

# AVL Trees: Deletion



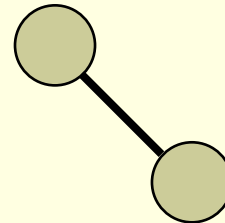
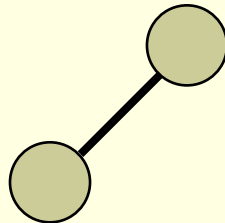
Computer Science Department  
University of Central Florida

*COP 3502 – Computer Science I*



# Insertion Revisited

- AVL Trees: Insertion
  - Let's take another look at insertion into AVL Trees
  - Hopefully this will be a bit easier than previous slides
  - Assuming you only have two nodes in your tree,
  - what are the two possible trees you may have?





# Insertion Revisited

- AVL Trees: Insertion

- Given these two trees, if we want to create an imbalance, where must we insert?



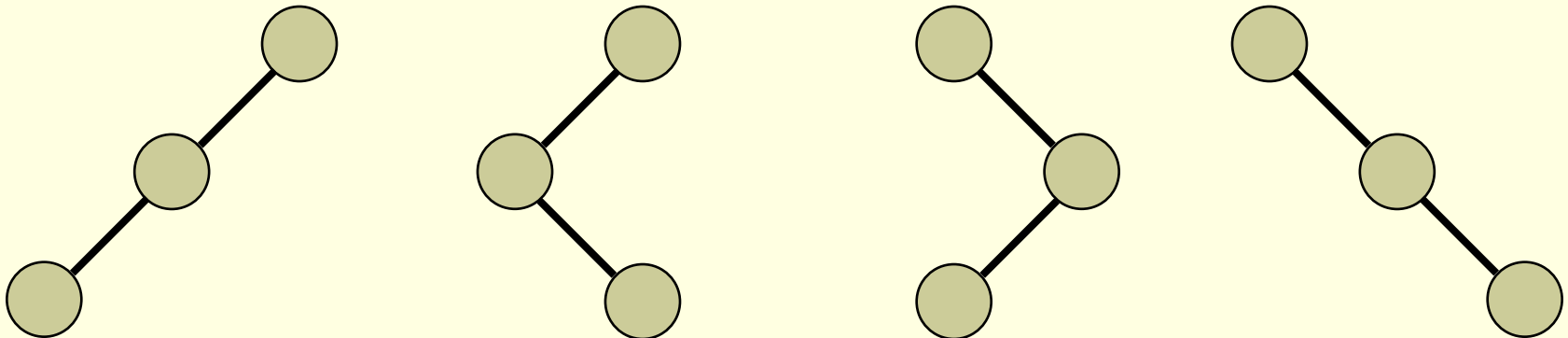
- Clearly, we must insert at the lower of the 2 nodes
- This will create a scenario where the left subtree has a height that is 2 greater than the right subtree
  - Or the opposite for the other tree depicted
- Now, from these two trees, draw all FOUR possible trees that can be created by inserting a new node



# Insertion Revisited

## ■ AVL Trees: Insertion

- Here are all four unbalanced trees that we can make from three nodes:



