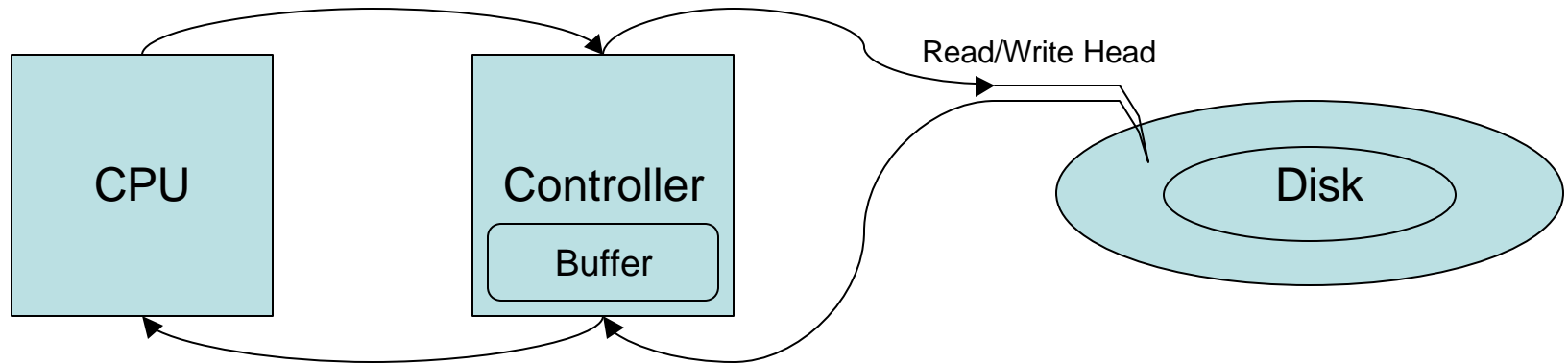
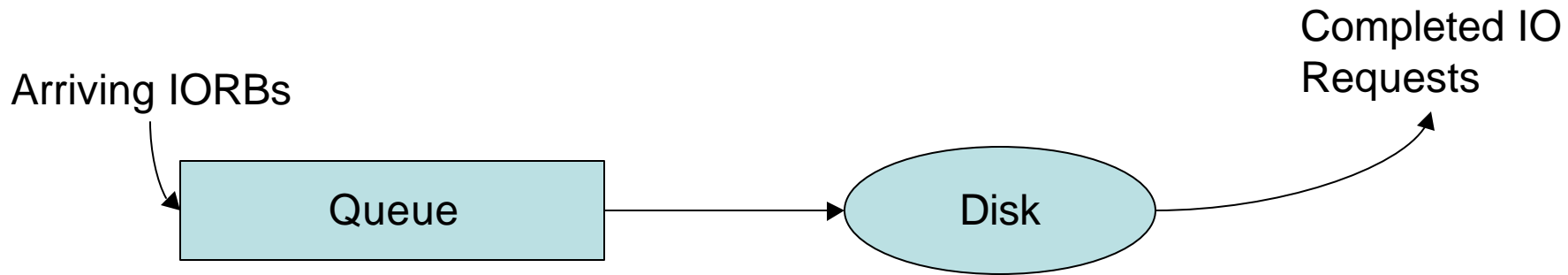
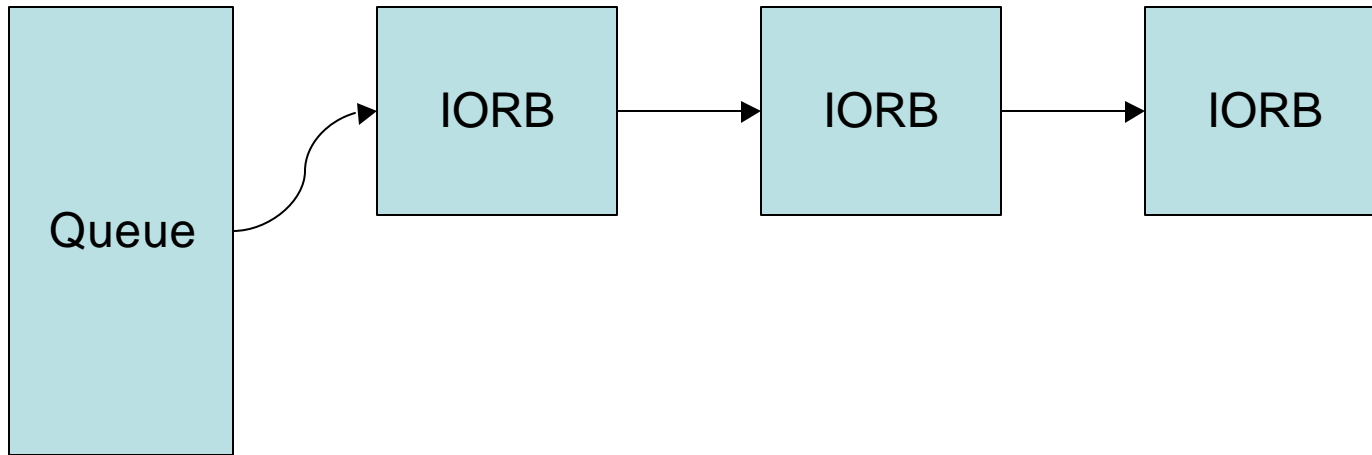


