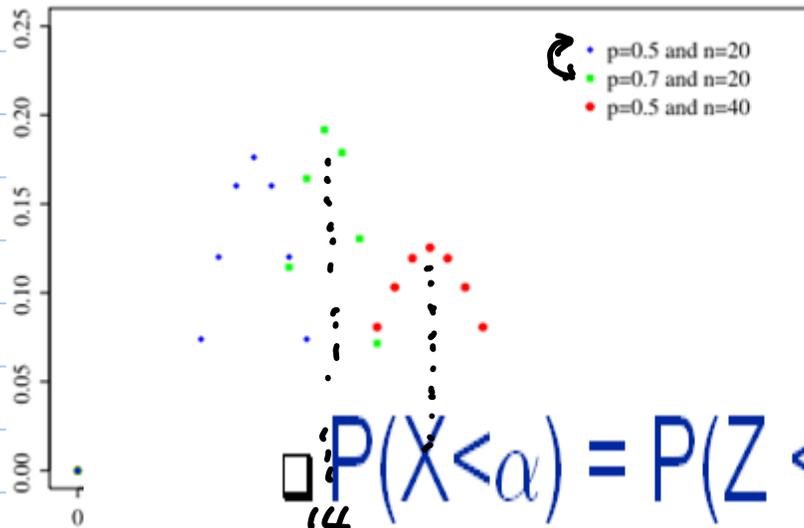


Note

$$P(X = k) \equiv f(k; n, p) = \binom{n}{k} p^k (1-p)^{n-k}$$

$P(A|B) = P(A) \cdot P(B)$   
 when A, B independent



$$E[X] = np$$

$P(Q) = P(\leq e \text{ bit error in } n \text{ bit trials})$

$$P(X < \alpha) = P(Z < \frac{\alpha - \mu}{\sigma}) = P(X=0) + P(X=1) + \dots + P(X=e)$$

v.v.  $X$ : # of error bits in  $n$  bits transmission

$$= \Phi\left(\frac{\alpha - \mu}{\sigma}\right)$$

□ 1 Mb/s link

□ each user:

- 100 kb/s when "active"
- active 10% of time

$$n = 35$$

r.v.  $X$ : # of active users  
at a moment

$$X \sim B(n, p)$$

$$\sim B(35, 0.1)$$

$$P(\text{congestion}) = P(\text{more than 10 users active})$$

$$= P(X > 10)$$

$$= P(X=11) + P(X=12) + P(X=13) + \dots + P(X=35)$$

$$\approx 0.0004$$

$$\begin{aligned} \square P(X < \alpha) &= P(Z < (\alpha - \mu) / \sigma) \\ &= \Phi((\alpha - \mu) / \sigma) \end{aligned}$$

$$Z = \frac{X - \mu}{\sigma}$$

$$\begin{aligned} P(X < \alpha) &= P(Z\sigma + \mu < \alpha) \\ &= P(Z\sigma < \alpha - \mu) \\ &= P\left(Z < \frac{\alpha - \mu}{\sigma}\right) \end{aligned}$$