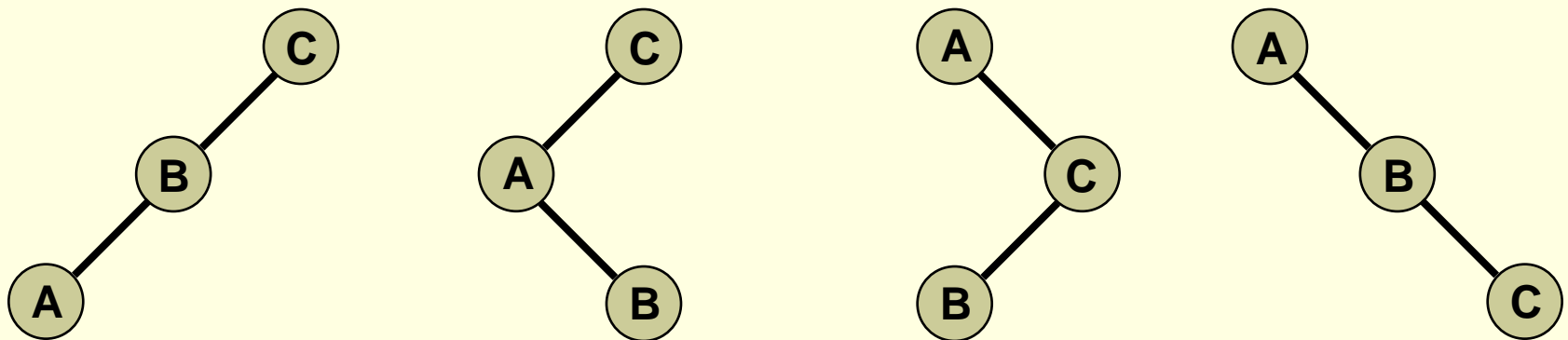
- Now, label these nodes with the labels A, B, and C
  - Where A is the smallest of the three nodes, B is the middle node, and C is the largest.
  - The inorder traversal of each tree should be A, B, C



# Insertion Revisited

## ■ AVL Trees: Insertion

- Here are all four trees with the node labels in their inorder listing:



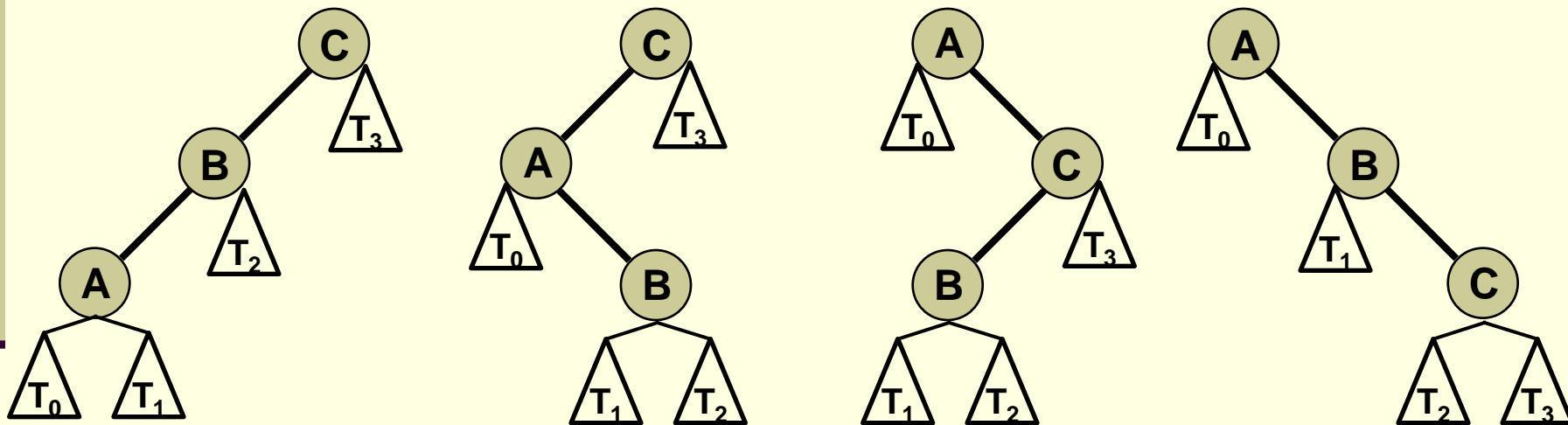
- Any time an imbalance occurs, it is localized to three nodes and their four subtrees
  - These are the four possibilities
  - Now we add in the depiction of the four subtrees of A, B, and C



# Insertion Revisited

- AVL Trees: Insertion

- Here are all four trees with the node labels in their **inorder** listing with subtrees in their **inorder** listing:



- We denote the four subtrees as  $T_0$ ,  $T_1$ ,  $T_2$ , and  $T_3$
- And they are listed in their inorder listing



