

# Gene expression & Clustering

# Determining gene function

- Sequence comparison tells us if a gene is similar to another gene, e.g., in a new species
  - Dynamic programming
  - Approximate pattern matching
- Genes with similar sequence likely to have similar function
- Doesn't always work.
  - “Homologous” genes may not be similar enough at the sequence level, to be detected this way
- New method to determine gene function: directly measure gene activity (DNA arrays)

# DNA Arrays--Technical Foundations

- An array works by exploiting the ability of a given mRNA molecule to hybridize to the DNA template.
- Using an array containing many DNA samples in an experiment, the expression levels of hundreds or thousands genes within a cell by measuring the amount of mRNA bound to each site on the array.
- With the aid of a computer, the amount of mRNA bound to the spots on the microarray is precisely measured, generating a profile of gene expression in the cell.

# Gene expression

- Microarray gives us an  $n \times m$  expression matrix  $I$ 
  - Each of  $n$  rows corresponds to a gene
  - Each of  $m$  columns corresponds to a condition or time point
  - Each column comes from one microarray
- $I(j,k)$  is the expression level of gene  $j$  in condition/experiment  $k$
- If two genes (rows) have similar “expression profiles”, then
  - they may be related in function
  - they may be “co-regulated”

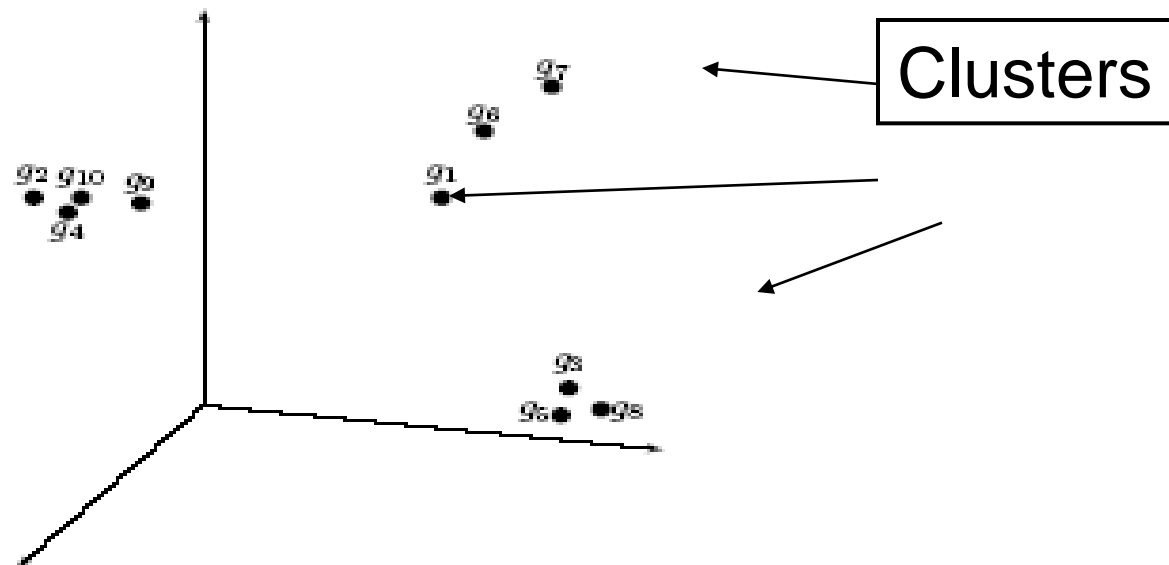
# Clustering of Microarray Data

Time	1 hr	2 hr	3 hr
$g_1$	10.0	8.0	10.0
$g_2$	10.0	0.0	9.0
$g_3$	4.0	8.5	3.0
$g_4$	9.5	0.5	8.5
$g_5$	4.5	8.5	2.5
$g_6$	10.5	9.0	12.0
$g_7$	5.0	8.5	11.0
$g_8$	2.7	8.7	2.0
$g_9$	9.7	2.0	9.0
$g_{10}$	10.2	1.0	9.2

