

1. Mix Col Program
2. Mix Cols by Hand
3. Key Schedule

$a = 17$ $b = 6$ $16b + b$
 $r \rightarrow 0110$ $res = 000$ $a = 10001 =$
 $cur = 0110$

when lsb is on $res = res \wedge cur$, add in this many copies of b

$cur = \text{leftshift}(cur) \rightarrow$ multiplies cur by 2^x essentially

row \ col	col 0	col 1	col 2	col 3
row 0	02	03	01	01
row 1	01	02	03	01
row 2	01	01	02	03
row 3	03	01	01	02

AES Poly

$$x^8 + x^4 + x^3 + x + 1$$

$02 \times D3 = BD$
 $03 \times 7E = 82$
 $01 \times B5 = B5$
 $01 \times 29 = 29$

 A3

$02 \times (11010011) = 10100110$
 $\oplus 1101$

 10111101
 B D

$03 \times (01111110)$
 $= (02 + 01) \times (01111110) = \oplus 11111100 \rightarrow 02 \times 7E$

 1000010
 8 2

AES Key Schedule

128 bit key input

creates 10 round keys

initial XOR is w/ the original key.

Notes for key expansion

$w[0..3] = \text{orig key}$

$w[4..7] = R1 \text{ key}$

$w[8..11] = R2 \text{ key}$

⋮

$w[40..43] = R10 \text{ key}$

Key Exp(byte key[16], word w[44])

~~tmp =~~

$w[0..3] = \text{key}[0..15]$ // copy key into R0 key

for ($i=4; i < 44; i++$)

temp = $w[i-1]$

if ($(i \% 4 == 0)$)

temp = SubWord(RotWord(temp) \oplus (Rcon[$i/4$] 000000))

$w[i] = w[i-4] \oplus \text{temp}$

R1 key ends in ... 49 36 B2 E7 $w[7]$ $w[10]$

~~R2 key~~

R2 key is

AF 01 68 74 42 36 B8 89 EF 29 36 24

find $w[11]$ →

	49	36	B2	E7
\oplus	EF	29	36	24
	A6	1F	84	C3

