

## Fall 2023 COP 3502 Section 1 Final Exam - Part B (100 pts) 12/8/2023 Solution

1) (10 pts) Write a function that takes in a positive integer  $n$ , dynamically allocates a 2D integer array of size  $n$  by  $n$ , fills it with a design of integers in between 1 and  $n$ , where each cell stores how many items are on its “backward diagonal”, and returns a pointer to the array. For example, for  $n = 5$ , the array should be filled as follows:

1	2	3	4	5
2	3	4	5	4
3	4	5	4	3
4	5	4	3	2
5	4	3	2	1

Fill in the function prototype below. You may use any functions from the math library as needed.

```
int** diagcount(int n) {  
  
    int** sq = calloc(n, sizeof(int*));  
    for (int i=0; i<n; i++)  
        sq[i] = calloc(n, sizeof(int));  
  
    for (int d=0; d<2*n-1; d++) {  
        int start = d < n ? 0 : d-n+1;  
        int draw = d < n ? d+1 : 2*n-1-d;  
        for (int i=start; d-i>=0 && i<n; i++)  
            sq[i][d-i] = draw;  
    }  
  
    return sq;  
}
```

```
// Alternate solution using 2 loops.  
int** diagcount_alt(int n) {  
  
    int** sq = calloc(n, sizeof(int*));  
    for (int i=0; i<n; i++)  
        sq[i] = calloc(n, sizeof(int));  
  
    for (int d=0; d<n; d++)  
        for (int i=0; i<=d; i++)  
            sq[i][d-i] = d+1;  
  
    for (int d=n; d<2*n-1; d++)  
        for (int i=d-n+1; i<n; i++)  
            sq[i][d-i] = 2*n-1-d;  
  
    return sq;  
}
```

**Grading: 3 pts for mem alloc, 2 pts for loop structure, 4 pts for logic to fill in grid, 1 pt for return, many other solutions possible.**

2) (10 pts) Write a function that takes in a pointer to a linked list and uses it to create a new linked list of the same size storing the values in the input linked list in the **reverse order** of the original list, returning a pointer to the front of this newly created list. **Write your function iteratively, repeatedly creating a new node and inserting it into the front of the new list you are creating.** (For example, if the input list to the function stores 2, 6, 5, 7, then your function will create a new list, inserting 2 to the front, then 6 to the front, then 5 to the front and finally 7 to the front to generate a list in the order 7, 5, 6, 2, returning a pointer to the node storing 7.)

```
typedef struct node {
    int data;
    struct node* next;
} node;

node* newrevlist(node* front) {

    node* cur = NULL;
    while (front != NULL) {
        node* tmp = malloc(sizeof(node));
        tmp->data = front->data;
        tmp->next = cur;
        cur = tmp;
        front = front->next;
    }

    return cur;
}
```

**Grading: works in null case – 1 pt**  
**Works in one node case – 2 pts**  
**Loops through input list – 2 pts**  
**Mallocs new nodes in loop – 1 pt**  
**Copies data from input list into - 1 pt**  
**List building works – 2 pts**  
**Return – 1 pt**

3) (6 pts) Evaluate the value of the following postfix expressions. No need to show the stack, you'll only be graded on the final answer.

(a) 18 3 / 3 2 + 4 5 - + \* 24

(b) 6 5 4 3 2 1 + \* + \* + 71

(c) 8 2 5 \* + 9 3 - / 7 \* 21

**Grading: 2 pts for each, all or nothing**

4) (4 pts) What does the function  $f(13, 12)$  return, where  $f$  is defined below?

```
int f(int a, int b) {
    if (b == 0) return 0;
    int res = a * (b & 1);
    return res + (f(a, b>>1) << 1);
}
```

156

Note: What the code is doing is just multiplying  $a$  by  $b$ . To see this, notice that each time, if the least significant bit is 1, then  $a$  is added, otherwise it's not. What the recursive call does is the same operation as before, but chops off the least significant bit of  $b$ . Then, this result is left shifted and added to the original. This ensures that the "contribution" of each copy of  $a$  is multiplied by the appropriate power of 2, based on where it shows up in  $b$ 's binary representation. Essentially, since  $12 = 1100$  in binary, this code will have the effect of adding  $8 \times 13 + 4 \times 13 = 156$ .

