



# An Introduction to Graph-Cut

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## An Introduction to Graph-Cut

- Overview
  - Min-cut
  - Normalized-Cut
- Applications

## An Introduction to Graph-Cut

- Graph-cut is an algorithm that finds a globally optimal segmentation solution.
- Also known as Min-cut.
- Equivalent to Max-flow. [1]

[1] Wu and Leahy: An Optimal Graph Theoretic Approach to Data Clustering:...

## What is a “cut”?

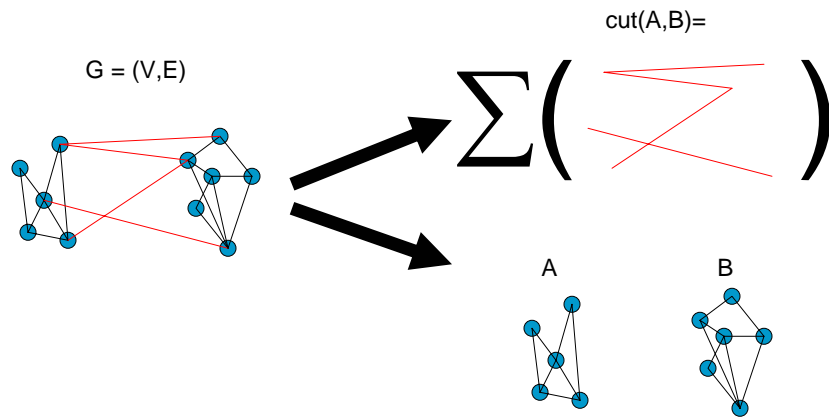
A graph  $G = (V, E)$  can be partitioned into two disjoint sets,  $A, B$ ,  $A \cup B = V$ ,  $A \cap B = \emptyset$  by simply removing edges connecting the two parts.

The degree of dissimilarity between these two pieces can be computed as total weight of the edges that have been removed. In graph theoretic language it is called the *cut*.

$$cut(A, B) = \sum_{u \in A, v \in B} w(u, v) \quad [2].$$

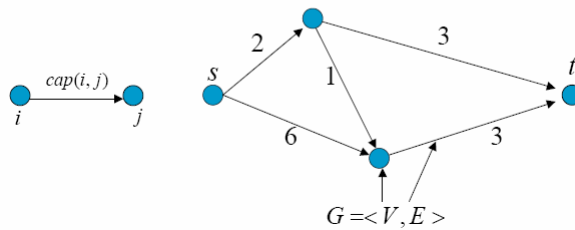
[2] Shi and Malik: Normalized cuts and image segmentation.

## Example cut



## Finding the Minimum-cut

- A source node and a sink node
- Directed edge (Arc)  $\langle i, j \rangle$  from node  $i$  to node  $j$
- Each arc  $\langle i, j \rangle$  has a nonnegative capacity  $\text{cap}(i, j)$
- $\text{cap}(i, j) = 0$  for non-existent arcs



## Finding the Minimum-cut

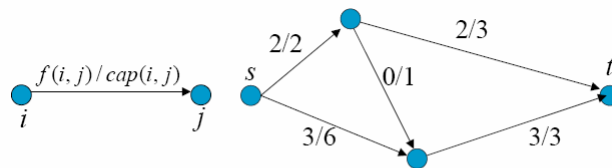
- Flow is a real value function  $f$  that assign a real value  $f(i, j)$  to each arc  $\langle i, j \rangle$  under :

- Capacity constraint :  $f(i, j) \leq \text{cap}(i, j)$
- Mass balance constraint:

$$\sum_{\langle i, j \rangle \in E} f(i, j) - \sum_{\langle k, i \rangle \in E} f(k, i) = \begin{cases} 0 & i \in V - \{s, t\} \\ |f| & i = s \\ -|f| & i = t \end{cases}$$

$|f|$  is the value of flow  $f$

## Finding the Minimum-cut

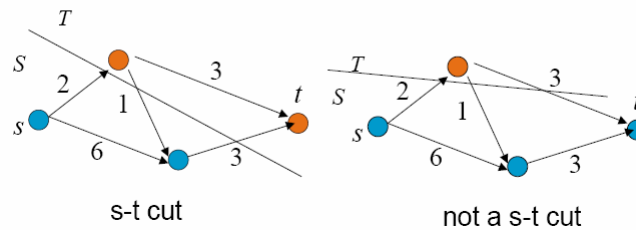


An example of flow

- maximum flow is the flow has maximum value among all possible flow functions

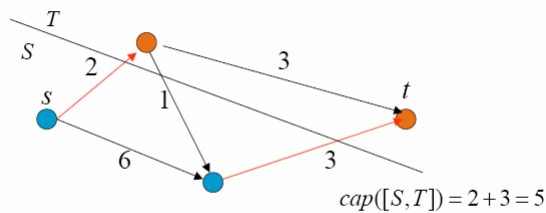
## Finding the Minimum-cut

- A cut is a partition of node set  $V$  which has two subsets  $S$  and  $T$
- A cut is a s-t cut iff  $s \in S, t \in T$