(a) Intensity matrix,  $I$

	$g_1$	$g_2$	$g_3$	$g_4$	$g_5$	$g_6$	$g_7$	$g_8$	$g_9$	$g_{10}$
$g_1$	0.0	8.1	9.2	7.7	9.3	2.3	5.1	10.2	6.1	7.0
$g_2$	8.1	0.0	12.0	0.9	12.0	9.5	10.1	12.8	2.0	1.0
$g_3$	9.2	12.0	0.0	11.2	0.7	11.1	8.1	1.1	10.5	11.5
$g_4$	7.7	0.9	11.2	0.0	11.2	9.2	9.5	12.0	1.6	1.1
$g_5$	9.3	12.0	0.7	11.2	0.0	11.2	8.5	1.0	10.6	11.6
$g_6$	2.3	9.5	11.1	9.2	11.2	0.0	5.6	12.1	7.7	8.5
$g_7$	5.1	10.1	8.1	9.5	8.5	5.6	0.0	9.1	8.3	9.3
$g_8$	10.2	12.8	1.1	12.0	1.0	12.1	9.1	0.0	11.4	12.4
$g_9$	6.1	2.0	10.5	1.6	10.6	7.7	8.3	11.4	0.0	1.1
$g_{10}$	7.0	1.0	11.5	1.1	11.6	8.5	9.3	12.4	1.1	0.0

(b) Distance matrix,  $d$



(c) Expression patterns as points in three-dimensional space.

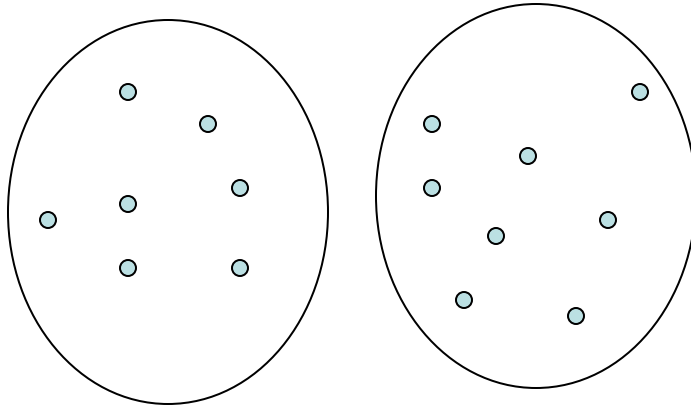
# Clustering

- Find groups of genes that have similar expression profiles to one another
- Such groups may be functionally related, and/or co-regulated
- Compute pairwise distance metric  $d(i,j)$  for every pair of genes  $i$  and  $j$
- This gives an  $n \times n$  “distance matrix”  $\mathbf{d}$

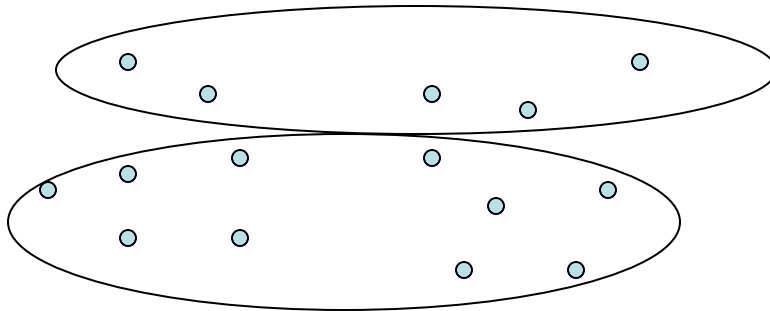
# Goal of clustering

- To group together genes into clusters such that
  - Genes within a cluster have highly similar expression profiles (small  $d(i,j)$ ): “**homogeneity**”
  - Genes in different clusters have very different expression profiles (large  $d(i,j)$ ): “**separation**”
- “Good” clustering is one that adheres to these goals
- A really “good” clustering is decided by biological interpretation of the clusters

