

DES

- talk about the function $f(A, J)$ that's in each round (S-box design)

- key schedule

- look at my code

$f(A = R_{i-1} \text{ 32 bits}, J = K_i \text{ 48 bits})$

1. Calculate $E(A)$ - expand A to 48 bits using table E (like a perm matrix)

A is ~~11E~~ was 7CF63D28

0111 1100 1111 0110 0011 1101 0011 1000

$E(A) = 32, 1, 2, 3, 4, 5, 4, 5$

00111111 (3F is 1st byte output)

2. Calc $E(A) \oplus J$ (round key xor w/round key)

hex $E(A) = 3F \dots$

$\oplus J = 8C$

B3

3. let step 2 output be $B = \underbrace{B_1 B_2 \dots B_8}_{6 \text{ bits}}$

$$B_{\text{hex}} = \underline{B} \underline{3} \underline{C} \underline{C} \underline{0} \underline{7} \underline{4} \underline{9} \underline{C} \underline{F} \underline{D} \underline{2}$$

$$B_1 = \underline{101100}$$

$$B_5 = \underline{010010}$$

$$B_2 = \underline{111100}$$

$$B_6 = \underline{011100}$$

$$B_3 = \underline{110000}$$

$$B_7 = \underline{111111}$$

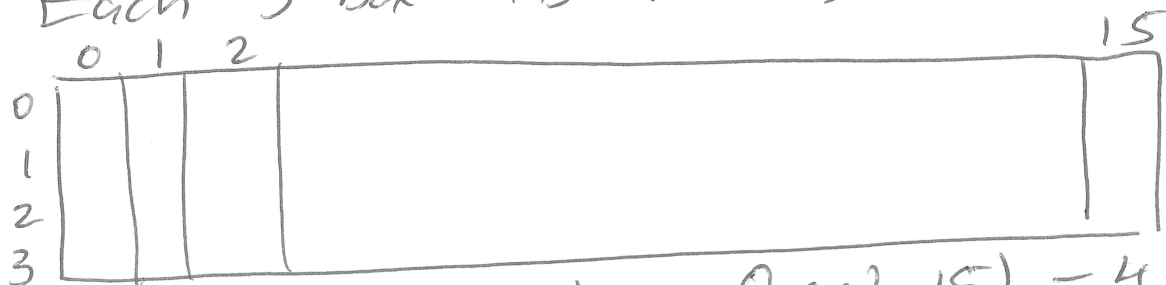
$$B_4 = \underline{000111}$$

$$B_8 = \underline{010010}$$

$$C_i = \underline{S_i(B_i)}$$

$$S_1(B_1), S_2(B_2), \dots, S_8(B_8)$$

Each S-box has 4 rows, 16 cols



outputs (values in box 0 and 15) - 4 bits
input = 6 bits output = 4 bits (only non-linear component)

$$B_i = b_1 b_2 b_3 b_4 b_5 b_6$$

$$\text{row} = b_1 b_6$$

$$\text{col} = b_2 b_3 b_4 b_5$$

$$S_1(\underline{101100})$$

$$\text{row} = 10 \quad (2)$$

$$\text{col} = 0110 \quad (6)$$

$$S_1[2][6] = 2 [0010]$$

$$S_2(\underline{111100}) = S_2[2][14] = 2 [0010] \quad \text{Output}$$

$$S_3(\underline{110000}) = S_3[2][8] = 11 [1011] \quad \text{15}$$

$$S_4(\underline{000111}) = S_4[1][3] = 5 [0101]$$

4. Return $P(C)$ $\{P$ is another perm matrix $\}$

S-box Design

P0: each row is perm of $0, 1, 2, \dots, 15$

P1: no s-box is a linear or affine function of its inputs.

P2: Changing one input bit to an S-box causes at least 2 output bits to change.

P3: For any Sbox and any input x , $S(x)$ and $S(x \oplus 001100)$ differ in at least 2 bits.

P4: for any s-box, any input x , and for $e, f \in \{0, 1\}$
 $S(x) \neq S(x \oplus 11ef00)$

P5: If we fix one input bit, and look at one output bit, there are 32 possibilities. The #s of 0s and 1s in that output must be in between 13 and 19.

Alg

1. $IP(X) = L_0 R_0$

2. Run 16 rounds:

$$L_i = R_{i-1}$$

$$R_i = L_{i-1} \oplus F(R_{i-1}, K_i)$$

3. Output $IP^{-1}(R_{16}L_{16})$

Will have key schedule for Monday.