

10/25/23 - RSA

Diffe-Hellman - Key Exchange (Discrete Log Problem)

Ron Rivest [↑] - new faculty member MIT (1975)
read paper.

ONE WAY FUNCTION → easy forwards
hard invert

Convinced Adi Shamir, Leonard ~~Adleman~~^{Adleman}

GOAL: Public Key Cryptosystem (asymmetric)

Rivest had ideas

Shamir = ideas

Adleman shot them down.

More than a year!

Passover Party 1977 (Mar-Apr)

Overnight → sketched out RSA encryption

Patent → Company → Sold

RSA

Alice wants people to be able to send secure msg to her.

Alice generates 2 private keys: p, q large primes.

Calculate $n = pq$, n will public key
(security is based on factoring being hard)

Calculate $\phi(n) = (p-1)(q-1)$ // Private

Alice picks a random e , $\gcd(e, \phi(n)) = 1$
 e is another public key.

Alice computes $d = e^{-1} \bmod \phi(n)$ via EEA.

d is another private key.

Public: n, e

Private: $p, q, d, (\phi(n))$

$$\begin{aligned} ed &\equiv 1 \bmod \phi(n) \\ ed &= k\phi(n) + 1 \text{ for some int } k \\ &\rightarrow \text{rel prime } n. \end{aligned}$$

Send to Alice: Pick M $0 < M < n$

Bob computes $C = M^e \bmod n$
sends C to Alice.

Alice receives C . Computes $C^d \equiv \boxed{M} \pmod{n}$

$$C^d = (M^e)^d = M^{ed} = M^{k\phi(n)+1} = (M^{\phi(n)})^k \cdot M \equiv 1 \cdot M \pmod{n}$$

$$n = 11 \times 17 = 187$$

$$\phi(n) = 10 \times 16 = 160 = 2^4 \times 2 \times 5 = 2^5 \times 5$$

$$e = 21 \quad \text{What is } d?$$

$$d = 21^{-1} \pmod{160}$$

$$160 = 7 \times 21 + 13$$

$$21 = 13 + 8$$

$$13 = 8 + 5$$

$$8 = 5 + 3$$

$$5 = 3 + 2$$

$$3 = 2 + 1$$

$$d = 61$$

$$3 - 2 = 1$$

$$3 - (5 - 3) = 1$$

$$2 \times 3 - 1 \times 5 = 1$$

$$2(8 - 5) - 1 \times 5 = 1$$

$$2 \times 8 - 3 \times 5 = 1$$

$$2 \times 8 - 3(13 - 8) = 1$$

$$5 \times 8 - 3 \times 13 = 1$$

$$5(21 - 13) - 3 \times 13 = 1$$

$$5 \times 21 - 8 \times 13 = 1$$

$$5 \times 21 - 8(160 - 7 \times 21) = 1$$

$$5 \times 21 - 8 \times 160 + 56 \times 21 = 1$$

$$61 \times 21 - 8 \times 160 = 1$$

$$\underline{61} \times 21 \equiv 1 \pmod{160}$$

$$M = 97 \quad C = 97^{21} \equiv 20 \pmod{187} \quad (\text{Alice})$$

$$M = 20^{61} \equiv 97 \quad \text{Alice computes}$$

Opponent sees $n = 187$, $e = 21$ They have difficulty computing d .

if factor, then $\Rightarrow \phi(n) \Rightarrow d$.

(a) How knowing $\phi(n)$ can help you break it.

(b) Code how to convert text/bits $\Leftrightarrow M$

$$\begin{aligned}n &= 187 = pq \\ \phi(n) &= 160 = (p-1)(q-1) \\ &= pq - p - q + 1\end{aligned}$$

$$n - \phi(n) = 27 = p + q - 1$$

$$28 = p + q$$

$$q = (28 - p)$$

$$187 = p(28 - p)$$

$$187 = 28p - p^2$$

$$p^2 - 28p + 187 = 0$$

Quadratic Formulas will work w/ large ^{ints} ~~ints~~

RSA2. java

$$\begin{array}{c} \text{'C'} \text{'A'} \text{'T'} \\ 2 \ 0 \ 19 \end{array} = 2 \times 26^2 + 0 \times 26^1 + 19 \times 26^0$$

$$\text{HOME} = 7 \times 26^3 + 14 \times 26^2 + 12 \times 26^1 + 4 \times 26^0$$