The Points are in some multi-dimensional space



Good Clustering



Bad Clustering



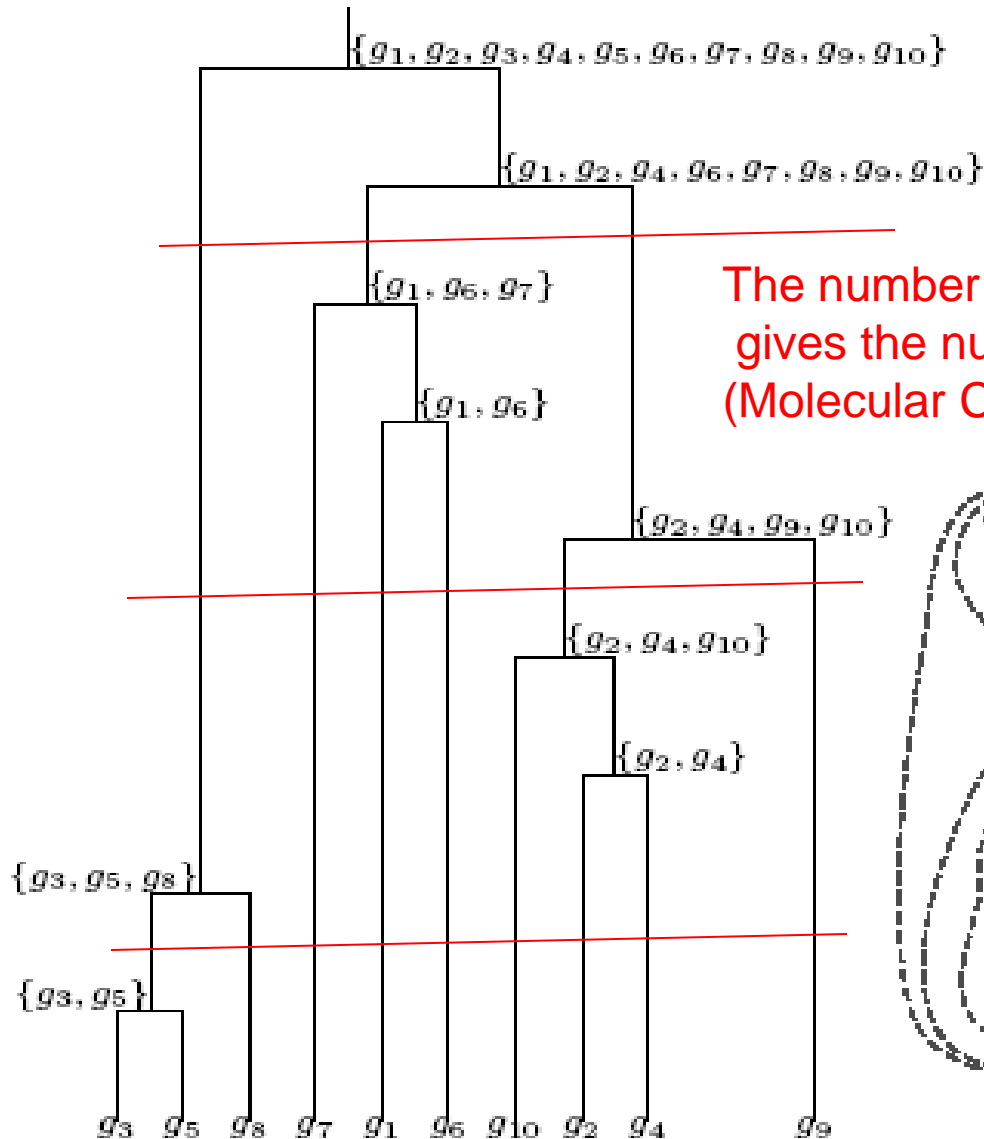
# Clustering problems

- How to measure distance/similarity ?
- How many clusters ?
- Very large data sets: ~10,000 gene, ~100 conditions create computational difficulties