I/O Hardware



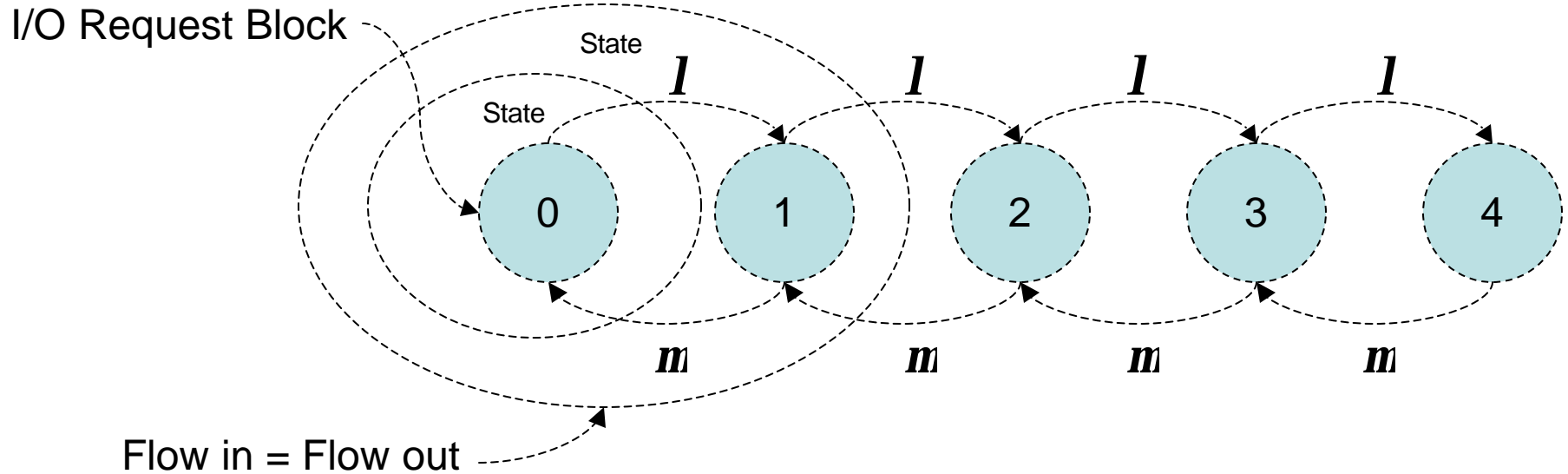
- CPU sends I/O request Blocks to controller, which in turn sends signals to the Read/Write Head to retrieve information from the disk
 - Disks are divided into tracks, (concentric disks on the surface of the media)
 - It is possible to arrange data in such a manner as to optimize the movement of the R/W Head across the tracks (i.e., rather than visiting tracks in the order 3-7-2, visit in the order 2-3-7 to allow one smooth motion)
- When the information has been retrieved, the controller sends an I/O Exception to the CPU

I/O Queue

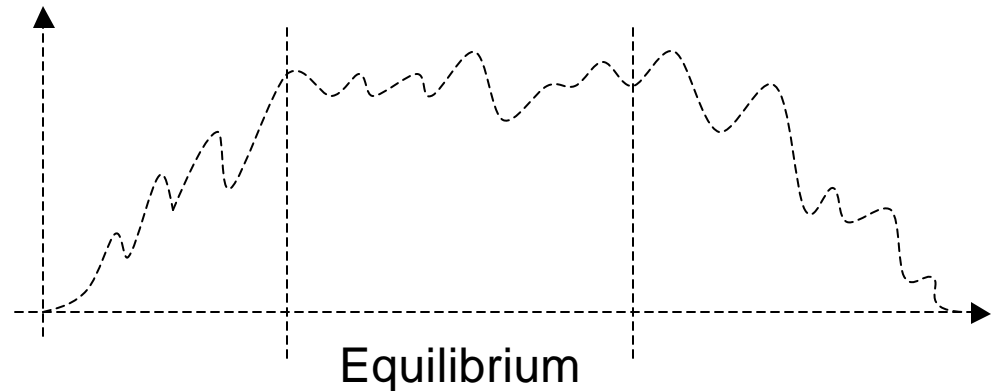


The queue is the final order of I/O's that is to occur. It may be reordered when new requests come in to make best use of the R/W Head

