

# CIS 3362 9/11/24

- (1) Finish Hill
- (2) Start ADFGVX

$$\begin{matrix} \rightarrow \\ M = \end{matrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix} \xrightarrow{\text{Inverse matrix (corresponding dec key)}} \underbrace{(ad-bc)^{-1} \pmod{26}} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$$

$$M \begin{pmatrix} p_1 \\ p_2 \end{pmatrix} = \begin{pmatrix} c_1 \\ c_2 \end{pmatrix}$$

$$M^{-1} \begin{pmatrix} c_1 \\ c_2 \end{pmatrix} = \begin{pmatrix} p_1 \\ p_2 \end{pmatrix}$$

$$\underline{\underline{M^{-1}M}} \begin{pmatrix} p_1 \\ p_2 \end{pmatrix} = \begin{pmatrix} p_1 \\ p_2 \end{pmatrix}$$

$$M^{-1} \times M = \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}}_{\text{Identity Matrix}} \begin{pmatrix} p_1 \\ p_2 \end{pmatrix} = \begin{pmatrix} 1 \times p_1 + 0 \times p_2 \\ 0 \times p_1 + 1 \times p_2 \end{pmatrix} = \begin{pmatrix} p_1 \\ p_2 \end{pmatrix}$$

$$\begin{array}{c}
 \begin{array}{cc}
 \cancel{a} & \cancel{b} \\
 \cancel{c} & \cancel{d}
 \end{array} \\
 \begin{array}{cc}
 \text{row} & \text{col} \\
 \begin{array}{cc}
 d & -b \\
 -c & a
 \end{array} & \begin{array}{cc}
 \text{col 1} & \text{col 2} \\
 \begin{array}{c} a \\ c \end{array} & \begin{array}{c} b \\ d \end{array}
 \end{array}
 \end{array}
 = \begin{pmatrix} ad-bc & bd-bd \\ -ac+ac & -bc+cd \end{pmatrix} \\
 \begin{array}{cc}
 2 \times 2 & 2 \times 2
 \end{array}
 = \begin{pmatrix} ad-bc & 0 \\ 0 & ad-bc \end{pmatrix}
 \end{array}$$

$$\underbrace{(ad-bc)^{-1}}_{\text{Determinant}} \pmod{26} \begin{pmatrix} ad-bc & 0 \\ 0 & ad-bc \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \checkmark$$

easily scales to larger "block sizes".

In modern ciphers, they usually now give directions to extend the key size if necessary.

AES  $\rightarrow$  128 bit  $\rightarrow$  192 bits  $\rightarrow$  256 bits

One other idea about Hill

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} 4 \\ 0 \end{pmatrix} = \begin{pmatrix} 4a+0b \\ 4c+0d \end{pmatrix} = \begin{pmatrix} 12 \\ 14 \end{pmatrix} \pmod{26}$$

cipher  plain

In theory 2 pairs of matching plain/cipher text will give you a total of 4 equations:

$$\left. \begin{array}{l} 4a \equiv 12 \pmod{26} \\ 4c \equiv 14 \pmod{26} \end{array} \right\} \begin{array}{l} 2 \text{ solutions} \\ \text{valid} \\ 2 \text{ solutions} \\ \text{valid} \end{array}$$

2 with a and b } If you're lucky these equations will have a unique sol.  
 2 with c and d }

$$4a \equiv 12 \pmod{26}$$

$$\Rightarrow 26 \mid (4a-12) \quad \text{def mod}$$

$$4a-12 = 26n, \quad n \in \mathbb{Z} \text{ (integer)}$$

$$2a-6 = 13n \Rightarrow 13 \mid (2a-6)$$

$$\Rightarrow 2a \equiv 6 \pmod{13}$$

$$a \equiv 3 \pmod{13}$$

$$a \equiv 3 \text{ or } 16 \pmod{26}$$

$$4c \equiv 14 \pmod{26}$$

$$26 \mid (4c-14)$$

$$4c-14 = 26n$$

$$2c-7 = 13n$$

$$2c \equiv 7 \pmod{13}$$

$$7(2c) \equiv 7(7) \pmod{13}$$

$$c \equiv 10 \text{ or } 23 \pmod{26}$$

$$\boxed{c \equiv 49 \equiv 10 \pmod{13}}$$

# ADFGVX

Cipher Germans used during WWI.

Initially it was ADFGX.

ADFGVX are the only letters in the ciphertext.

They were chosen because they are less likely to be confused for each other when transmitted as Morse code.

A = • —

G = — — •

F = • • — •

D = ~~•~~ — ••

X = — •• —

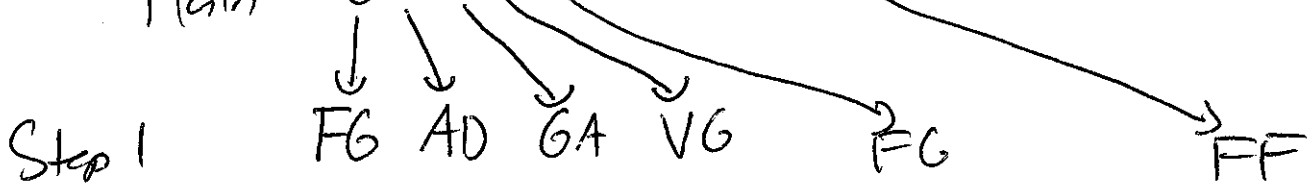
V = ••• —

RC	A	D	F	G	V	X
A	<del>0</del>	O	H	9	T	A
D	E	C	W	Z	V	S
F	R	8	4	G	2	1
G	K	J	U	X	P	Y
V	6	S	B	N	D	Q
X	F	3	M	I	7	L

→ KEY

map  
char → loc

Plan: GOKNIGHTS2024



Use grid, for each char, find its row label and col label and write that down in order.

Step 2: Second key (key word) used for transposition cipher.

↓	↓	↓	↓	↓	↓	↓
K	E	Y	W	O	R	D
(3)	(2)	(7)	(6)	(4)	(5)	(1)
F	G	A	D	G	A	V
G	X	G	F	G	A	F
A	V	V	D	F	V	
A	F	F				

A for repeated PAPER (3)(1)(4)(2)(5) letters

Ciphertext: VFAGXVFFGAA GGFAAVDFOAGVF  
 read columns down in order

- (1) SUBSTITUTE
- (2) MOVE AROUND

CLAUDE SHANNON  
 CONFUSION +  
 DIFFUSION