

Linear Search vs Binary Search



Computer Science Department
University of Central Florida

COP 3502 – Computer Science I



More Announcements

- Additional Announcements:
 - Best way to get in touch with me:
 - Again, for anything Program/Assignment related, contact the TA
 - but when you need to contact me, email is the best way
 - jcazalas@cs.ucf.edu
 - I check email often and respond quickly
 - If you want a quick response, **do NOT email through Webcourses**, as I check that weekly, at best



More Announcements

- Additional Announcements:
 - Webcourses is up!
 - What is webcourses for?
 - Only used for submitting assignments, taking quizzes, and used to store your grades

 - The course website is where ALL content for the course will be stored



More Announcements

- Additional Announcements:
 - Program 1 is now assigned
 - Main purpose of Program 1:
 - WAKE UP call to those who have forgotten C
 - It is not a difficult program.
 - May take only 3 hours for some of you
 - May take upwards of 20+ for others

- Silver Lining:
 - If you put in the time investment to correctly do Program 1, you will be prepared for any program that is assigned during the semester



More Announcements

- Additional Announcements:
 - Class attendance:
 - The reality is that those who come to class
 - And are MENTALLY present as well as physically
 - Do end up with a better grasp of the material
 - Therefore, to encourage attendance, I will randomly pick seven days throughout the semester and will take attendance on those days.
 - You will simply print and sign your name on a piece of paper
 - To avoid excuses for missing days, only five of the days will count...allowing you two free misses
 - This will count as five bonus points on the Final Exam



More Announcements

- Additional Announcements:
 - PPT slides:
 - A quick comment on these PPT slides
 - I hate when professors talk for 50 minutes and then when you go to study their slides for the exam, you find 8 slides in the form of an outline.
 - So unless you took notes like crazy that day, or have a great memory, you have no idea what was talked about
 - So you will notice these slides are the opposite
 - What could be done with 20 slides, I may use 60.
 - Simply because I want to give you a good study aid when it comes to the exams.
 - So, yeah, the slides are EPIC...ENJOY!



Linear Search

- Searching from C-Programming class
 - In COP 3223, we studied how to find a value in an array
 - Look at each value in the array
 - Compare it to what we're looking for
 - If we see the value we are searching for,
 - Return that we've found it!
 - Otherwise, if we've iterated through the entire array and haven't located the value,
 - Return that the value isn't in the array



Linear Search

- Searching from C-Programming class
 - Your code should look something like this:

```
int search(int array[], int len, int value) {  
  
    int i;  
    for (i=0; i<len; i++) {  
        if (array[i] == value)  
            return 1;  
    }  
    return 0;  
}
```




Linear Search

- Searching from C-Programming class
 - Analyze code:
 - Clearly, if the array is unsorted, this algorithm is optimal
 - They ONLY way to be sure that a value isn't in the array is to look at every single spot of the array
 - Just like you can't be sure that you DON'T have some piece of paper or form unless you look through ALL of your pieces of paper
 - But we ask a question:
 - Could we find an item in an array faster if it were already sorted?



Binary Search

- Number Guessing Game from childhood
 - Remember the game you most likely played as a child
 - I have a secret number between 1 and 100.
 - Make a guess and I'll tell you whether your guess is too high or too low.
 - Then you guess again. The process continues until you guess the correct number.
 - Your job is to MINIMIZE the number of guesses you make.



Binary Search

- Number Guessing Game from childhood
 - What is the first guess of most people?
 - 50.
 - Why?
 - No matter the response (too high or too low), the most number of possible values for your remaining search is 50 (either from 1-49 or 51-100)
 - Any other first guess results in the risk that the possible remaining values is greater than 50.
 - Example: you guess 75
 - I respond: too high
 - So now you have to guess between 1 and 74
 - 74 values to guess from instead of 50



Binary Search

- Number Guessing Game from childhood
 - Basic Winning Strategy
 - Always guess the number that is halfway between the lowest possible value in your search range and the highest possible value in your search range

- Can we now adapt this idea to work for searching for a given value in an array?