**Grading: 4 or 0, unfortunately it's all or nothing**

5) (8 pts) Consider the problem of earning money until you get at least a target amount of dollars. You have a job where if on a particular day you earn  $d$  dollars, then the next day you will earn  $d + x$  dollars, where  $x$  is a constant. Write a **recursive function** that takes in integers  $target$ ,  $d$  and  $x$ , where  $target$  is how much is left to earn,  $d$  is the pay for today, and  $x$  is the increase factor between days, and returns the minimum number of days (whole number) of work needed to earn the required amount of money.

```
int earnmoney(int target, int d, int x) {
    if (target <= 0) return 0;
    return 1 + earnmoney(target-d, d+x, x);
}
```

**Grading: 2 pts if for base case, 1 pt return in base case**

**1 pt add 1, 1 pt each parameter (only get credit if rec call for these), 1 pt return**

6) (10 pts) Please give **an exact solution** (a function that  $T(n)$  satisfies in terms of  $n$ ) to the following recurrence relation using the iteration technique:

$$T(n) = 2T(n-1) + 2^n, \text{ for } n > 0$$
$$T(0) = 4$$

$$T(n) = 2T(n-1) + 2^n$$

$$T(n) = 2(2T(n-2) + 2^{n-1}) + 2^n$$

$$T(n) = 4T(n-2) + 2^n + 2^n$$

$$T(n) = 4T(n-2) + 2(2^n)$$

$$T(n) = 4(2T(n-3) + 2^{n-3}) + 2(2^n)$$

$$T(n) = 8T(n-3) + 2^n + 2(2^n)$$

$$T(n) = 8T(n-3) + 3(2^n)$$

After  $k$  iterations, we have

$$T(n) = 2^k T(n-k) + k(2^n)$$

Plug in  $k = n$ :

$$T(n) = 2^n T(0) + n(2^n) = 4 \times 2^n + n(2^n) = (n+4)2^n$$

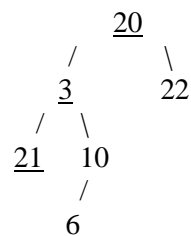
**Grading: 1 pt first iteration, 2 pts second iteration, 2 pts third iteration, 2 pt general, 3 pts to finish**

7) (6 pts) Show the contents of the following array right after a partition is performed on it using index 0 as the partition element. Please use the partition algorithm shown in class.

index	0	1	2	3	4	5	6	7	8	9	10	11
orig array	13	16	9	14	17	23	5	6	22	1	8	22
after part	5	8	9	1	6	13	23	17	22	14	16	22

**Grading: floor(correct/2)**

8) (10 pts) Consider the following problem: given a binary tree, find the maximum possible value that can be obtained by including the values stored in each node on a single path from the root node to any leaf node. For example, in the tree below, the nodes on the path that provides the maximum sum are underlined.



Write a function that takes in a pointer to the root of a binary tree and returns this maximal sum. Your function **must be recursive**.

```

typedef struct treenode {
    int data;
    struct treenode* left;
    struct treenode* right;
} treenode;

int summaxpath(treenode* root) {

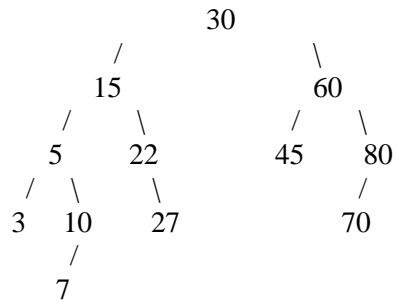
    if (root == NULL) return 0;

    int left = summaxpath(root->left);
    int right = summaxpath(root->right);

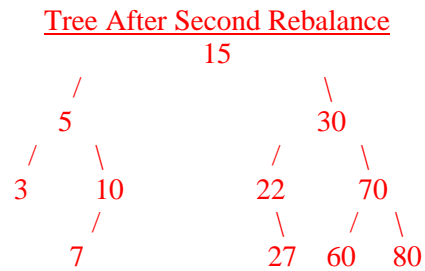
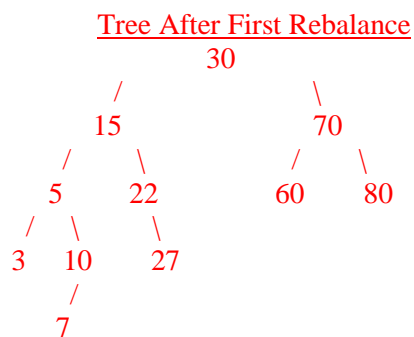
    if (left > right)
        return left + root->data;
    return right + root->data;
}
  
