

# More Recursion



Computer Science Department  
University of Central Florida

*COP 3502 – Computer Science I*



# Recursion

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- What is Recursion? (*reminder from last time*)
  - From the programming perspective:
  - Recursion solves large problems by **reducing** them to **smaller** problems of the **same form**
  - Recursion is a function that invokes itself
    - Basically **splits** a problem into **one or more SIMPLER versions of itself**
    - And we must have a way of stopping the recursion
    - So the function must have some sort of calls or conditional statements that can actually terminate the function



# Recursion - Factorial

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- Example: Compute Factorial of a Number
  - What is a factorial?
    - $4! = 4 * 3 * 2 * 1 = 24$
    - In general, we can say:
    - $n! = n * (n-1) * (n-2) * \dots * 2 * 1$
    - Also,  $0! = 1$ 
      - (just accept it!)



# Recursion - Factorial

## ■ Example: Compute Factorial of a Number

### ■ Recursive Solution

- Mathematically, factorial is already defined recursively
  - Note that each factorial is related to a factorial of the next smaller integer

- $4! = 4 * 3 * 2 * 1 = 4 * (4-1)! = 4 * (3!)$

- Right?

- Another example:

- $10! = 10 * 9 * 8 * 7 * 6 * 5 * 4 * 3 * 2 * 1$

- $10! = 10 * (9!)$

This is clear right?  
Since 9! clearly is equal to  
 $9 * 8 * 7 * 6 * 5 * 4 * 3 * 2 * 1$



# Recursion - Factorial

## ■ Example: Compute Factorial of a Number

### ■ Recursive Solution

- Mathematically, factorial is already defined recursively
  - Note that each factorial is related to a factorial of the next smaller integer
- Now we can say, in general, that:
- $n! = n * (n-1)!$
- But we need something else
  - We need a stopping case, or this will just go on and on and on
  - NOT good!
- We let  $0! = 1$

■ So in “math terms”, we say

- $n! = 1$  if  $n = 0$
- $n! = n * (n-1)!$  if  $n > 0$



# Recursion - Factorial

- How do we do this recursively?
  - We need a function that we will call
    - And this function will then call itself (recursively)
      - until the stopping case ( $n = 0$ )

```
#include <stdio.h>

void Fact(int n);
int main(void) {
    int factorial = Fact(10);
    printf("%d\n", factorial);
    return 0;
}
```

## Here's the Fact Function

```
int Fact (int n) {
    if (n = 0)
        return 1;
    else
        return (n * fact(n-1));
}
```

- This program prints the result of  $10*9*8*7*6*5*4*3*2*1$ :
  - 3628800

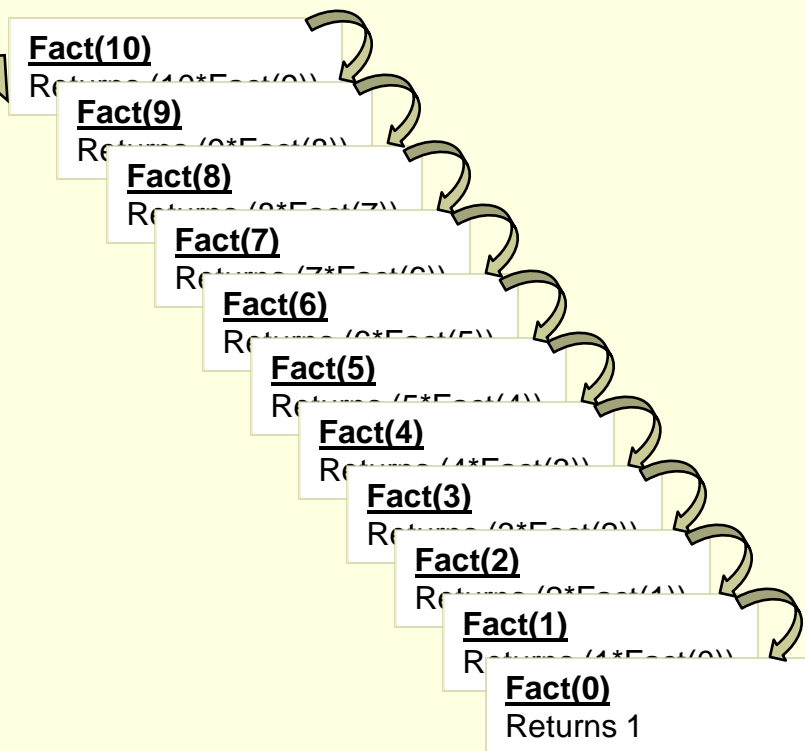


# Recursion - Factorial

- Here's what's going on...in pictures

```
#include <stdio.h>

void Fact(int n);
int main(void) {
    int factorial = Fact(10);
    printf("%d\n", factorial);
    return 0;
}
```



