

# DES 9/24/25

Early 1970s

Competition to find a standard for symmetric key encryption.

↳ 2 people communicating share a private key.

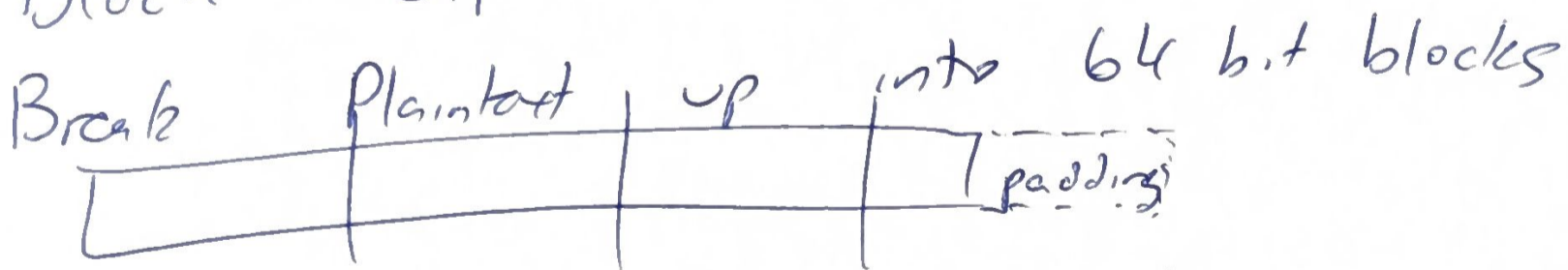
Winning System: Lucifer used internally @ IBM  
Horst Feistel

↳ NSA took the original algorithm + modified it.  
adopted ~ 1977ish

was soon standard from 70s to 1999

Best Paper Found a chosen plaintext-attack of size  $2^{48}$  to find a key of 56 bits.

## Block Cipher

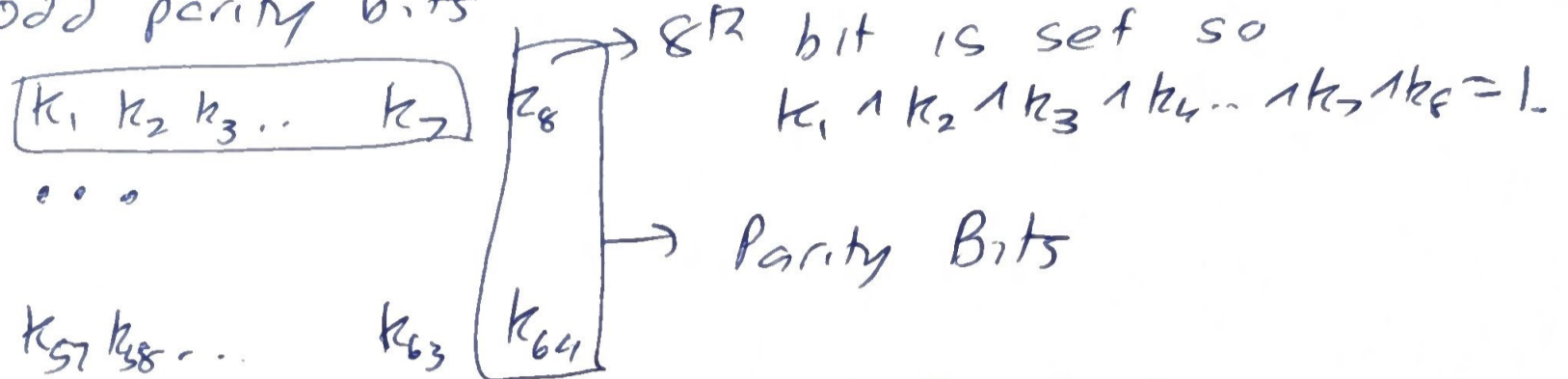


Several "Modes" of running block cipher.

Only talk about encrypting 1 block with a key.

