

# DES Day 2

- S-box review
- review function
- Key generation
- Look @ Code

A single S-box input = 6 bits  $\Rightarrow$  output = 4 bits

8 boxes labels  $S_1$  through  $S_8$

$S_1$  (bits 1-6)

$S_2$  (bits 7-12)

$\vdots$

$S_8$  (bits 43-48)

input =  $b_1 b_2 b_3 b_4 b_5 b_6$

row =  $b_1 b_6$  (0-3)

col =  $b_2 b_3 b_4 b_5$  (0-15)

$S_1$ (101110)	$\rightarrow$	row = 10 (2)	$\rightarrow$	<span style="border: 1px solid black; padding: 2px;">11</span>
$S_2$ (010111)	$\rightarrow$	row = 01 (1)	$\rightarrow$	<span style="border: 1px solid black; padding: 2px;">10</span>
$S_3$ (111101)	$\rightarrow$	row = 11 (3)	$\rightarrow$	<span style="border: 1px solid black; padding: 2px;">2</span>
$S_4$ (000110)	$\rightarrow$	row = 00 (0)	$\rightarrow$	<span style="border: 1px solid black; padding: 2px;">3</span>

## S-box Criteria (Cryptography Theory + Practice Stinson)

- P0) each row permutation 0 to 15.
- P1) No S-box linear or affine function of its input.
- P2) Change 1 bit of input at least 2 bits of output change.
- P3) For any S-box any input  $x$ ,  $S(x)$  and  $S(x \oplus 001100)$  differ in at least 2 bits.

P4) For any S-box  $S$ , any input  $x$  and bits  $e, f$ ,  $S(x) \neq S(x \oplus 11ef00)$

P5) In any S-box, if we fix one input bit and look @ the distribution of a fixed output bit, the # of 0s + 1s will each be in between 13 and 19, inclusive

64 possible input

bit 4 = 1, input for  $x$  abcdef.

Look @ output bit 2, look @ bit 2

$$S(abcdef) = \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \end{array} \begin{array}{l} \geq 13 \text{ zeros} \\ \leq 19 \text{ zeros} \end{array}$$

all 32 inputs

## Round Review

$$L_i = R_{i-1}$$

$$R_i = L_{i-1} \oplus f(R_{i-1}, K_i)$$

$$f(A, J)$$

1. Calculate  $E(A)$  32 bits  $\rightarrow$  48 bits

2.  $E(A) \oplus J$  Xor w/ ~~bit~~ key

3. Split in 8 groups of 6 bits:  $B_1, \dots, B_8$

4.  $S_i(B_i) \rightarrow C_i$  48 bits  $\rightarrow$  32 bits

5. Return  $P(C)$

# DES Key Generation

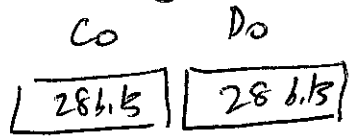
1.  $C_0 D_0 = PC^{-1}(K)$  (permute key)

2. for  $i = 1$  to  $16$ :

$$C_i = LS_i(C_{i-1})$$

$$D_i = LS_i(D_{i-1})$$

$$K_i = PC^{-2}(C_i D_i)$$



LS = cyclic left shift  
(1 bit or 2 bits)

$LS_i = 1$  if

$i = 1, 2, 9$  or  $16$

else  $LS_i = 2$

$$LS_{1bit}(1011010) = 0110101$$

$$LS_{2bits}(1011010) = 1101010$$

Will show you How to get direct mappings for each round key.

PC-1 left  $\frac{1}{2}$

<del>57</del>	(49)	41	(33)	25	(17)	9
1	58	50	42	(34)	26	18
(10)	2	59	(51)	43	35	27
19	11	3	(60)	52	44	36
(57)						

PC-2 strts 14, 17, 11, 24, 1, 5, 3, 28  
 $\downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow$   
 10 51 34 60 49 17 33 57