

CIS 3362 9/27/24

DES cont Begin

$$P_0 = L_0 R_0 = IP(P)$$

Rounds

for $i = 1$ to 16

$$L_i = R_{i-1}$$

$$R_i = L_{i-1} \oplus f(R_{i-1}, K_i)$$

End

Round 16 $L_{16} R_{16}$

$$\text{Cipher} = IP^{-1}(R_{16} L_{16})$$

$f(\text{Input } A \text{ (text), Input } J \text{ (key round)})$

1. $E(A) \rightarrow 32 \text{ bits to } 48 \text{ bits}$

2. $B = E(A) \oplus J$ (xor w/ round key)

3. $B = B_1 B_2 \dots B_8$ (split into 6 bit blocks)

4. for $i = 1$ to 8

$$C_i = S_i(B_i)$$

4 bits = 6 bits

$$B_i = b_1 b_2 b_3 b_4 b_5 b_6$$

row = $b_1 b_6$

col = $b_2 b_3 b_4 b_5$

$S_i(B_i)$ is the entry

in row $b_1 b_6$ col $b_2 b_3 b_4 b_5$

5. return $P(C)$.

$$S_1(\underline{101011}) - \begin{array}{l} \text{row} = 11 \rightarrow 3 \\ \text{col} = 0101 \rightarrow 5 \end{array} S_1[3][5] = 9(1001)$$

$$S_4(\underline{011101}) - \begin{array}{l} \text{row} = 01 \rightarrow 1 \\ \text{col} = 1110 \rightarrow 14 \end{array} S_4[1][14] = 14(1110)$$

$$S_6(\underline{000110}) - \begin{array}{l} \text{row} = 00 \rightarrow 0 \\ \text{col} = 0011 \rightarrow 3 \end{array} S_6[0][3] = 15(1111)$$

$$S_8(\underline{101100}) - \begin{array}{l} \text{row} = 10 \rightarrow 2 \\ \text{col} = 0110 \rightarrow 6 \end{array} S_8[2][6] = 14(1110)$$

S box criteria

1. Each row is a permutation of 0 to 15.
2. No s-box is a linear or affine function of its inputs.
3. Changing 1 bit in input forces at least 2 output bits to change.
4. For all x , $S(x)$ and $S(x \oplus 001100)$ differ in at least 2 bits
5. $S(x) \neq S(x \oplus 11e f 00)$ for all e, f .
6. If you fix a single input bit and analyze a fixed output bit there should be "no worse than" a 13-19 split btw 0s + 1s.

Verify #5 for S_2 , $x = 101100$

$x \oplus 11100$

$$S_2(011100) : \begin{matrix} r=0 \\ c=4 \end{matrix} \rightarrow \begin{matrix} 5 \\ 0101 \end{matrix}$$

$$S_2(011000) : \begin{matrix} r=0 \\ c=12 \end{matrix} \rightarrow \begin{matrix} 12 \\ 1100 \end{matrix}$$

$$S_2(010100) : \begin{matrix} r=0 \\ c=10 \end{matrix} \rightarrow \begin{matrix} 2 \\ 0010 \end{matrix}$$

$$S_2(\underline{010000}) \rightarrow \begin{matrix} r=0 \\ c=8 \end{matrix} \rightarrow \begin{matrix} 3 \\ 0011 \end{matrix}$$

flipped 4
pos
same

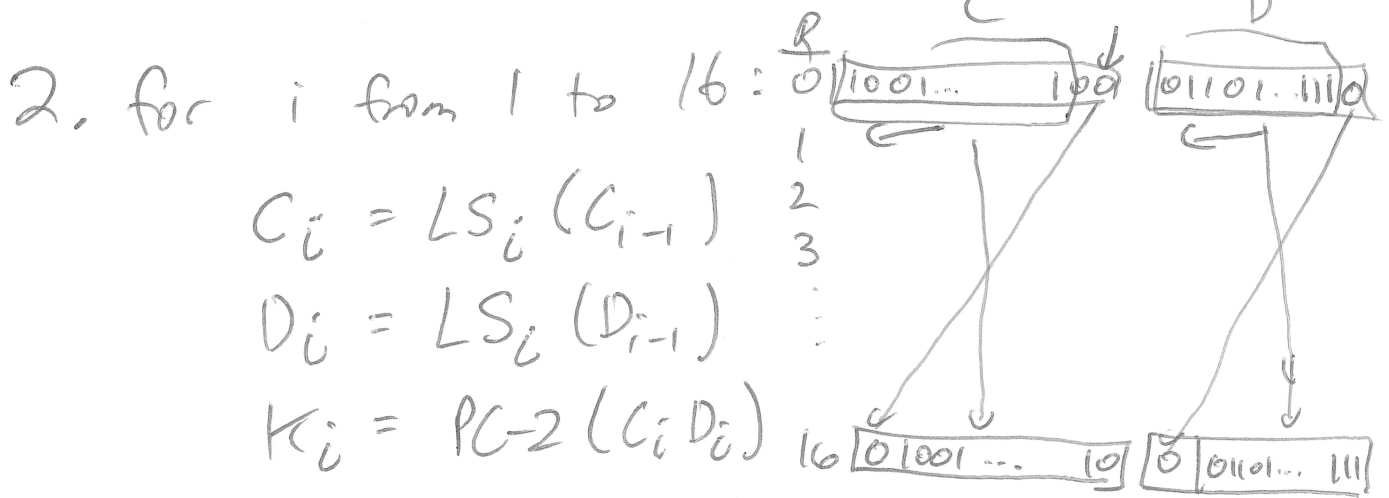
$$S_2(\underline{101100}) = 13$$

row = 2
col = 6
1101

Key Schedule

Given K (56 bits w/ 8 parity bits)
 calculate $K_1, K_2, K_3 \dots K_{16}$ the 16
 round keys each which use 48 bits.

1. $C_0 D_0 = PC-1(K)$
 ↓ 28 bits → 28 bits → 56 bits (won't "have" bits 8, 16, 24, ...)



PC-1:

	57	49	41	33	25	17	9	1
C_{ϕ}	58	50	42	34	26	18	10	2
	59	51	43	35	27	19	11	3
	60	52	44	36	57	55	47	39
	31	23	15	7	62	54	46	38
D_{ϕ}	30	22	14	6	61	53	45	37
	29	21	13	5	28	20	12	4
								63

PC-2 = 14, 17, 11, 24, 1, 5, 3, 28, 15, 6, 21, 10, 23, 19, 12, 4
 ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓
 10 51 34 60 49 17 33 57 2 9 19 42 3 35 26 25