Input Plaintext is 64 bits

Key is 56 bits but is transmitted with 8 odd parity bits



High Level Published  
 $= [L_0 R_0]$   $L_0$  is left 32 bits  
 $R_0$  is ~~left~~ right 32 bits

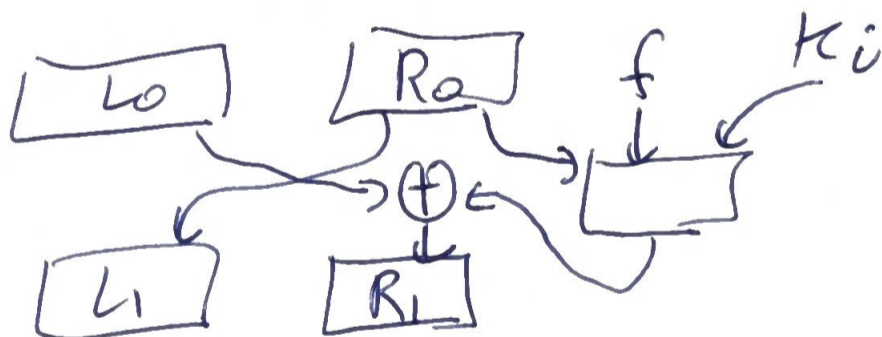
1.  $X \rightarrow IP(X)$ ,  $IP$  is a permutation of bits

2. for  $i=1$  to 16 #Rounds

$$L_i = R_{i-1}$$

$$R_i = L_{i-1} \oplus f(R_{i-1}, K_i)$$

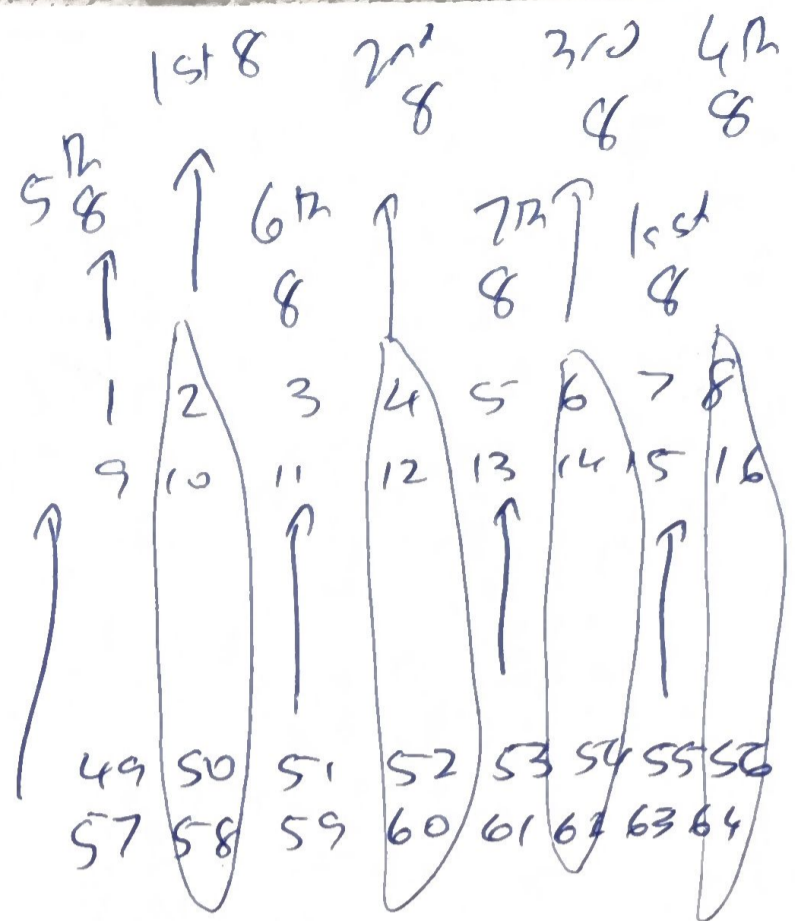
48 bit Round key generated from 56 bit key



3.  $IP^{-1}(R_{16} L_{16})$   
 $16 \quad 16$

# IP Patterns

58	50	...	2
60	52	...	4
62	54	...	6
64	56	...	8
57	49	...	1
59	51	...	3
61	53	...	5
63	55	...	7



Definition of Inverse Function

$$f^{-1}(f(x)) = x$$

Decryption
Encrypting

$$f(R_{i-1}, K_i) \rightarrow 48 \text{ bits}$$

1.  $E(R_{i-1}) \parallel (\text{Expansion Matrix})$

2.  $K_i \oplus E(R_{i-1})$       6 bits   6 bits      6 bits

3. Let  $K_i \oplus E(R_{i-1}) = b_1 b_2 b_3 b_4 \dots b_8$   
 Break these 48 bits into 8 6-bit blocks

Output  $c_i = S_i(b_i)$  for  $i=1$  to 8

$\uparrow$  4 bits       $\uparrow$  6 bits  
 Only non-linear part of alg.