I/O Equilibrium



After a certain period of build-up the system works into an equilibrium, shown by the IORB x Time graph to the right, and the diagram above. In each state, the Flow of Incoming IORB equals the Flow of completed IORB's.



Some Definitions

A = Number of IO Request arrivals

C = Number of completed IO Requests

T = Length of time system is observed

B = Length of time the device is busy

S = Service time (per request)

Arrival Rate: $I = \frac{A}{T}$

Throughput: $x = \frac{C}{T}$

Utilization: $U = \frac{B}{T} = \frac{I}{m}$

Service Time: $s = \frac{1}{m}$

Probabilities of IO System

P_0 Probability of having 0 I/O requests in the system

P_1 Probability of having 1 I/O requests in the system

P_2 Probability of having 2 I/O requests in the system

P_i Probability of having i I/O requests in the system

Assuming an equal flow-in rate and flow-out rate we get the following...

$$mP_1 = lP_0 \Rightarrow P_1 = \frac{l}{m}P_0$$

$$mP_2 = lP_1 \Rightarrow P_2 = \frac{l}{m}P_1$$

\vdots

$$mP_i = lP_{i-1} \Rightarrow P_i = \frac{l}{m}P_{i-1}$$

Note: $\frac{l}{m} = U$

where U is the utilization of IO system

Finding P_i

Given that:

$$P_2 = \frac{1}{m} P_1$$

substituting the following $P_1 = \frac{1}{m} P_0$ and we get

$$P_2 = \left(\frac{1}{m}\right)\left(\frac{1}{m}\right)P_0 = \left(\frac{1}{m}\right)^2 P_0$$

Continue the pattern for P_i we get:

$$P_i = \left(\frac{1}{m}\right)^i P_0 = U^i P_0 \quad \text{where } i \text{ is any given state}$$

Sum of Probabilities

$$P_0 + P_1 + P_2 + \dots + P_k = \sum_{i=0}^k P_i = 1$$

substituting $P_i = U^i P_0$ we get

$$\sum_{i=0}^k U^i P_0 = 1 \Rightarrow P_0 = \frac{1}{\sum_{i=0}^k U^i}$$

Geometric Series

Please recall the following about a geometric series

$$\sum_{i=0}^{\infty} r^i = \frac{1}{1-r} \quad \text{where } |r| < 1$$

Assume there exists an infinite number of states

$$P_o = \frac{1}{\sum_{i=0}^k U^i} = \frac{1}{1-U} = 1-U$$

where $|U| < 1$

In other words, the utilization, U , must always be less than 100% or else the system will not function

Average System State

\bar{N} is the average system state at a finite small amount of time

N = The sum of the states times the probability of the state occurring, or...

$$\bar{N} = \lim_{t \rightarrow 0} \sum_{i=0}^k iP_i = \lim_{t \rightarrow 0} \sum_{i=0}^k i(1-U)U^i$$

$$\bar{N} = \lim_{t \rightarrow 0} U(1-U) \sum_{i=0}^k iU^{i-1}$$

$$iU^{i-1} = \frac{\partial(U^i)}{\partial U}$$

Noting that the first derivative of $U^i = iU^{i-1}$ we can substitute this into the equation

$$\bar{N} = \lim_{t \rightarrow 0} U(1-U) \sum_{i=0}^k \left(\frac{\partial U^i}{\partial U} \right)$$

Average System State

Solving the for the geometric sum and the derivative of the result we get the following

$$\bar{N} = \lim_{t \rightarrow 0} U(1-U) \left(\frac{\partial \left(\frac{1}{1-U} \right)}{\partial U} \right) = \lim_{t \rightarrow 0} \frac{U(1-U)}{(1-U)^2}$$

Which, in turn, gives us the final equation:

$$\bar{N} = \lim_{t \rightarrow 0} \frac{U}{(1-U)}$$