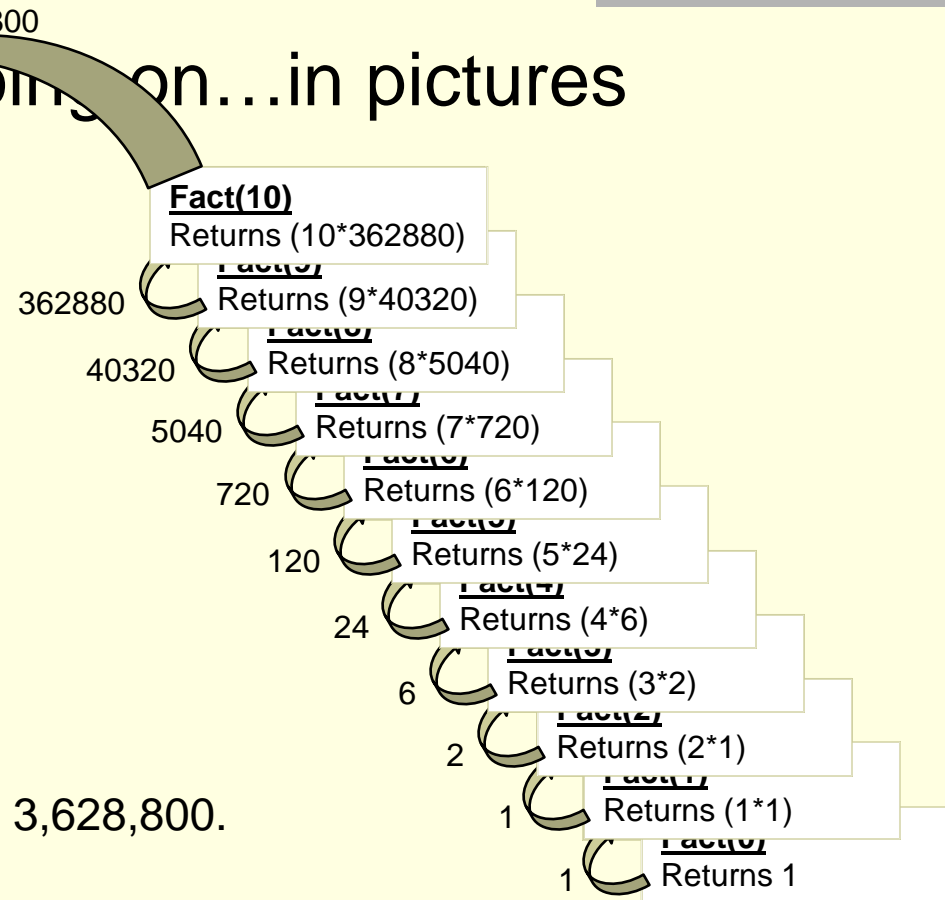


# Recursion - Factorial

- Here's what's going on...in pictures

```
#include <stdio.h>

void Fact(int n);
int main(void) {
    int factorial = Fact(10);
    printf("%d\n", factorial);
    return 0;
}
```



- Now factorial has the value 3,628,800.





# Recursion: General Structure

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- General Structure of Recursive Functions:
  - What we can determine from previous examples:
    - When we have a problem, we want to break it into chunks
    - Where one of the chunks is a smaller version of the same problem
  - Factorial Example:
    - We utilized the fact that  $n! = n \cdot (n-1)!$
    - And we realized that  $(n-1)!$  is, in essence, an easier version of the original problem
    - Right?
    - We all should agree that  $9!$  is a bit easier than  $10!$



# Recursion: General Structure

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- General Structure of Recursive Functions:
  - What we can determine from previous examples:
    - Eventually, we break down our original problem to such an extent that the small sub-problem becomes quite easy to solve
    - At this point, we don't make more recursive calls
    - Rather, we directly return the answer
    - Or complete whatever task we are doing
  - This allows us to think about a general structure of a recursive function



# Recursion: General Structure

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- General Structure of Recursive Functions:
  - Basic structure has 2 main options:
    - 1) Break down the problem further
      - Into a smaller sub-problem
    - 2) OR, the problem is small enough on its own
      - Solve it
  - In programming, when we have two options, we use an if statement
  - So here are our two constructs of recursive functions...



# Recursion: General Structure

- General Structure of Recursive Functions:

- 2 general constructs:

- **Construct 1:**

```
if (terminating condition) {
    DO FINAL ACTION
}
else {
    Take one step closer to terminating condition
    Call function RECURSIVELY on smaller subproblem
}
```

- Functions that return values take on this construct



# Recursion: General Structure

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- General Structure of Recursive Functions:

- 2 general constructs:

- **Construct 2:**

```
if (!(terminating condition) ) {  
    Take a step closer to terminating condition  
    Call function RECURSIVELY on smaller subproblem  
}
```

- void recursive functions use this construct



# Recursion: General Structure

## ■ Example using Construct 1

- Our function (Sum Integers):
  - Takes in one positive integer parameter,  $n$
  - Returns the sum  $1+2+\dots+n$
  - So our recursive function must sum all the integers up until (and including)  $n$
- How do we do this recursively?
  - We need to solve this in such a way that part of the solution is a sub-problem of the EXACT same nature of the original problem.



