**
- **C. How to iterate through all possible subsets w/bitmask.**

XVI. Backtracking

- **A. Use in Eight Queens Problem**
- **B. Idea for Sudoku Solving**
- **C. Trying all possibilities and stopping at a dead end.**

Mathematical Background

a) With respect to binary, remember the algorithm to convert to and from binary and decimal. To go from decimal to binary, use repeated division and mod by 2. In general, to go from base 10 to another base, use repeated division and mod by that base.

b) Make sure you know how to apply some basic log rules, including adding and subtracting two logs, the power rule, and that the log and exponent functions are inverses of each other.

c) You should be able to handle sums between various bounds of a constant and a linear function.

d) To set up the Big-Oh problems given in this course, make sure you set up a function that solves for the running time of an algorithm and use the given information to solve for the unknown constant.

e) I will definitely test recurrence relations - this is essentially a part of the mathematics background for the course. You will need to know how to solve them using the iteration technique.

Structs

Just make sure you can handle the various different modes in which structs are used. (By themselves, in an array, inside of another struct, etc.) Also, pay attention to the difference between a struct and a pointer to a struct. The latter is more typically used. Also make sure that you know the difference between an array of struct and an array of pointers to structs.

Recursion

I am likely to ask a recursive coding question and a recursive tracing question.

Remember, that often, recursion fits into one of two constructs:

1) if (!terminating condition) { do work }

2) if (terminating condition) { finish } else { do work, call rec }

However, not all recursive algorithms, follow these two constructs. Consider the permutation algorithm. It does not just make one recursive call or even two.

The main idea behind recursion is to take a problem of a certain size, do some work and finish solving the problem by solving a different version of the same problem of a smaller size.

The toughest part is "seeing" how you can break a problem down into a smaller recursive solution.

My favorite analogy is imagining that someone else has written a function to solve the task already, and your job is to write a function to solve the task at hand, but you can call the function that someone else has written as an aid, just not with the same parameters.

The more complicated recursive patterns, such as floodfill and permutation do have loops in them, because there are many "directions" to recursively move in, so to speak.

Algorithm Analysis

The key goal here is to determine the number of simple statements that are run by a segment of code. Typically, a summation can be set up to determine this, in terms of some input parameter. For recursive functions, we have to set up recurrence relations

Linked Lists

Make sure you look at all the mechanics involved in inserting and deleting nodes from a linked list. Also, consider slight "twists" in the design of a linked list, like a circular list or a doubly linked list.

The most important pieces of information dealing with linked lists:

1) Watch out for NULL pointer errors
2) Make sure you don't "lose" the list.
3) Make sure you connect the links in the proper order.
4) Don't forget to "patch" everything up for some operation.
5) Determine when it is necessary and not necessary to use an extra helper pointer.
6) Determine when it is necessary and not necessary for a function to return a pointer to the beginning of the list.

Stacks and Queues

Stacks are last in, first out structures and Queues are first in, first out structures.

Typically, stack and queue operations occur in O(1) time if the structure is efficiently implemented.

The array implementation of a stack just needs the array, its size and an integer storing the index to the top of the stack.

In a queue, you need more information. Typically, we also need to store the number of items currently in the queue as well as the index to the front of the queue.

Make sure you understand how the implementation here affects run-time. (For example, if we implemented a queue with an array but always had the front of the queue be index 0, a dequeue could potentially take O(n) time to move each element forward one slot.)

In a linked list implementation of a queue, a pointer is needed to the back of the queue as well as the front. But, for a stack, only the latter is needed.

Binary Search Trees

Many of the concerns necessary with linked lists translate over to binary trees. One key point about binary trees:

Recursion is even more important/useful for binary trees than linked lists. In particular, it's very difficult to think about how to iteratively go through all the nodes in a binary tree, but recursively, the code is reasonably concise and simple.

AVL Trees

All I will test upon for these is how to do rotations after insertions and deletions. Remember the following ideas:

1) All insertions can be fixed with a single rotation.
2) Deletions may need more than one rotation to be fixed. But all errors are propagated up the tree.

To find WHERE to do a rotation, start at the inserted node or the parent of the deleted node, going "up" the tree, node by node, until you find an offending node. Then perform the appropriate rotation. From there, continue up the tree.

Binary Heaps

A binary heap is an efficient data structure to use if the main operations that need to be handled are inserting items and deleting the minimum item. Both tasks can be done in O(lg n) time, where n represents the number of items stored in the heap.

Typically, a binary heap is implemented using an array. The implementation I showed in class stores the first element of the heap in index 1. In this convention, the left child of a node at index i is at index 2i and the right child is at index 2i+1. A node stored at index i has its parent at index i/2.

The key to getting a binary heap working are the subroutines identified in class as "percolateUp" and "percolateDown".

Hash Tables

A hash table is an efficient data structure that easily allows for inserting items and searching for items. The main problem with hash tables is collisions, since hash functions are many-to-one functions (meaning that two different input values can hash to the same output location.) There are three ways to deal with collisions discussed in class:

a) Linear Probing
b) Quadratic Probing
c) Linear Chaining Hashing

The first couple are reasonable so long as the table is no more than half full. The last is most probably the best way to deal with the issue.

Backtracking

I spent a decent amount of time looking at three examples of code using backtracking: Eight Queens, Sudoku and the Digit Divisibility problem. I may have you fill in an incomplete function that executes some sort of backtracking. The key to backtracking is that it is very similar to brute force, except we *skip* building off of partial solutions that we know are doomed to fail. This is how a Sudoku puzzle which looks like it has something like 9^{50} options (if 31 squares are given) can still be solved very quickly. Mostly, we set up our code to run a brute force solution, but then in the loop where we are filling in the next portion of the partial solution, we skip over unviable options. In the code examples I've shown you, I typically implement this with an if statement that screens out invalid cases inside the for loop. This if statement, when triggered, executes continue, which skips the ensuing recursive call that would have been made, building off of that partial solution. (For example, in the Digit Divisibility problem, when we see that 326 isn't divisible by 3, we skip trying *ANY* number that starts 326.)

Bit-Wise Operators

Bitwise and = &
Bitwise or = |
Bitwise xor = ^
Bitwise not = ~

When using bitwise operators, it's important to remember what each of the 32 bits in an int represents. In particular, the most significant bit is worth -2^{31}. The rest of the bits are worth their unsigned value.

We can use a single integer to store an array of true/false. In this manner, we can store items of a subset (if bit 0 is 1, then item 0 is in our subset, if bit 7 is 0, then item 7 is not in our subset, and so forth).

Lots of real world problems can efficiently be encoded with the bitwise operators, such as finding the intersection of two sets (bitwise and), union of two sets (bitwise or), or exclusive or of two sets (bitwise exclusive or).

We can code up a brute force solution without recursion going through each subset of a set using bitwise operators as well.