# Insertion Revisited

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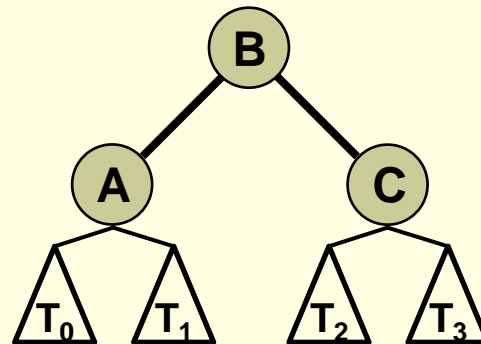
- AVL Trees: Insertion
  - So what is the purpose of all this?
  - We said this method is supposedly MUCH easier than dealing with the various rotations of the tree
  - So we've done all this labeling
    - Finding nodes 'A', 'B', and 'C' and labeling them as such
  - How the heck does this help us???
  
- Here ya go...



# Insertion Revisited

## ■ AVL Trees: Insertion

- Part 1: Once an insertion causes an imbalance, find and label the nodes 'A', 'B', and 'C'
- Part 2: Once the nodes are labeled, no matter what structural imbalance occurred, they can all be fixed the same way:



- Simply restructure those three nodes, and their four respective subtrees, as shown above, and the imbalance will be corrected!

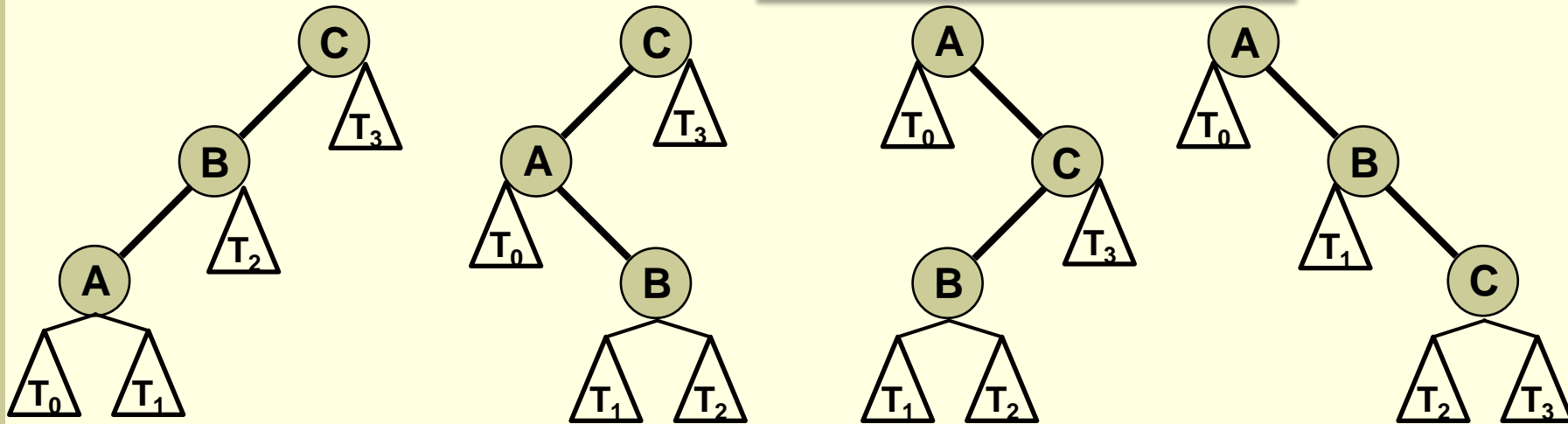




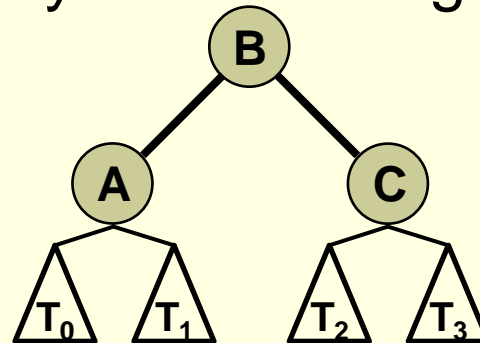
# Insertion Revisited

## ■ AVL Trees: Insertion

All 4 of these trees:



■ Can be fixed by restructuring into this:





# Brief Interlude: FAIL Picture





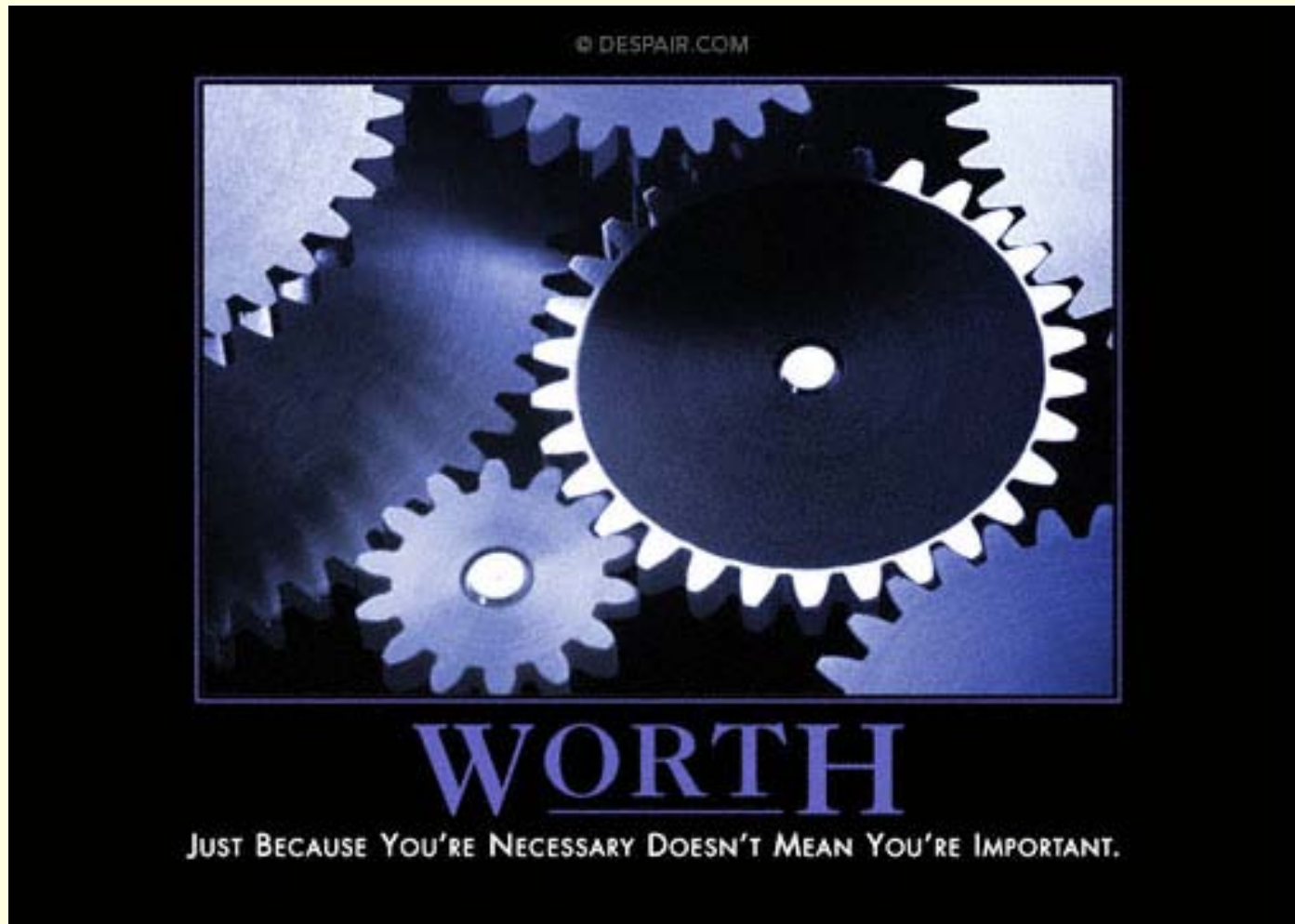
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**WASN'T  
THAT  
MOMENTOUS!**



# Daily Demotivator



# AVL Trees: Deletion



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