# Recursion: General Structure

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- Example using Construct 1
  - Our function:
    - Using  $n$  as the input, we define the following function
      - $f(n) = 1 + 2 + 3 + \dots + n$ 
        - Hopefully it is clear that this is our desired function
      - Example:
        - $f(10) = 1 + 2 + 3 + 4 + 5 + 6 + 7 + 8 + 9 + 10$
  - So the question is:
  - Given this function,  $f(n)$ , how do we make it recursive???



# Recursion: General Structure

- Example using Construct 1
  - Our function:
    - Using  $n$  as the input, we define the following function
      - $f(n) = 1 + 2 + 3 + \dots + n$
  - REMEMBER:
    - We want a function that solves this same problem
    - But we want that problem to be recursive:
      - It should solve  $f(n)$  by reducing it to a smaller problem, **but of the same form**
    - Just like the factorial example:  $n! = n * (n-1)!$ 
      - $(n-1)!$  was a smaller form of  $n!$
    - So think, what is a “smaller form” of our function,  $f(n)$





# Recursion: General Structure

## ■ Example using Construct 1

### ■ Our function:

- Using  $n$  as the input, we define the following function

- $f(n) = 1 + 2 + 3 + \dots + n$

- So to make this recursive, can we say:

- $f(n) = 1 + (2 + 3 + \dots + n)$



- Does that “look” recursive?
- Is there a sub-problem that is the EXACT same form as the original problem?
  - NO!
- $2+3+\dots+n$  is NOT a sub-problem of the form  $1+2+\dots+n$



# Recursion: General Structure

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- Example using Construct 1
  - Our function:
    - Using  $n$  as the input, we get the following function
      - $f(n) = 1 + 2 + 3 + \dots + n$
    - Let's now try this:
      - $f(n) = 1 + 2 + \dots + n = n + (1 + 2 + \dots + (n-1))$
    - AAAHHH.
    - Here we have an expression
      - $1 + 2 + \dots + (n-1)$
    - which IS indeed a sub-problem of the same form



# Recursion: General Structure

- Example using Construct 1
  - Our function:
    - Using  $n$  as the input, we get the following function
      - $f(n) = 1 + 2 + 3 + \dots + n$
    - So now we have:
      - $f(n) = 1 + 2 + \dots + n = n + (1 + 2 + \dots + (n-1))$
    - Now, realize the following:
      - Use an example:
        - $f(10) = 1 + 2 + \dots + 10 = 10 + (1 + 2 + \dots + 9)$
        - And what is  $(1 + 2 + \dots + 9)$ ? **It is f(9)!**
        - So look at what we can say:
        - We can say that,  **$f(10) = 10 + f(9)$**



# Recursion: General Structure

## ■ Example using Construct 1

### ■ Our function:

- Using  $n$  as the input, we get the following function

- $f(n) = 1 + 2 + 3 + \dots + n$

- So now we have:

- $f(n) = 1 + 2 + \dots + n = n + (1 + 2 + \dots + (n-1))$

- Now, realize the following:

- So, in general, we have:  **$f(n) = n + f(n-1)$**

- Right?

- Just like  $f(10) = 10 + f(9)$

- So, we've defined our function,  $f(n)$ , to be in terms of a smaller version of itself...in terms of  $f(n-1)$



# Recursion: General Structure

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- Example using Construct 1
  - Our function:
    - Using  $n$  as the input, we get the following function
      - $f(n) = 1 + 2 + 3 + \dots + n$
    - So now we have:
      - $f(n) = 1 + 2 + \dots + n = n + (1 + 2 + \dots + (n-1))$
    - Now, realize the following:
      - So here is our function, defined recursively
      - **$f(n) = n + f(n-1)$**



# Recursion: General Structure

- Example using Construct 1
  - Our function (now recursive):
    - $f(n) = n + f(n-1)$
    - Reminder of construct 1:

```
if (terminating condition) {  
    DO FINAL ACTION  
}  
else {  
    Take one step closer to terminating condition  
    Call function RECURSIVELY on smaller subproblem  
}
```



# Recursion: General Structure

- Example using Construct 1
  - Our function:
    - $f(n) = n + f(n-1)$
    - Reminder of construct 1:
    - So we need to determine the terminating condition!
    - We know that  $f(0) = 0$ 
      - So our terminating condition can be  $n = 0$
    - Additionally, our recursive calls need to return an expression for  $f(n)$  in terms of  $f(k)$ 
      - for some  $k < n$
    - We just found that  $f(n) = n + f(n-1)$
    - So now we can write our actual function...