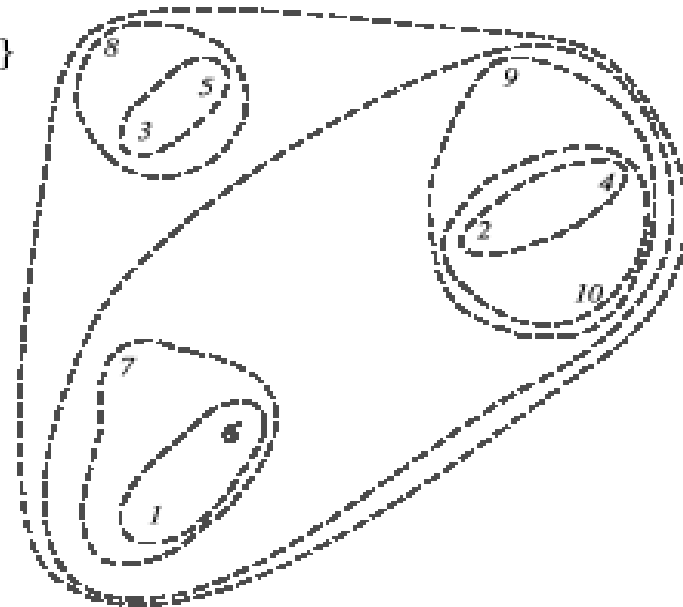
# Hierarchical clustering

- One approach to clustering
- Does not explicitly partition genes into groups
- Organizes genes into a tree; genes are at the leaves of the tree
- Edges have lengths
- Total path length between two genes (leaves) correlates with the distance between the genes

# Hierarchical Clustering



The number of vertical lines intersecting a red line gives the number of sub-clusters at that point.  
(Molecular Clock : length of vertical lines)



# Hierarchical Clustering Algorithm

1. Hierarchical Clustering ( $d, n$ )
2. Form  $n$  clusters each with one element
3. Construct a graph  $\mathcal{T}$  by assigning one vertex to each cluster
4. **while** there is more than one cluster
5.     Find the two closest clusters  $C_1$  and  $C_2$
6.     Merge  $C_1$  and  $C_2$  into new cluster  $C$  with  $|C_1| + |C_2|$  elements
7.     **Compute distance from  $C$  to all other clusters**
8.     Add a new vertex  $C$  to  $\mathcal{T}$  and connect to vertices  $C_1$  and  $C_2$
9.     Remove rows and columns of  $d$  corresponding to  $C_1$  and  $C_2$
10.     Add a row and column to  $d$  corresponding to the new cluster
11.      $C$   
return  $\mathcal{T}$

**Different ways to define distances between clusters may lead to different clusterings**

# Hierarchical Clustering: Recomputing Distances

- $d_{min}(C, C^*) = \min d(x,y)$   
*for all elements  $x$  in  $C$  and  $y$  in  $C^*$* 
  - Distance between two clusters is the **smallest** distance between any pair of their elements
- $d_{avg}(C, C^*) = (1 / |C^*||C|) \sum d(x,y)$   
*for all elements  $x$  in  $C$  and  $y$  in  $C^*$* 
  - Distance between two clusters is the **average** distance between all pairs of their elements

# *K*-means clustering

- Another popular solution to the clustering problem
- Guess a number  $k$ , which is the number of clusters that will be reported
- Finds explicit clusters (unlike hierarchical clustering)

# The Problem

- Given a set of  $n$  data points  $V = (v_1, v_2, \dots, v_n)$ 
  - Each data point is a vector in an  $m$ -dimensional space and an integer  $k$ , the problem is to determine a set of  $k$  points or *centers* in  $m$ -dimensional space that minimizes the squared error distortion as defined below.
- Let  $X = (x_1, x_2, \dots, x_k)$  be a set of  $k$  points in the same vector space
- For each  $v_i$ , find  $x \in X$  that is closest to it, i.e., the Euclidian distance  $d(v_i, x)$  is least
- Sum the square of this Euclidian distance, over all  $v_i$

- Formally,  $d(v_i, X) = \min_{x \in X} d(v_i, x)$

$$d(V, X) = \frac{\sum_{i=1}^n d(v_i, X)^2}{n}$$

# Objective function

- Find  $X$  such that  $d(V, X)$  is minimized
- **Input:** A set,  $V$ , consisting of  $n$  points and a parameter  $k$
- **Output:** A set  $X$  consisting of  $k$  points (*cluster centers*) that minimizes the  $d(V, X)$  over all possible choices of  $X$
- An NP-complete problem for  $k > 1$



# K-Means Clustering: Lloyd Algorithm

1. Lloyd Algorithm
2. Arbitrarily assign the  $k$  cluster centers
3. **while** the cluster centers keep changing
4. Assign each data point to the cluster  $C_i$  corresponding to the closest cluster representative (center) ( $1 \leq i \leq k$ )
5. After the assignment of all data points, compute new cluster representatives according to the center of gravity of each cluster, that is, the new cluster representative is  
$$\Sigma v / |C| \text{ for all } v \text{ in } C \text{ for every cluster } C$$

\*This may lead to merely a locally optimal clustering.