Binary Search

■ Array Search

- We are given the following sorted array:

index	0	1	2	3	4	5	6	7	8
value	2	6	19	27	33	37	38	41	118

- We are searching for the value, 19
- So where is halfway between?
 - One guess would be to look at 2 and 118 and take their average (60).
 - But 60 isn't even in the list
 - And if we look at the number closest to 60
 - It is almost at the end of the array



Binary Search

- Array Search

- We quickly realize that if we want to adapt the number guessing game strategy to searching an array, we **MUST** search in the middle **INDEX** of the array.
- In this case:
 - The lowest index is 0
 - The highest index is 8
 - So the middle index is 4



Binary Search

■ Array Search

■ Correct Strategy

- We would ask, “is the number I am searching for, 19, greater or less than the number stored in index 4?”
 - Index 4 stores 33
- The answer would be “less than”
- So we would modify our search range to in between index 0 and index 3
 - Note that index 4 is no longer in the search space
- We then continue this process
 - The second index we’d look at is index 1, since $(0+3)/2=1$
 - Then we’d finally get to index 2, since $(2+3)/2 = 2$
 - And at index 2, we would find the value, 19, in the array



Binary Search

■ Binary Search code:

```
int binsearch(int a[], int len, int value) {  
  
    int low = 0, high = len-1;  
    while (low <= high) {  
        int mid = (low+high)/2;  
        if (value < a[mid])  
            high = mid-1;  
        else if (value > a[mid])  
            low = mid+1;  
        else  
            return 1;  
    }  
  
    return 0;  
}
```




Binary Search

- Binary Search code:
 - At the end of each array iteration, all we do is update either low or high
 - This modifies our search region
 - Essentially halving it



Binary Search

■ Efficiency of Binary Search

■ Analysis:

- Let's analyze how many comparisons (guesses) are necessary when running this algorithm on an array of n items

First, let's try $n = 100$

- After 1 guess, we have 50 items left,
- After 2 guesses, we have 25 items left,
- After 3 guesses, we have 12 items left,
- After 4 guesses, we have 6 items left,
- After 5 guesses, we have 3 items left,
- After 6 guesses, we have 1 item left
- After 7 guesses, we have 0 items left.



Binary Search

■ Efficiency of Binary Search

■ Analysis:

■ Notes:

- The reason for the last iteration is because the number of items left represent the number of other possible values to search
 - We need to reduce this to 0.
- Also, when n is odd, such as when $n=25$
 - We search the middle element, # 13
 - There are 12 elements smaller than 13
 - And 12 elements bigger than 13
 - This is why the number of items is slightly less than $\frac{1}{2}$ in those cases



Binary Search

- Efficiency of Binary Search
 - Analysis:
 - General case:
 - After 1 guess, we have $n/2$ items left
 - After 2 guesses, we have $n/4$ items left
 - After 3 guesses, we have $n/8$ items left
 - After 4 guesses, we have $n/16$ items left
 - ...
 - After k guesses, we have $n/2^k$ items left



Binary Search

■ Efficiency of Binary Search

■ Analysis:

- General case:
- So, after k guesses, we have $n/2^k$ items left
- The question is:
 - How many k guesses do we need to make in order to find our answer?
 - Or until we have one and only one guess left to make?
- So we want to get only 1 item left
- If we can find the value that makes the above fraction equal to 1, then we know that in one more guess, we'll narrow down the item



Binary Search

■ Efficiency of Binary Search

■ Analysis:

- General case:
- So, after k guesses, we have $n/2^k$ items left
 - Again, we want only 1 item left
 - So set this equal to 1 and solve for k

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

- This means that a binary search roughly takes $\log_2 n$ comparisons when searching in a sorted array of n items



Binary Search

■ Efficiency of Binary Search

■ Analysis:

- Runs in logarithmic ($\log n$) time
- This is MUCH faster than searching linearly
- Consider the following chart:

<u>n</u>	<u>log n</u>
8	3
1024	10
65536	16
1048576	20
33554432	25
1073741824	30

- Basically, any $\log n$ algorithm is SUPER FAST.