# Recursion: General Structure

- Example using Construct 1
  - Our function:  $f(n) = n + f(n-1)$

```
// Pre-condition: n is a positive integer.  
// Post-condition: Function returns the sum  
// 1 + 2 + ... + n  
int sumNumbers(int n) {  
  
    if ( n == 0 )  
        return 0;  
    else  
        return (n + sumNumbers(n-1));  
}
```





# Recursion: General Structure

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- Another example using Construct 1
  - Our function:
    - Calculates  $b^e$ 
      - Some base raised to a power,  $e$
      - The input is the base,  $b$ , and the exponent,  $e$
      - So if the input was 2 for the base and 4 for the exponent
        - The answer would be  $2^4 = 16$
    - How do we do this recursively?
      - We need to solve this in such a way that part of the solution is a sub-problem of the EXACT same nature of the original problem.



# Recursion: General Structure

## ■ Another example using Construct 1

### ■ Our function:

- Using  $b$  and  $e$  as input, here is our function

- $f(b,e) = b^e$

- So to make this recursive, can we say:

- $f(b,e) = b^e = b * b^{(e-1)}$

- Does that “look” recursive?

- YES it does!

- Why?

- Cuz the right side is indeed a sub-problem of the original

- We want to evaluate  $b^e$

- And our right side evaluates  $b^{(e-1)}$

### **Example with numbers:**

$f(2,4) = 2^4 = 2 * 2^3$

---So we solve the larger problem ( $2^4$ ) by reducing it to a smaller problem ( $2^3$ ).



# Recursion: General Structure

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- Another example using Construct 1
  - Our function:
    - $f(b,e) = b*b^{(e-1)}$
    - Reminder of construct 1:

```
if (terminating condition) {  
    DO FINAL ACTION  
}  
else {  
    Take one step closer to terminating condition  
    Call function RECURSIVELY on smaller subproblem  
}
```



# Recursion: General Structure

---

- Another example using Construct 1
  - Our function:
    - $f(b,e) = b*b^{(e-1)}$
    - Reminder of construct 1:
    - So we need to determine the terminating condition!
    - We know that  $f(b,0) = b^0 = 1$ 
      - So our terminating condition can be when  $e = 0$
    - Additionally, our recursive calls need to return an expression for  $f(b,e)$  in terms of  $f(b,k)$ 
      - for some  $k < e$
    - We just found that  $f(b,e) = b*b^{(e-1)}$
    - So now we can write our actual function...



# Recursion: General Structure

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- Another example using Construct 1
  - Our function:

```
// Pre-conditions: e is greater than or equal to 0.  
// Post-conditions: returns be.  
int Power(int base, int exponent) {  
  
    if ( exponent == 0 )  
        return 1;  
    else  
        return (base*Power(base, exponent-1));  
}
```



# Recursion: General Structure

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## ■ Example using Construct 2

- Remember the construct:
  - This is used when the return type is void

```
if (!(terminating condition) ) {  
    Take a step closer to terminating condition  
    Call function RECURSIVELY on smaller subproblem  
}
```



# Recursion: General Structure

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- Example using Construct 2
  - Our function:
    - Takes in a string (character array)
    - Also takes in an integer, the length of the string
    - The function simply prints the string in REVERSE order
  - So what is the terminating condition?
    - We will print the string, in reverse order, character by character
    - So we terminate when there are no more characters left to print
    - The 2<sup>nd</sup> argument to the function (length) will be reduced until it is 0 (showing no more characters left to print)



# Recursion: General Structure

- Example using Construct 2

- Our function:

```
void printReverse(char word[], int length) {  
    if (length > 0) {  
        printf("%c", word[length-1]);  
        printReverse(word, length-1);  
    }  
}
```

- What's going on:

- Let's say the word is "computer"
      - 8 characters long
    - So we print word[7]
      - Which would refer to the "r" in computer





# Recursion: General Structure

- Example using Construct 2
  - Our function:

```
void printReverse(char word[], int length) {  
    if (length > 0) {  
        printf("%c", word[length-1]);  
        printReverse(word, length-1);  
    }  
}
```

- What's going on:
  - We then recursively call the function
  - Sending over two arguments:
    - The string, "computer"
    - And the length, minus 1



# Recursion: General Structure

- Example using Construct 2

- Our function:

```
void printReverse(char word[], int length) {  
    if (length > 0) {  
        printf("%c", word[length-1]);  
        printReverse(word, length-1);  
    }  
}
```

- What's going on:

- After the first recursive call, length is now 7
    - Therefore, word[6] is printed
      - Referring to the "e" in computer
    - Then we recurse (again and again) and finish once length  $\leq 0$



# Brief Interlude: Human Stupidity

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# Recursion – Practice Problem

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## ■ Practice Problem:

- Write a recursive function that:
  - Takes in two non-negative integer parameters
  - Returns the product of these parameters
    - But it does NOT use multiplication to get the answer
  - So if the parameters are 6 and 4
  - The answer would be 24
- How do we do this not actually using multiplication
- What another way of saying  $6*4$ ?
- We are adding 6, 4 times!
- $6*4 = 6 + 6 + 6 + 6$
- So now think of your function...