## Finding the Minimum-cut

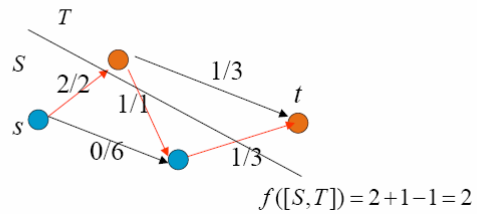
- $cap([S, T]) = \sum_{\langle i, j \rangle \in E, i \in S, j \in T} cap(i, j)$



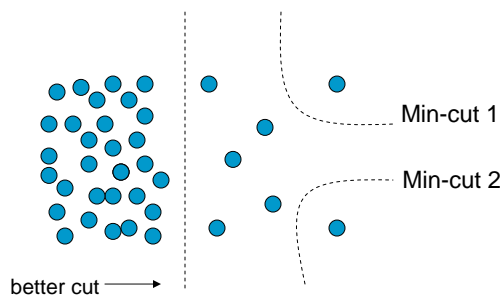
- Minimum cut is the s-t cut whose capacity is minimum among all possible s-t cuts

## Finding the Minimum-cut

$$f([S, T]) = \sum_{\langle i, j \rangle \in E, v_i \in S, v_j \in T} f(i, j) - \sum_{\langle j, i \rangle \in E, v_i \in S, v_j \in T} f(j, i)$$



## The problem with Min-cut



Need to account for cluster similarity

## Normalized-cut

- Instead use *normalized cut* (*Ncut*).

$$Ncut(A, B) = \frac{cut(A, B)}{assoc(A, V)} + \frac{cut(A, B)}{assoc(B, V)}$$

$$assoc(A, V) = \sum_{u \in A, t \in V} w(u, t)$$

$$Nassoc(A, B) = \left( \frac{assoc(A, A)}{assoc(A, V)} + \frac{assoc(B, B)}{assoc(B, V)} \right)$$

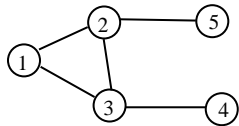
## Normalized-cut

$$\begin{aligned} Ncut(A, B) &= \frac{cut(A, B)}{assoc(A, V)} + \frac{cut(A, B)}{assoc(B, V)} \\ &= \frac{assoc(A, V) - assoc(A, A)}{assoc(A, V)} + \frac{assoc(B, V) - assoc(B, B)}{assoc(B, V)} \\ &= 2 - \left( \frac{assoc(A, A)}{assoc(A, V)} + \frac{assoc(B, B)}{assoc(B, V)} \right) \\ &= 2 - Nassoc(A, B) \end{aligned}$$

## Pixel labeling problem

Given

$$\mathcal{S} = \{1, \dots, n\} \quad \mathcal{N} \subseteq \mathcal{S} \times \mathcal{S}$$



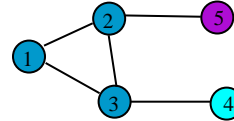
$$\mathcal{L} = \{l_1, \dots, l_m\}$$

*Assignment cost* for giving a particular label to a particular node. Written as  $D$ .

*Separation cost* for assigning a particular pair of labels to neighboring nodes. Written as  $V$ .

Find

$$\text{Labeling } f = (f_1, \dots, f_n)$$



Such that the sum of the assignment costs and separation costs (the energy  $E$ ) is small

## Energy Minimization

- Optimizing the labeling problem can be thought of as minimizing some energy function.

$$E(f) = \sum_{p \in P} D_p(f_p) + \sum_{p, q \in N} V_{p, q}(f_p, f_q)$$

$$D_p(f_p)$$

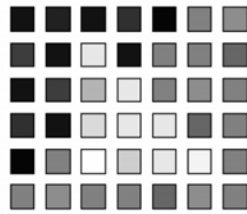
measure of image discrepancy

$$\sum_{p, q \in N} V_{p, q}(f_p, f_q)$$

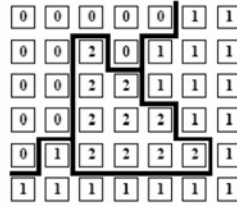
measure of smoothness or other visual constraints