# Conservative K-Means Algorithm

- Lloyd algorithm is fast but in each iteration it moves many data points, not necessarily causing better convergence.
- A more conservative method would be to move one point at a time only if it improves the overall **clustering cost**
  - The smaller the clustering cost of a partition of data points is, the better that clustering is
  - Different methods (e.g.,  $d(V, X)$  we saw earlier) can be used to measure this clustering cost

# K-Means “Greedy” Algorithm

1. ProgressiveGreedyK-Means( $k$ )
2. Select an arbitrary partition  $P$  into  $k$  clusters
3. **while** forever
4.      $bestChange \leftarrow 0$
5.     **for** every cluster  $C$
6.         **for** every element  $i$  not in  $C$
7.             **if** moving  $i$  to cluster  $C$  reduces its clustering cost
8.                 **if**  $(cost(P) - cost(P_{i \rightarrow C}) > bestChange$
9.                      $bestChange \leftarrow cost(P) - cost(P_{i \rightarrow C})$
10.                      $i^* \leftarrow i$
11.                      $C^* \leftarrow C$
12.             **if**  $bestChange > 0$
13.                 Change partition  $P$  by moving  $i^*$  to  $C^*$
14.             **else**
15.                 **return**  $P$

# UPGMA: Unweighted Pair Group Method with Arithmetic Mean

- UPGMA is a clustering algorithm that:
  - computes the distance between clusters using average pairwise distance
  - assigns a *height* to every vertex in the tree
- Does not require an additive distance matrix. Output tree does not necessarily match the distance matrix for every pair of nodes

# Clustering in UPGMA

- Given two disjoint clusters  $C_i$ ,  $C_j$  of sequences,

$$d_{ij} = \frac{1}{|C_i| \times |C_j|} \sum_{\{p \in C_i, q \in C_j\}} d_{pq}$$

- Algorithm is a variant of the hierarchical clustering algorithm

# UPGMA Algorithm

UPGMA ( $D, n$ )

1. Form  $n$  clusters each with a single element
2. Construct a graph  $T$  by assigning one vertex to each cluster
3. Assign height  $h(v)=0$  to every vertex  $v$  in the graph
- 4 **while** there is more than one cluster
5. Find the two closest clusters  $C_1$  and  $C_2$
6. Merge  $C_1$  and  $C_2$  into new cluster  $C$  with  $|C_1| + |C_2|$  elements
7. **for every cluster**  $C^* \neq C$  **compute**

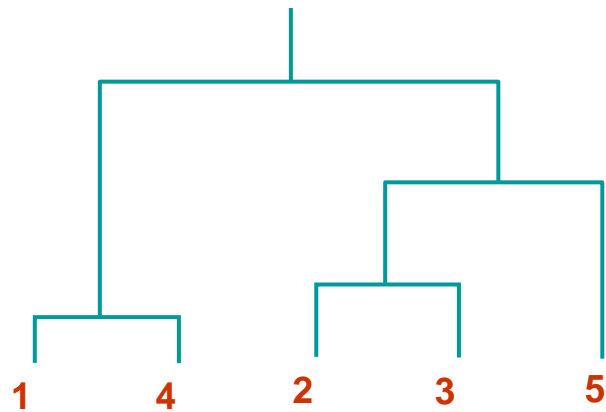
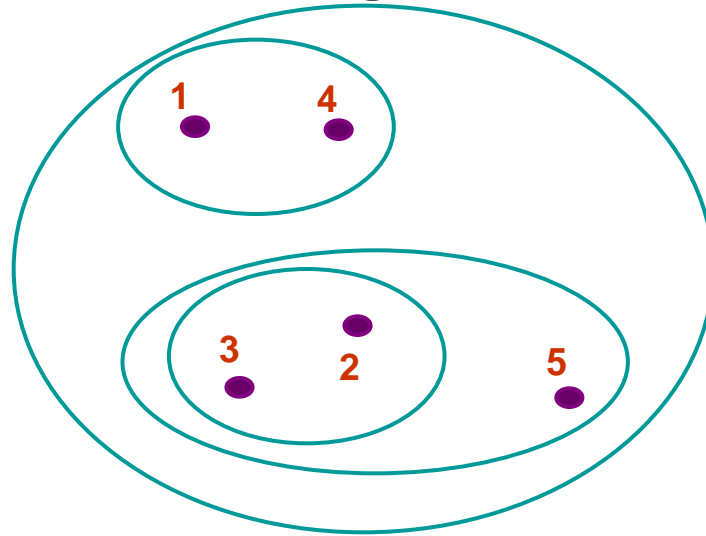
$$D(C, C^*) = \frac{1}{|C| |C^*|} \sum_{i \text{ in } C} \sum_{j \text{ in } C^*} D(i, j)$$