Binary Search

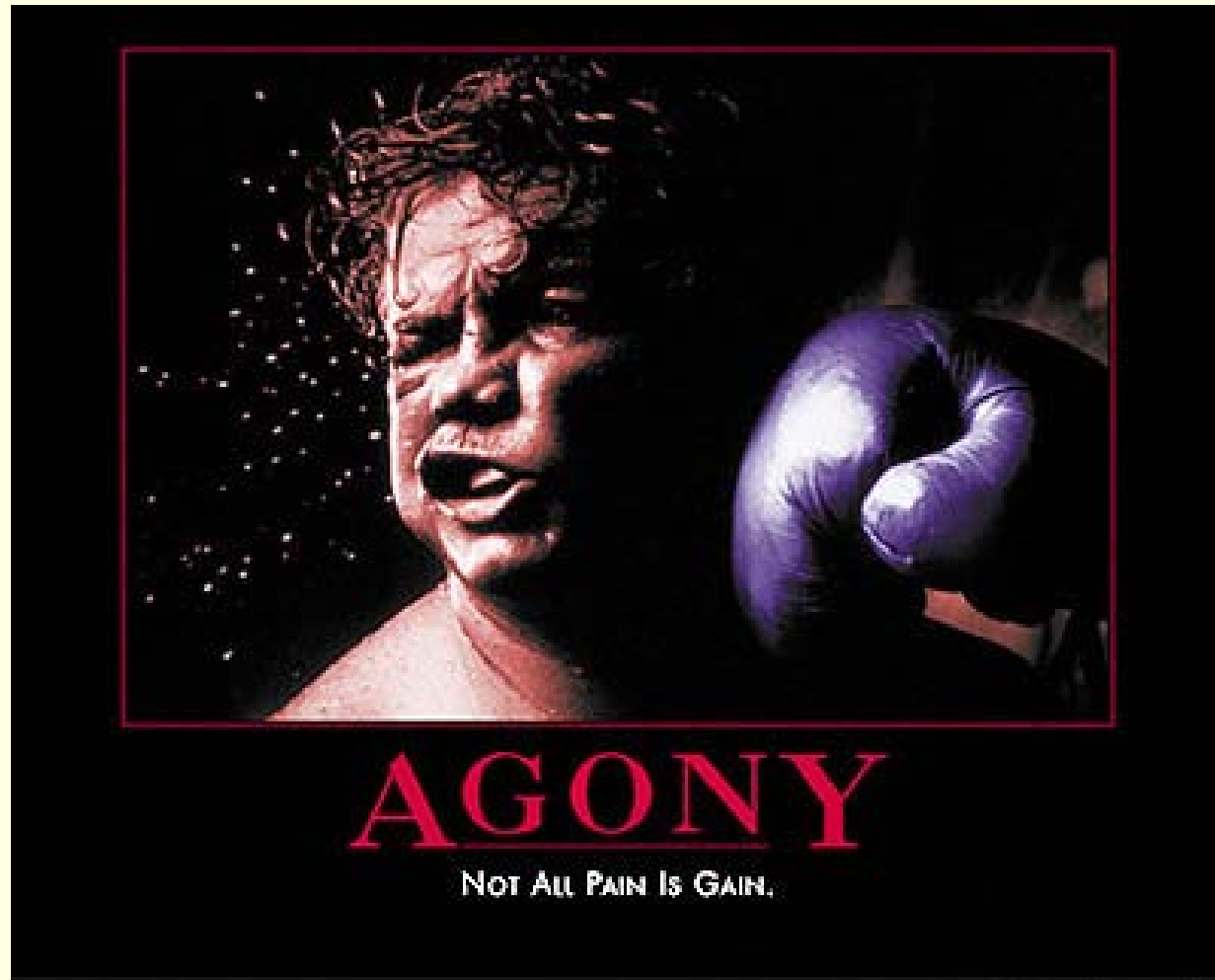
WASN'T

THAT

INCREDIBLE!



Daily Demotivator



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