# The Labeling Problem



(a) An image



(b) A labeling

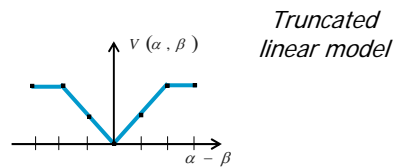
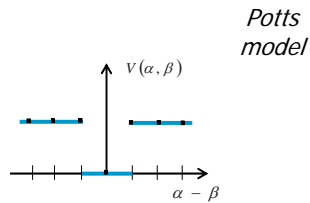
Common idea behind many Computer Vision problems

Assign labels to pixels based on noisy measurements (input images)

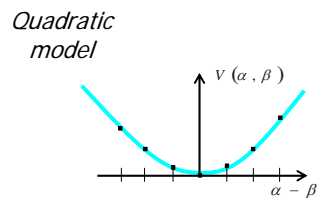
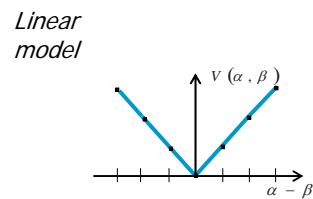
In the presence of uncertainties, find the best Labeling !

(Stereo, 3D Reconstruction, Segmentation, Image Restoration)

## Choices of $V$ Robust



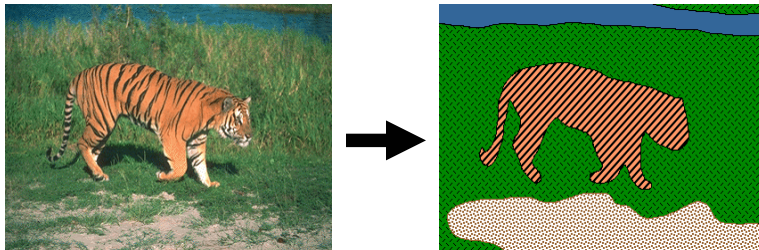
## Not robust



## What do graph cuts provide?

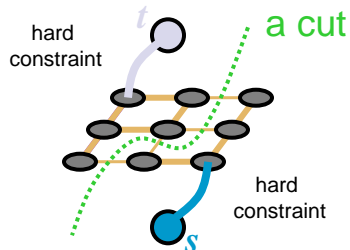
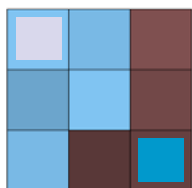
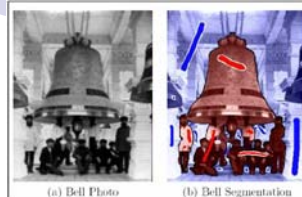
- For less interesting  $V$ , polynomial algorithm for global minimum!
- For a particularly interesting  $V$ , approximation algorithm
  - Proof of NP hardness
- For many choices of  $V$ , algorithms that find a “strong” local minimum
- Very strong experimental results

## Graph Cut based Segmentation



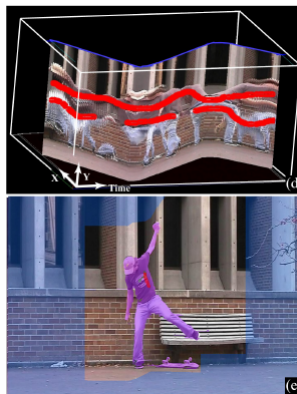
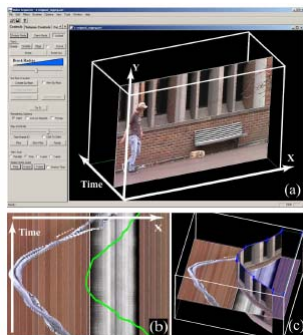
# Graph Cut based Segmentation

User Guided Segmentation;  
Specifies hard constraints.



# 3D Applications

2D and Time



[3] Wang et al.: Interactive Video Cutout

# Some Results



Input Image



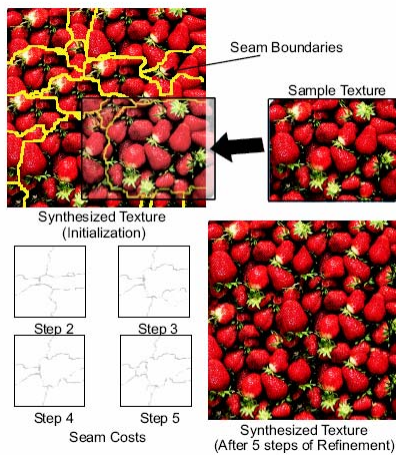
Efros and Freeman '01  
(Image Quilting)