8. Add a new vertex  $C$  to  $T$  and connect to vertices  $C_1$  and  $C_2$

9. 
$$h(C) \leftarrow \frac{D(C_1, C_2)}{2}$$

10. Assign length  $h(C) - h(C_1)$  to edge  $(C_1, C)$
11. Assign length  $h(C) - h(C_2)$  to edge  $(C_2, C)$
12. Remove rows and columns of  $D$  corresponding to  $C_1$  and  $C_2$
13. Add a row and column to  $D$  corresponding to the new cluster  $C$
14. return  $T$

# UPGMA Algorithm (cont'd)



# UPGMA's Weakness

- The algorithm produces an **ultrametric** tree : the distance from the root to any leaf is the same
- UPGMA assumes a constant molecular clock: all species represented by the leaves in the tree are assumed to accumulate mutations (and thus evolve) at the same rate. This is a major pitfall of UPGMA.



# Neighbor Joining Algorithm

The neighbor joining algorithm works well with additive trees, does not assume any molecular clock and generates clusters that satisfy homogeneity and separation criteria .  
For a cluster  $C$  , define  $u(C)$  as

$$u(C) = \frac{1}{\text{number of clusters} - 2} \sum_{\text{all clusters } C'} D(C, C')$$

To choose which two clusters to merge, the algorithm minimizes  $D(C_1, C_2) - u(C_1) - u(C_2)$

(The theory behind this algorithm is complex and is left out here.)

# NeighborJoining ( $D, n$ )

1. Form  $n$  clusters, each with a single element
2. Construct a graph  $T$  by assigning a vertex to each cluster
3. **while** there is more than one cluster
4. Find clusters  $C_1, C_2$  minimizing  $D(C_1, C_2) - u(C_1) - u(C_2)$
5. Merge  $C_1, C_2$  into a new cluster  $C$  of size  $|C_1| + |C_2|$
6. Compute  $D(C, C^*) = \frac{D(C_1, C) + D(C_2, C)}{2}$  to every other cluster  $C^*$
7. Add vertex  $C$  to  $T$  and connect to vertices  $C_1$  and  $C_2$
8. Assign length  $\frac{1}{2}D(C_1, C_2) + \frac{1}{2}(u(C_1) - u(C_2))$  to edge  $(C_1, C)$
9. Assign length  $\frac{1}{2}D(C_1, C_2) + \frac{1}{2}(u(C_2) - u(C_1))$  to edge  $(C_2, C)$
10. Remove rows and columns of  $D$  corresponding to  $C_1$  and  $C_2$
11. Add a row and column to  $D$  for the new cluster  $C$
12. **Return**  $T$

# Sources

- <http://www.math.tau.ac.il/~rshamir/ge/02/scribes/lec01.pdf>
- <http://bioinformatics.oupjournals.org/cgi/screenpdf/20/3/340.pdf>
- [http://www.absoluteastronomy.com/encyclopedia/M/Mi/Minimum\\_spanning\\_tree.htm](http://www.absoluteastronomy.com/encyclopedia/M/Mi/Minimum_spanning_tree.htm)
- Serafim Batzoglou (UPGMA slides)  
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