

1) Papers - come to office after or office hrs

2) El Commit & Digital Signature Scheme - instead

Alice  $M$  to Bob

she wants to prove she wrote it.

Alice has to have her own El Commit keys for (regular crypto scheme).

$$q = \text{prime \#}$$

$$\alpha = \text{primitive root}$$

$$X_A = \text{Alice's secret \#}$$

$$Y_A = \alpha^{X_A} \text{ mod } q \text{ (Alice's public key)}$$

To sign she wants to prove that used her private key w/o divulging any info about it.

Alice  $\rightarrow M \rightarrow$  Bob (encryption is independent / separate of signature)

How ALICE SIGNS  $M$ ,  $(S1, S2)$

1) Calculate  $m = H(M)$  using some agreed upon hash function,  $m$  is fixed size output

2) Choose random,  $K$ ,  $\gcd(K, q-1) = 1$ . (If this  $\gcd = 1$ , then  $\alpha^K$  is also a primitive root.)

$$3) S1 = \alpha^K \text{ mod } q$$

4) Compute  $k^{-1} \pmod{q-1}$ , can only exist iff  $\gcd(k, q-1) = 1$ .

$$5) S_2 = k^{-1} (m - X_a \cdot S_1) \pmod{q-1}$$

Bob received  $m, (S_1, S_2)$ , wants to verify Alice's signature.

1) Bob calculates  $m = H(M)$

2) To verify Bob calculates

$$V_1 = \alpha^m \pmod{q}$$

$$V_2 = (Y_a)^{S_1} \times S_2 \pmod{q}$$

Message is verified as a valid signature ONLY IF  $V_1 = V_2$ ,

$$V_2 = (Y_a)^{S_1} \times S_2 \pmod{q}$$

$$= (\alpha^{X_a \cdot S_1}) \times \alpha^{k \cdot S_2} \pmod{q}$$

$$= \alpha^{X_a \cdot S_1 + k \cdot S_2} \pmod{q}$$

find this mod  $q-1$

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$$\equiv \alpha^{m \pmod{q-1}} \pmod{q}$$

$$\equiv \alpha^m \pmod{q} \equiv V_1 \quad \checkmark$$

$V_1$  was this

Via Fermat's Theorem

$$2^B \equiv 2^{\frac{B \bmod q-1}{}} \pmod{q}$$

$$2^{B_1} \equiv 2^{B_2} \pmod{q} \quad \text{iff}$$

if  $B_1 \equiv B_2 \pmod{q-1}$ , then

$$2^{B_1} \equiv 2^{B_2} \pmod{q}$$

$$X_a \cdot S_1 + k \cdot S_2 = X_a \cdot d^k + \underbrace{k \cdot k^{-1} (m - X_a \cdot S_1)}_{\pmod{q-1}}$$

$$\equiv X_a \cdot d^k + m - X_a \cdot d^k$$

$$\equiv \boxed{m \pmod{q-1}}$$

$$y^2 \equiv x^3 + ax + b \pmod{p}$$

Find  $(x, y)$  that works

$$C \equiv x^3 + ax + b$$

→ Compute this.

$$y^2 \equiv c \pmod{p}$$

Has 2 solutions iff

$$c^{\frac{p-1}{2}} \equiv 1 \pmod{p}$$

Has 1 solution iff

$$c^{\frac{p-1}{2}} \equiv 0 \pmod{p}$$

else

has 0 solutions

→ Ans  $c^{\frac{p+1}{4}} \pmod{p}$

Reason

$$\left(c^{\frac{p+1}{4}}\right)^2 = c^{\frac{p+1}{2}} = c^{\frac{p-1}{2} + 1}$$

$$\equiv c^{\frac{p-1}{2}} \times c \pmod{p}$$

$$\equiv 1 \times \underline{\underline{c}} \pmod{p}$$