Graph-cut

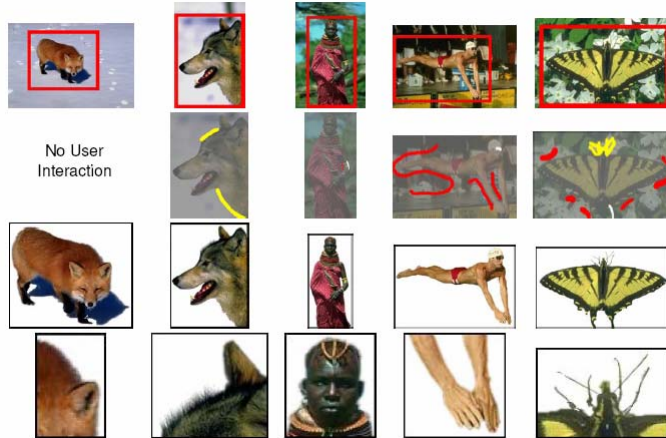
[3] Kwatra et al.: Graphcut Textures: Image and Video Synthesis Using Graph Cuts

# Some Results



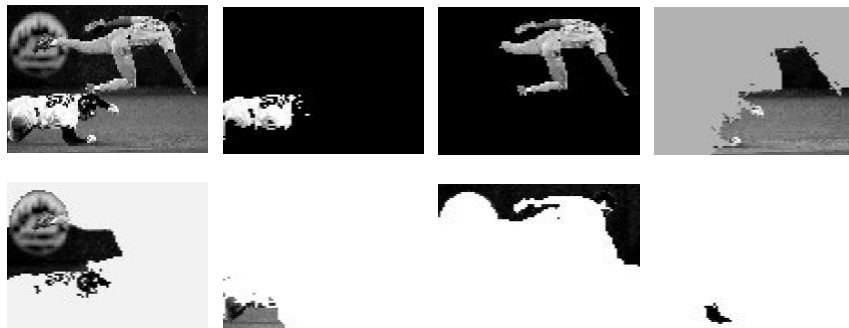
[4] Kwatra et al.: Graphcut Textures: Image and Video Synthesis Using Graph Cuts

## Some Results

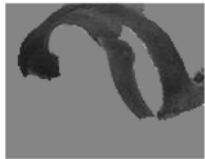
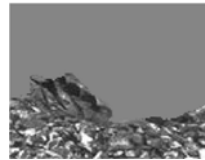


[5] Rother, Kolmogorov and Blake:  
"GrabCut" - Interactive Foreground Extraction using Iterated Graph Cuts

## Some Results


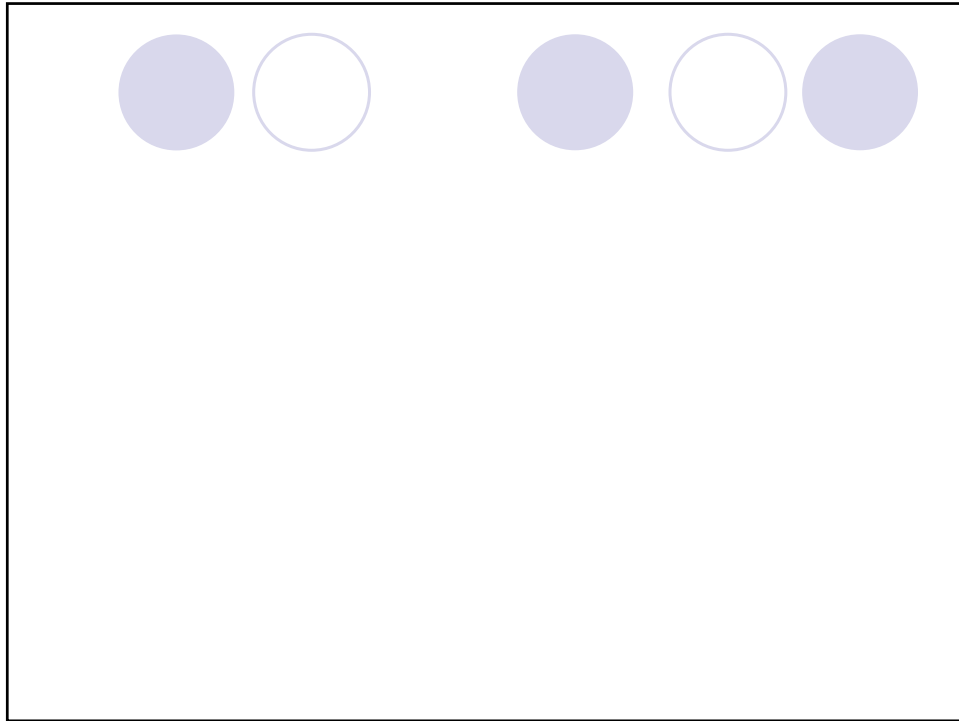


Some Results



Some Results





- So far weight measurement has been distance, but could also use appearance, texture, or other information to calculate similarity measure.