# Recursion – Practice Problem

- Practice Problem:
  - Solution:

```
// Precondition: Both parameters are
// non-negative integers.
// Postcondition: The product of the two
// parameters is returned.
function Multiply(int first, int second) {
    if ( ( second == 0 ) || ( first = 0 ) )
        return 0;
    else
        return (first + Multiply(first, second-1));
}
```

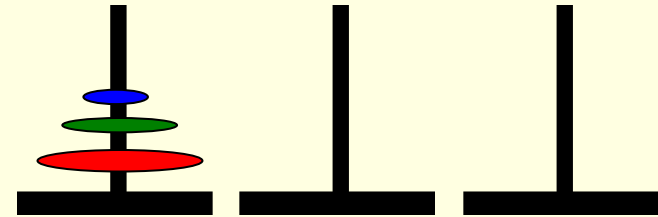


# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Here's the problem:

- There are three vertical poles
- There are 64 disks on tower 1 (left most tower)
  - The disks are arranged with the largest diameter disks at the bottom
- Some monk has the daunting task of moving disks from one tower to another tower
  - Often defined as moving from Tower #1 to Tower #3
    - Tower #2 is just an intermediate pole
  - He **can only move ONE disk at a time**
  - And he MUST follow the rule of **NEVER putting a bigger disk on top of a smaller disk**



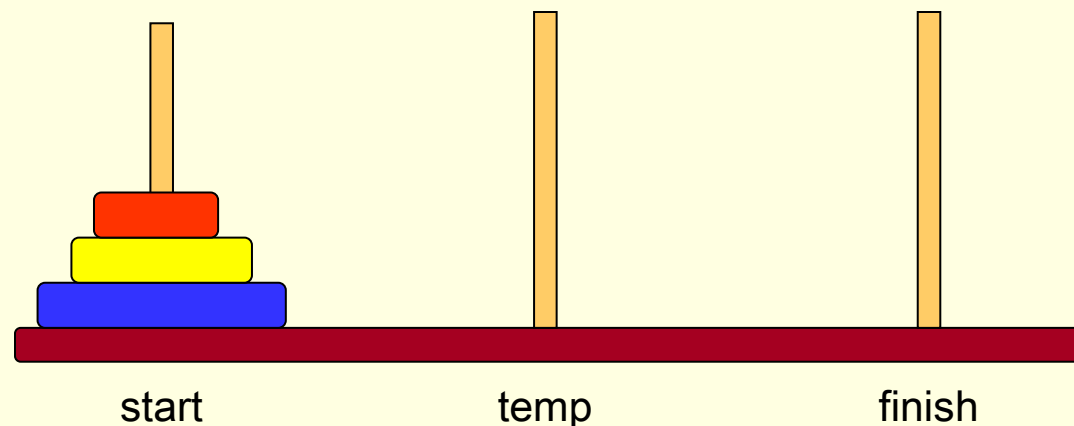


# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Solution:

- We must find a recursive strategy
- Thoughts:
  - Any tower with more than one disk must clearly be moved in pieces
  - If there is just one disk on a pole, then we move it





# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Solution:

- Irrespective of the number of disks, the following steps MUST be carried out:
  - The bottom disk needs to move to the destination tower
    - 1) So step 1 must be to move all disks above the bottom disk to the intermediate tower
    - 2) In step 2, the bottom disk can now be moved to the destination tower
    - 3) In step 3, the disks that were initially above the bottom disk must now be put back on top
      - Of course, at the destination
  
- Let's look at the situation with only 3 disks...





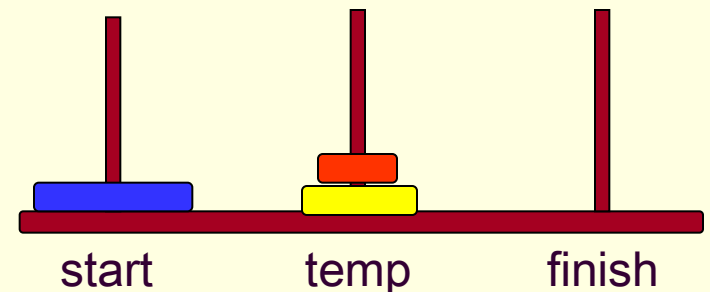
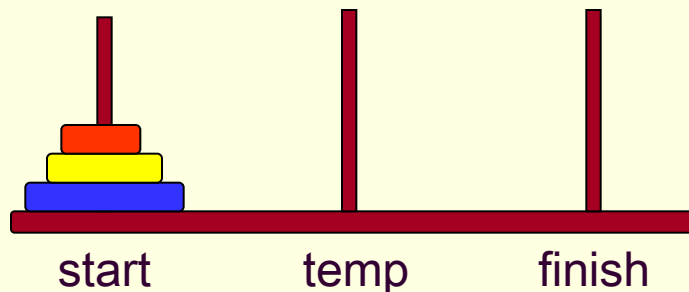
# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Solution:

