

Discrete Log -

$$y = b^x \pmod{p}$$

Given  $b, y$  and  $p$  find  $x$ .

table seq  $b^1, b^2, b^3 \dots b^{p-1}, b^{p-1} \equiv 1 \pmod{p}$   
and after that it repeats, so all rows have  
a cycle length that is either  $p-1$ , or divides  
evenly into  $p-1$ .

Special bases are ones where the row has each  
value from 1 to  $p-1$  exactly once (cycle len =  
 $p-1$ .)

primitive roots of generators mod  $p$  (prime)

2 is a generator mod 29

$2^1, 2^2, 2^3 \dots 2^{28} \pmod{29}$  is a seemingly  
random permutation of 1 to 28.

\* Why this will be useful

\* Patterns in chart

\* Algorithm to solve Disc Log a little bit  
faster.

for encryption we like processes that are "easy"  
calculate forwards, hard to "invert" or undo.

~~Discrete~~ Discrete Log is exactly this!

# ONE-WAY FUNCTION!

If a generator,  $g$ , exists for a prime  $p$ , we can prove that there are exactly  $\phi(p-1)$  generators.

Given:  $g^1, g^2, g^3, \dots, g^{p-1}$  are all unique mod  $p$   
exactly  $\phi(p-1)$  integers in the range 1 to  $p-1$  do NOT share a common factor with  $p-1$ .

Let  $a$  be an arbitrary int  $1 \leq a \leq p-1$  s.t.  $\gcd(a, p-1) = 1$ .

Prove  $g^a$  is also a generator.

$$g^a, g^{2a}, g^{3a}, g^{4a}, \dots, g^{(p-1)a}$$

We know that  $g^x \equiv g^{x \pmod{p-1}} \pmod{p}$

$$p=7 \quad g^{20} \equiv g^{2 \pmod{6}} \pmod{7}$$

because  $20 \equiv 2 \pmod{6}$

$$g^{20} = (g^6)^3 g^2 \equiv 1^3 \cdot g^2 \equiv g^2 \pmod{p}$$

What is  $24^{20} \pmod{7}$ , since  $24 \equiv 3 \pmod{7}$   
 $\equiv 3^{20} \pmod{7}$ , since  $20 \equiv 2 \pmod{6}$   
 $\equiv 3^2 \pmod{7}$   $\uparrow$   
 $\phi(6)$

Need to Prove exponents:  $a, 2a, 3a, \dots, (p-1)a$  are all distinct mod  $p-1$ .

$$\gcd(a, p-1) = 1.$$

Assume the contrary,

$$a_i \equiv a_j \pmod{p-1}, \quad i \neq j, \quad 1 \leq i, j \leq p-1$$

$$a_i - a_j \equiv 0 \pmod{p-1}$$

$$a(i-j) \equiv 0 \pmod{p-1}$$

$$\rightarrow (p-1) \mid a(i-j), \text{ but } \text{since } \gcd(a, p-1) = 1$$

$$\rightarrow (p-1) \mid (i-j), \text{ contradicts fact that } 1 \leq |i-j| \leq p-2$$

CONCLUSION is # generators has to be equal to  $\phi(p-1)$ . (last part if  $\gcd(a, p-1) > 1$ , then  $g^a$  isn't generator.)

$$p = 29 \quad p-1 = 28 \quad \phi(28) = (2^2 - 2^1)(7-1)$$

$$= 2 \times 6 = 12 \text{ generators}$$

How many bases have

$$\text{cycle length } 14 \rightarrow \phi(14) = (2-1)(7-1) = 6$$

$$\sum \phi(d) = p-1$$

$d \in \text{Divisor}(p-1)$

$p \in \text{Prime}$

$$\phi(4) = (2^2 - 2) = 2$$

$$\phi(7) = (7-1) = 6$$

$$\phi(2) = (2-1) = 1$$

$$\phi(1) = 1$$

# Discrete Log in $\sqrt{p}$ time

Req DiscLog

```

int disclog(int b, int y, int p) {
    int ans = 1;
    for (int i = 0; i < p; i++) {
        if (ans == y) return i;
        ans = (ans * b) % p;
    }
}

```

$O(p)$

$p = 29$

base  $\sqrt{p}$  ish

$$g^a = g^{6x+y}$$

$$= g^{6x-y}$$

$0 \leq x < 5$   
 $0 \leq y \leq 5$

can do + or - smaller term!

Calculate  $g^6, g^{12}, g^{18}, g^{24}, g^{30}$  ] run time is  $O(\sqrt{p})$

$g^6 \rightarrow 6$

$g^{12} \rightarrow 7$

$g^{18} \rightarrow 13$

$g^{24} \rightarrow 20$

$g^{30} \rightarrow 4$

store answers in a table

$2^{16} \equiv 25 \pmod{29}$

$2^{6x} \cdot 2^y \equiv 25 \pmod{29}$

$2^{6x-24} \equiv 1 \pmod{29}$

$$2^{6x} \cdot 2^y \equiv 25 \pmod{29}$$

$$\cancel{2^{-y}} \times 2^{6x} \times \cancel{2^y} \equiv 25 \times 2^{-y} \pmod{29}$$

$$2^{6x} \equiv 25 \times 2^{-y} \pmod{29}$$

MAP

$$6 \rightarrow 6$$

$$7 \rightarrow 12$$

$$13 \rightarrow 18$$

$$20 \rightarrow 24$$

$$4 \rightarrow 30$$

GOAL

$$2^y \cdot 2^{6x} \cdot 2^{-y} \equiv 2^y 25 \pmod{29}$$

$$2^{6x} \equiv 2^y \cdot 25 \pmod{29}$$

$$y=0 \rightarrow 25$$

$$y=1 \rightarrow 50 \equiv 21 \pmod{29}$$

$$y=2 \rightarrow 42 \equiv 13$$

$$18 \equiv 2^2 \cdot 2^5 \pmod{29}$$

$$18 - 2 = 16$$