

19/3/19 (1)

Fermat's Formula $p \in \text{Prime}$ $\gcd(a, p) = 1$ $a^{p-1} \equiv 1 \pmod p$

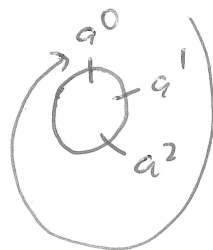
Euler ϕ func: $\phi(n) = \#$ values $\{1, 2, \dots, n\}$
rel prime w/n.

Euler's formula $\gcd(a, n) = 1$, $a^{\phi(n)} \equiv 1 \pmod n$

$\text{ord}_{a,n} = \min\{k\} \mid a^k \equiv 1 \pmod n.$
 $k > 0$

for all a, n $\gcd(a, n) = 1$, $\text{ord}_{a,n} \mid \phi(n)$.

In modular exponentiation,
each base has a cycle.



For primes^p, the longest possible
cycle is length $p-1$.

For all primes^p, there exist values a ,
such that $\text{ord}_{a,p} = p-1$.

But for compositesⁿ, there may or may not
be an a such that $\text{ord}_{a,n} = \phi(n)$.

last time we showed for all a , $n=15$
 $\text{ord}_{a,n} = 1, 2$ or 4 none $= 8$.

Primality Testing (we need large
primes to run RSA + other public
key schemes)

— Trial Division is TOO slow!

Need to Differentiate btw primes, composites.

for all primes $a^{p-1} \equiv 1 \pmod p$.

Idea: Given n , Pick random a .

10/23/19 @

Calculate $a^{n-1} \bmod n$. If answer $\neq 1$,
proof n is composite.

Problem: for some composites and some choices
of a , perhaps $a^{n-1} \equiv 1 \bmod n$.

→ Call these pseudoprimes!

In general for random choice of n, a
where n is composite, this is true
less than 50% of time.

Unfortunately, for a very few values of
 n , it turns out that for ALL a 's $\gcd(a, n) = 1$
 $a^{n-1} \equiv 1 \bmod n$. (Carmichael Numbers)
561, 1729, ...?

A Pretty Good Primality Test

Repeat 100 (or) times:

(1) Pick random a $1 < a < n-1$

(2) Calculate $a^{n-1} \bmod n$, Calculate $\gcd(a, n)$.

(3) If value $\neq 1$, return composite

* Return "IS PROBABLY PRIME"

(2b) If $\gcd \neq 1$
return false

As long as $n \neq$ Carmichael, this generally
works!

Main Primality Test Used: Miller - Rabin ^{10/23/19③}

Improves on "pretty good test" by trying to catch Carmichael #s.

Miller Rabin (int n)

- $n = 101$
 $n-1 = 100$
 $= 2^2 \cdot 25$
 $(a^{25})^2 \stackrel{50}{=} a^{100}$
1. $n-1 = 2^k n'$, $n' \in \text{odd}$, $k \geq 1$
2. ~~Calculate~~ Pick random a
3. if $\gcd(a, n) \neq 1$, return false (composite)
4. Calculate $X = a^{n'} \bmod n$.
- $n = 97$
 $n-1 = 96$
 $= 2^5 \cdot 3$
 $a^3, a^6, a^{12}, a^{24}, a^{48}, a^{96} = 1$
5. If value in step 4 $\neq 1$, return is prob prime.
6. for ($i = 0$; $i < k-1$; $i++$) {
 $X = X^2 \bmod n$
 if ($X \neq -1$) is prob prime
 if ($X \neq 1$) return composite

for composite chance fails 1 trial is $\leq \frac{1}{4}$.

Discrete Log Problem

Given a base a , and a result X , determine the exponent e s.t.

$$a^e \equiv X \bmod p.$$

$p \in \text{prime}$

$$\log_2 64 = 6$$

$$2^6 \equiv 64 \bmod 97.$$

This is a HARD problem. No one yet knows how to do this efficiently (poly time). 10/23/1994

for each prime p , $\exists a \in \{1, 2, \dots, p-1\}$ s.t

$\text{ord}_{a,p} = p-1$. These values of a are called generators, primitive roots.

$p=11$

base \ exp	1	2	3	4	5	6	7	8	9	10
2	2	4	8	5	10	9	7	3	6	1
8	8	9	6	4	10	3	2	5	7	1
7	7	5	2	3	10	4	6	9	8	1
6	6	3	7	9	10	5	8	4	2	1
4	4	16	64	256						
	4	5	9	3	1					

$$a^{p-1} \equiv 1$$

mod but not small

$$(a^{(k) \cdot p-1}) \equiv 1$$