#### ■ Step 1:

- Move 2 disks from start to temp using finish Tower.
- To understand the recursive routine, let us assume that we know how to solve 2 disk problem, and go for the next step.
  - Meaning, we “know” how to move 2 disks appropriately





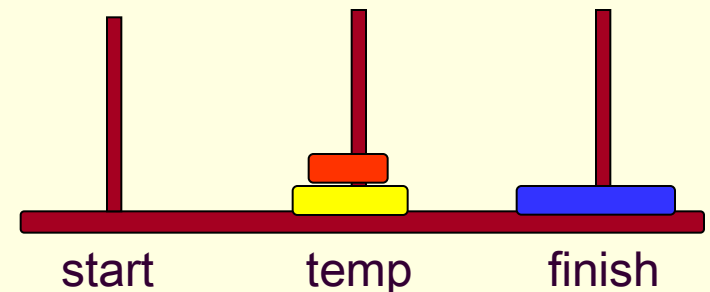
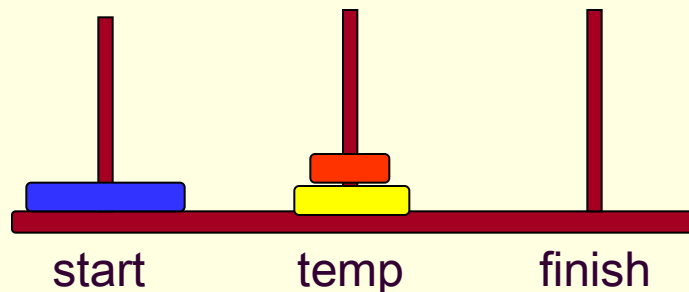
# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Solution:

#### ■ Step 2:

- Move the (remaining) single disk from start to finish
- This does not involve recursion
  - and can be carried out without using temp tower.
- In our program, this is just a print statement
  - Showing what we moved and to where





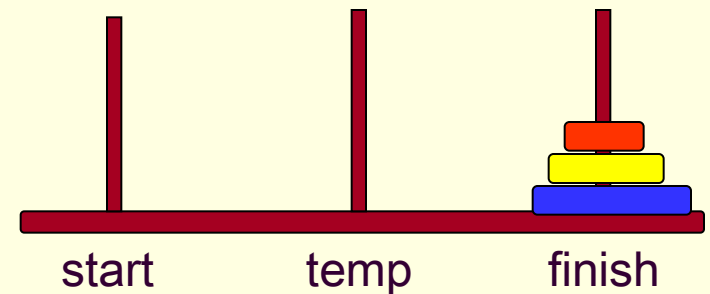
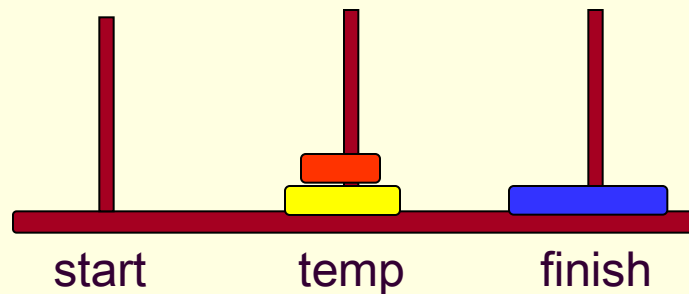
# Recursion – Towers of Hanoi