```

**Grading: 2 pts root NULL case,  
 2 pts each rec call  
 2 pts if or ? for larger side  
 1 pt to add root->data to left call  
 1 pt to add root->data to right call**

9) (8 pts) Show the result of deleting the value 45 from the AVL tree shown below:



Since there are two tree rebalances, please draw a box around the state of the tree after the first rebalance, and then again after the second/final rebalance. **So you should have two boxes around two trees.**



**Grading: First Pic → 1 pt remove 45, 3 pts for positions of 60, 70 and 80. (0 if not BST)  
 Second Pic → 1 pts 15 at root, 1 pt 5 left, 1 pt 30 right, 1 pts attach rest (0 if not BST)**

10) (10 pts) Consider the task of printing out (in lowercase letters) the **last word** alphabetically stored in a trie. (You may assume if a link exists in a trie, that some word exists further down that link.) Write a **void recursive** function to accomplish this task. (Note: Different recursive calls will print different letters in this word.)

```
typedef struct trienode {
    int isWord;
    struct trienode next[26];
} trienode;

void printlast(trienode* root) {

    if (root == NULL) return;

    for (int i=25; i>=0; i--) {
        if (root->next[i] != NULL) {
            printf("%c", (char)('a'+i));
            printlast(root->next[i]);
            break;
        }
    }
}
```

**Grading: 2 pts base case**

**3 pts loop backwards 25 to 0**

**1 pts check for NULL**

**2 pts print letter**

**1 pt rec call**

**1 pt break or return here or some mechanism to not print too much**

11) (5 pts) Convert 917 in base 10 to base 5.

```
5 | 917
5 | 183 R 2
5 | 36 R 3
5 | 7 R 1
5 | 1 R 2
5 | 0 R 1
```

**12132**

**Grading: 1 pt each “digit”, if correct digits but wrong order do 4/5. Must do division, remainder process to get points.**

12) (10 pts) In the game of NIM, a player is given several piles of stones and must select one of the piles to take 1 or more stones from. A player (currently moving) can win if the bitwise XOR of the number of stones in all the piles is any non-zero number. Any move the player makes to change this bitwise XOR to 0 (after her move) is an optimal move. In particular, let  $x$  = the bitwise XOR of the number of stones in all the piles and let  $\text{piles}[i]$  be the number of stones in  $\text{piles}[i]$ . As long as the bitwise XOR between  $x$  and  $\text{piles}[i]$  (call this xor computation  $y$ ) is less than  $\text{piles}[i]$ , then an optimal move would be take  $y$  stones from pile  $i$ . Write a function that takes in an array,  $\text{piles}$ , and integer  $n$ , representing the number of piles, and returns an integer array of size 2, where index 0 stores which pile (0 based) to remove stones from, and index 1 stores how many stones to take for an optimal move. If no move exists to return the state of the piles to have a bitwise XOR of 0, then return NULL. Some of the code has been given.

```
int* getwinmove(int* piles, int n) {

    int xor = 0;
    /* Fill in code here. */
    for (int i=0; i<n; i++)           // Grading: 1 pt
        xor ^= piles[i];             // Grading: 2 pts

    if (xor == 0) return NULL;

    int* res = calloc(2, sizeof(int));

    /* Fill in code here */
    for (int i=0; i<n; i++) {         // Grading: 1 pt
        int leave = piles[i]^xor;     // Grading: 2 pts
        if (leave < piles[i]) {      // Grading: 2 pts
            res[0] = i;               // Grading: 1 pt
            res[1] = piles[i]-leave;  // Grading: 1 pt
        }
    }

    return res;
}
```

13) (3 pts) George Boole passed away on December 8<sup>th</sup>, 1864, exactly 159 years ago. What type of expression, found in if statements, is named after him?

**Boolean (Give to All)**