

Presentation Overview

- Introduction to the problem
- 3D rendering techniques using sketch based interfaces
- SMARTPAPER's features
- Processing Pipeline
 - 2D processing
 - 3D Geometry Reconstruction
- Feedback/Cutting/Joining/Rendering
- Conclusion

Introduction

- It is natural to sketch out a 3D object conceptually
- Current methods involve taking a conceptual sketch and using a complex system (such as CAD) to create it
- A better design would be to allow the sketched interface to immediately translate to a 3D object





- Uses several gestures
- Simple to create primitives
- May be unintuitive (depends on gesture set)





Introducing SMARTPAPER

- A combination of the other techniques
- Allow free form drawing
- Allow gesture manipulation
- Render from 2D graphs



















Pre-Processing: Over Tracing

- Remove Over Tracing
 - A pair of strokes A,B are over tracing if:
 - They have nearly equal slopes
 - One End Point of A lies in the X and Y ranges of the endpoints of B
- Once found, a connection is made in two passes, A's starting point to B's ending point



Over tracing Example

Pre-Processing: Gesture Recognition

- Arrow was recognized as proof of concept
 - Has to be drawn in two strokes
 - Both recognizer and cutting modules query if this gesture was drawn
 - Either recognizer or cutting module is called depending on the operation

2D Graph Generation

- Graph is generated with a connectivity matrix, each vertex contains (x,y) coordinates
- Imagine taking the endpoints from each line to create vertices
- If performing extrusion, a copy of one graph is made along the direction of the extrusion arrow and edges are connected
- Then use "clustering" to determine proximity and combine graphs
- We assume all sketches are closed objects
 - All vertices must have at least a degree of 3
 - Remove any vertex that does not meet this rule



3D Reconstruction: Face Determination

- All Faces are Cycles
- Not all cycles are faces though
- Using the closed object assumption
 - All edges of graph G are part of exactly two faces
 - The shortest path of any two vertices V1, V2 is the same length as the path in the face F
 - Proof by contradiction (Available in backup slides)
- Two algorithms are used to determine faces
 - Edge Coherence Algorithm
 - Modified Dijkstra Algorithm







Dijkstra Pseudo Code

S: Set of edges that are not part of at least 2 faces R: Set of edges that are part of at least 2 faces M: Matrix having faces as rows and edges as columns.

//Initialize S to all edges that are not part of at least 2 faces. If a current object is being augmented, this set is a proper subset of the edge set of the graph.

while S is not empty do

 $e \in S, e = (v_1, v_2)$

//if $e' \in R$ shares an end vertex with e, then mark end vertices of e' as traversed

//Find all edge-disjoint v_1, v_2 -paths, mark M accordingly.

//Traverse M column-wise and mark all edges are that are part of more than 2 faces. Then traverse M row-wise to delete all faces consisting of only marked edges.

G = G - e end while



More on 3D Reconstruction

- Using hints, create three Z layers
 - One layer consists of vertices to which only hidden edges are incident
 - Second layer is of vertices to which only visible edges are incident
 - Third layer is all other vertices
- These layers give each vertex an initial Z value
- The 3D object is finally represented as a new graph, similar to boundary representation (B-Rep)
 - Boundary Representation is a representation of faces, edges, and vertices. It is currently used in CAD systems.
- Use this graph to reproject when using different viewpoints















Conclusions

- Seems useful, though not compared with other sketching systems only CAD
- Cannot evaluate curved 3D objects
- Could make use of more gestures
- Thresholds and constraints often used
- Focuses on architectural design techniques



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Proof of Definition A

Definition A: All edges of graph G are part of exactly two faces. Every valid face F of G is such that for all pairs $v_1, v_2 \in V(G)$ that are in F, the shortest v_1, v_2 -path in G is of the same length as the v_1, v_2 -path in F.

Justification: The first statement is a property of closed, non-laminar, rigid objects. If the second statement is not true, let P be the shortest v_1, v_2 -path in G and let P' be the shorter v_1, v_2 -path in F. There is at least one edge in P not in P'. PP' thus creates a smaller closed walk and hence a smaller cycle C containing v_1 and v_2 . The edge set $E(C) - (E(C) \cap E(F))$ divides face F into two or more different planes, which is a contradiction as F is a valid face and hence is planar.