## ■ Towers of Hanoi:

### ■ Solution:

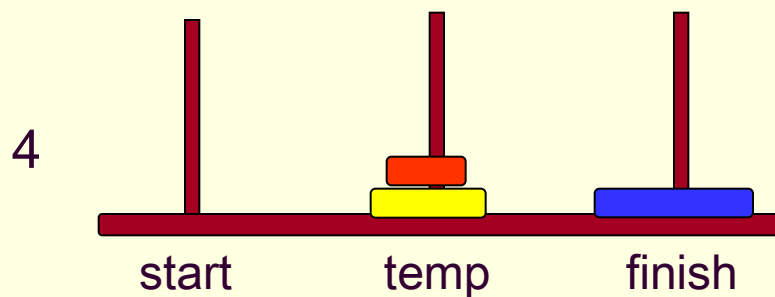
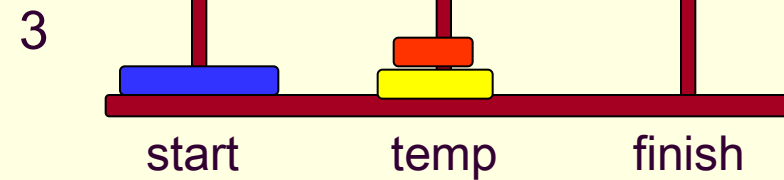
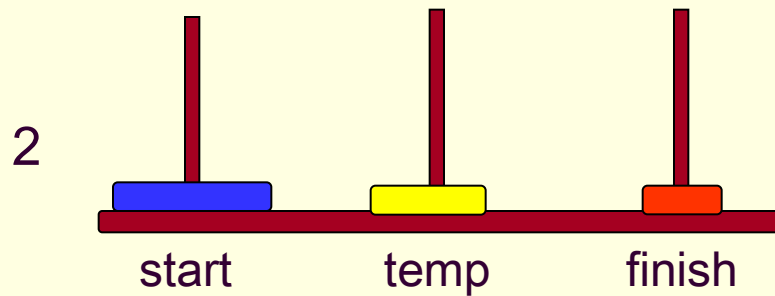
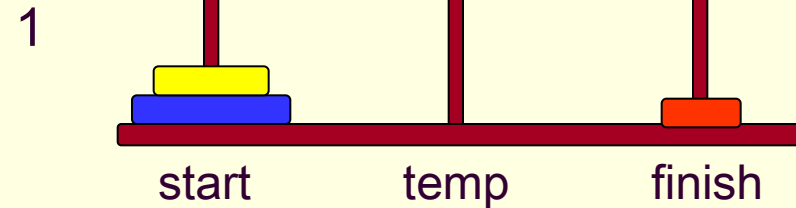
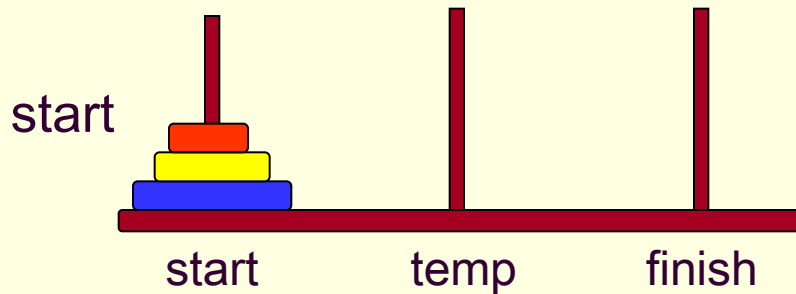
#### ■ Step 3:

- Now we are at the last step of the routine.
- Move the 2 disks from temp tower to finish tower using the start tower
  - This is also done recursively



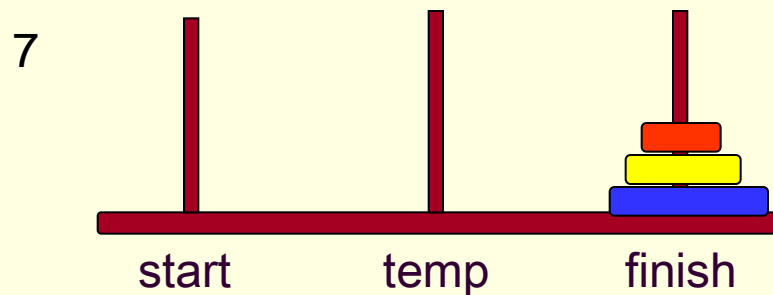
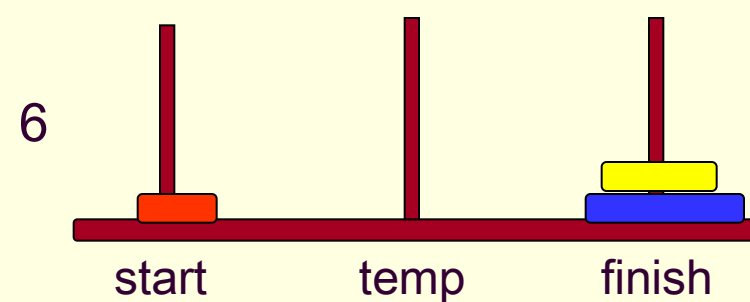
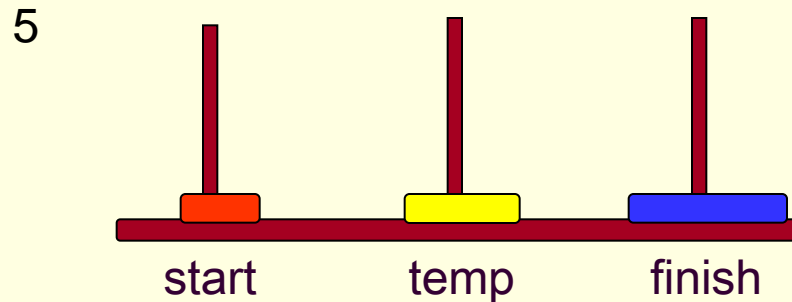


# Recursion – Towers of Hanoi





# Recursion – Towers of Hanoi



- # of steps needed:
  - We had 3 disks requiring seven steps
  - 4 disks would require 15 steps
  - n disks would require  $2^n - 1$  steps
    - HUGE number



