

Number Theory

prime numbers $\Rightarrow$ Fermat's Thm $\rightarrow$ Euler's Thm $\rightarrow$ Miller Rabin
 mod, mod inverse (already seen) Primality Test
 $\rightarrow$ RSA requires the use of 2 large primes +
 reason it works is based on Euler's Thm.

* Discrete Log Problem (El Gamal Public Key System)

What's a regular log? $\log_2 64 = 6$
 because $2^6 = 64$

log is inverse of the exponent problem.

Input (Given): Ans, Base

Output : Exponent

New Problem

Input: Ans, Base, Mod

Output: exponent

What's smallest
 positive int x
 that makes this true?

$$5^x \equiv 6 \pmod{11}$$

Exp	0	1	2	3	4	5
Ans	1	5	3	4	9	1

$25 \equiv 3$
 $15 \equiv 4$
 $20 \equiv 9$

OOPS no ans!

$$3^x \equiv 6 \pmod{11}$$

Exp	0	1	2	3	4	5
Ans	1	3	9	5	4	1

$$2^x \equiv 6 \pmod{11} \quad \left. \vphantom{2^x} \right\} \textcircled{x=9}$$

Exp	0	1	2	3	4	5	6	7	8	9	10
Ans	1	2	4	8	5	10	9	7	3	6	1

answer

This is the discrete log problem

for prime numbers, p , we can calculate the cycle length for each possible base in modular exponentiation. We previously proved that the cycle length, c , for any base divides evenly into $p-1$.

$p=11$, $p-1=10$, all cycle lengths are either 1, 2, 5 or 10

Did program for $p=11$

* cycle len = 10	2, 6, 7, 8	$\phi(10)=4$
cycle len = 5	3, 4, 5, 9	$\phi(5)=4$
cycle len = 2	10	$\phi(2)=1$
cycle len = 1	1	$\phi(1)=1$

→ these values have property that they produce each possible non-zero mod value in a reasonably unpredictable order.

→ primitive roots OR generators

Once we know a primitive root exists, then we can prove there are exactly $\phi(p-1)$ of them.

$$2^{10} \equiv 1 \pmod{11}$$

and $2^2 \equiv 4 \pmod{11}$

$$4^5 \equiv (2^2)^5 \equiv 2^{10} \equiv 1 \pmod{11}$$

The Discrete Log Problem will be use
 with $\text{mod} = \text{prime}$, $\text{base} = \text{primitive root}$
 $\Rightarrow$ each possible ans exists.

Easiest / Slowest alg to solve

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    b, a, p mod
    res = 1
    for (int i = 0; i < p; i++) {
        if (res == a) return i;
        res = (res * b) % p;
    }
    Run time  $O(p * \text{oper})$ 
    
```

An alg that is a little bit faster

let $n = \lceil \sqrt{p} \rceil$ ^{ceiling (least int $\geq \sqrt{p}$)}

goal: find integers c and d such that

$$\underline{a}^{\underline{nc-d}} \equiv \underline{b} \pmod{\underline{p}}$$

Imagine $p=97$ $n=10$

loop through
 $c \begin{matrix} 1 \\ 0 \end{matrix} \begin{matrix} 10 \\ 10 \end{matrix}$

if $c=1$, set $d=0, 1, 2, 3, \dots, 10$

$\begin{matrix} 10-0 & 10-1 & 10-2 & \dots & 10-10 \\ a & a & a & \dots & a \end{matrix}$
 $c=2 \quad \begin{matrix} 20-0 & 20-1 & \dots & 20-10 \\ a & a & \dots & a \end{matrix}$

- ① Calculate a^{nc} for $c=0,1,2,\dots,n$
store answers in a chart
- ② Calculate a^d for $d=0,1,2,\dots,n$

$$a^d (a^{nc-d}) \equiv (b) a^d \pmod{p}$$

$$\Rightarrow a^{nc} \equiv b \cdot a^d \pmod{p}$$

for each d mult
 $b \cdot a^d$ and see

if a^{nc} is in the lookup
chart

TABLE

$a^{nc} \rightarrow c$

Step 1 takes $O(\sqrt{p})$ time
Step 2 takes $+ O(\sqrt{p})$ time

$O(\sqrt{p})$