# Recursion – Towers of Hanoi

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- Towers of Hanoi:
  - Solution:

```
// Function Prototype
void moveDisks(int n, char start, char finish, char temp);

void main() {
    int disk;
    int moves;
    printf("Enter the # of disks you want to play with:");
    scanf("%d",&disk);
    // Print out the # of moves required
    moves = pow(2,disk)-1;
    printf("\nThe No of moves required is=%d \n",moves);
    // Initiate the recursion
    moveDisks(disk,'A','C','B');
}
```



# Recursion – Towers of Hanoi

- Towers of Hanoi:
  - Solution:

```
void moveDisks(int n, char start, char finish, char temp) {  
    if (n == 1) {  
        printf("Move Disk from %c to %c\n", start, finish);  
    }  
    else {  
        moveDisks(n-1, start, temp, finish);  
        printf("Move Disk from %c to %c\n", start, finish);  
        moveDisks(n-1, temp, finish, start);  
    }  
}
```



# Recursion

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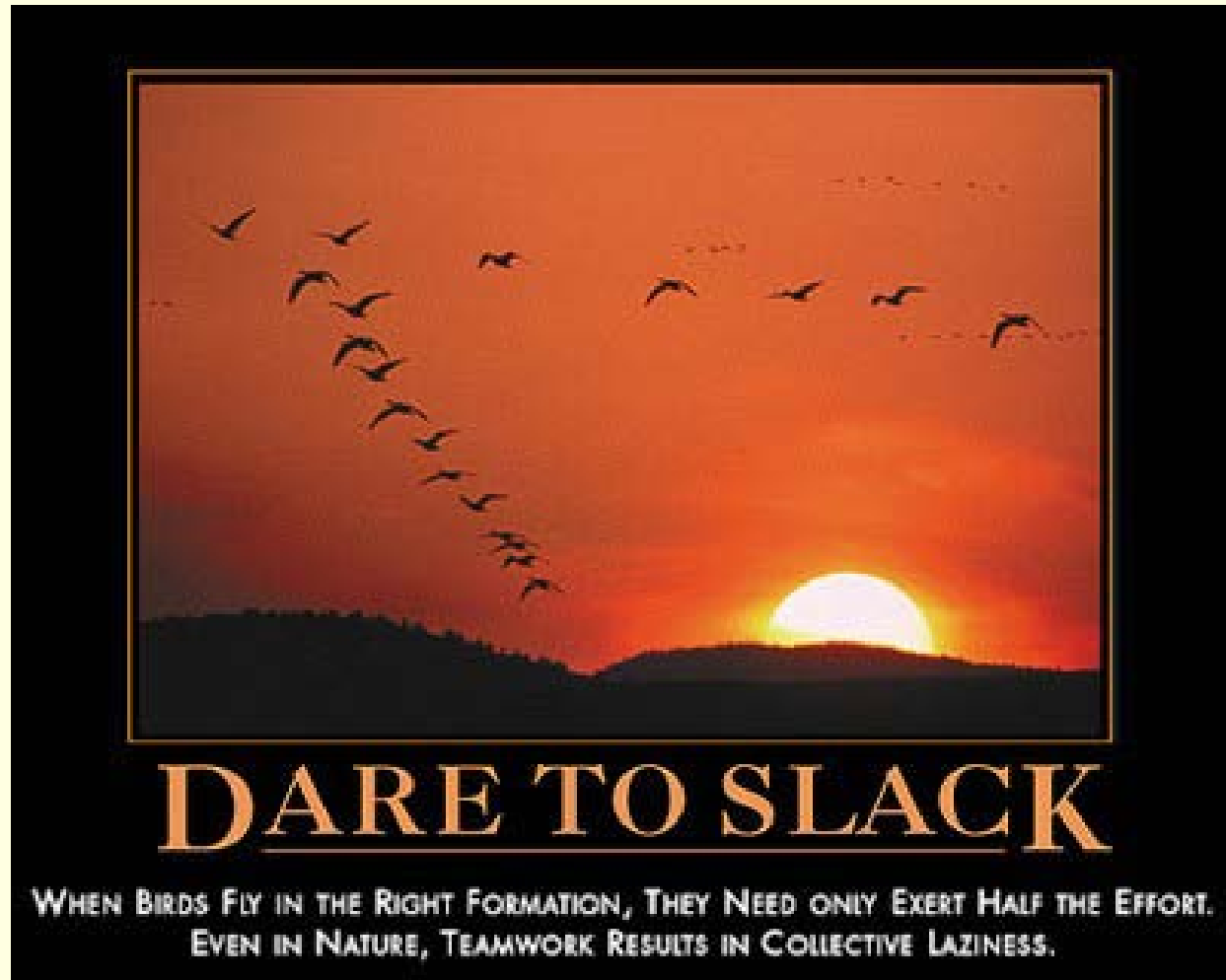
**WASN'T  
THAT  
ENCHANTING!**

(Sorry, wanted a “word of the day”, and this is what I got from the wife!)





# Daily Demotivator



# More Recursion



Computer Science Department  
University of Central Florida

*COP 3502 – Computer Science I*